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Ambiguity Attitudes and Self-Confirming Equilibrium in Sequential Games*

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Abstract

We consider a game with sequential moves played by agents who are randomly drawn from large populations and matched. We assume that, when players are uncertain about the strategy distributions of the opponents, preferences over actions at any information set admit a smooth-ambiguity representation in the sense of Klibanoff, Marinacci, and Mukerji (*Econometrica*, 2005). This may induce dynamically inconsistent preferences and calls for an appropriate definition of sequential best response. We take this into account in our analysis of self-confirming equilibrium (SCE) and rationalizable SCE in sequential games with feedback played by agents with non-neutral ambiguity attitudes. Battigalli, Cerreia-Vioglio, Maccheroni, and Marinacci (*Amer. Econ. Rev.*, 2015) show that the set of SCE's of a simultaneous-move game with feedback expands as ambiguity aversion increases. We show by example that SCE in a sequential game is not equivalent to SCE applied to the strategic form of such game, and that the previous monotonicity result does not extend to general sequential games. Still, we provide sufficient conditions under which the monotonicity result holds for (rationalizable) SCE.

KEYWORDS: Sequential games with feedback, smooth ambiguity, self-confirming equilibrium, rationalizable self-confirming equilibrium.

1 Introduction

Self-confirming equilibrium When a game is played recurrently and the learning dynamic has reached a rest point, each agent chooses a best reply to his subjective belief,

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which may be incorrect and yet confirmed by the evidence available to him. A profile of strategies and beliefs with this property is a self-confirming equilibrium (henceforth, SCE).¹ The standard definition of SCE assumes that agents are subjective expected utility maximizers, i.e., that they are ambiguity neutral.² Yet, a large body of empirical evidence supports the ambiguity aversion hypothesis. This is particularly relevant when agents only have scarce evidence of opponents' behavior and face therefore strategic uncertainty. Following Battigalli, Cerreia-Vioglio, Maccheroni and Marinacci (2015, henceforth BCMM), we analyze the SCE concept in games played by ambiguity averse agents. Unlike BCMM, who essentially restrict their attention to simultaneous-move games, we consider games with sequential moves, represented in extensive form. In a sequential game, agents have evidence, at best, of how opponents play on the equilibrium path, but no evidence of how they would play off the path. Thus, sequential games constitute a natural context for self-confirming equilibrium analysis. In the rest of this introduction, we describe how our paper adds to the previous literature in general and to BCMM in particular.

SCE and model uncertainty BCMM analyze SCE in simultaneous-move population games played recurrently by agents with non-neutral attitudes toward ambiguity, which is the imperfect quantifiability of the relevant risks. Specifically, agents are assumed to have smooth-ambiguity preferences in the sense of Klibanoff, Marinacci, and Mukerji (2005, henceforth KMM). This decision model is flexible and analytically convenient for game theoretic applications. First, it separates ambiguity attitudes, a stable personal trait like risk attitudes, from the perception of uncertainty, which is a property of subjective beliefs affected by the game situation; second, it provides a parameterization of ambiguity aversion analogous to the parameterization of risk aversion, which simplifies comparative static exercises.³

SCE is defined as follows: Let I denote the set of player roles (e.g., buyer and seller) and S_i the set of pure strategies of any agent playing in role i . A profile of strategy distributions $\sigma^* = (\sigma_i^*)_{i \in I} \in \times_{i \in I} \Delta(S_i)$ is a SCE if, for each $i \in I$ and each $s_i^* \in S_i$ with $\sigma_i^*(s_i^*) > 0$, there is a belief μ_i about the strategy distributions σ_{-i} of the opponents that justifies s_i^* as a KMM-best response and is consistent with the long-run distribution of ex-post observations for i generated by s_i^* and σ_{-i}^* (e.g., the (s_i^*, σ_{-i}^*) -induced distribution of terminal nodes). Since the distribution of observations may not reveal the true underlying distribution of strategies σ_{-i}^* , agents may be uncertain about it.

Assuming that the own-payoff relevant consequences (e.g., one's own monetary payoff) are observed by each agent after each play, BCMM prove a monotonicity result: higher ambiguity aversion entails a larger set of equilibria. Intuitively, for each agent, the strategy played repeatedly in equilibrium yields known risks because the agent observes its long-run

¹This terminology is due to Fudenberg and Levine (1993).

²For a perspective on ambiguity attitudes that applies to this paper, see Cerreia-Vioglio et al. (2013) and the survey by Marinacci (2015).

³See the survey by Marinacci (2015).

distribution of payoffs, while deviations are untested and may be perceived as ambiguous; therefore, higher ambiguity aversion penalizes deviations relative to the equilibrium choice. This monotonicity result implies that greater ambiguity aversion entails less predictability of strategies in the long-run, because the set of possible steady states is larger.

Our contribution The scope of BCMM’s analysis is essentially limited to simultaneous-move games and, possibly, games played in strategic form—such as experimental games played in the lab with the so-called “strategy method”⁴—because it is well known that ambiguity aversion may make preferences over strategies dynamically inconsistent (e.g., Siniscalchi 2011). As a consequence, there would be incentives to make covert commitments if such commitment moves were available.⁵ Thus, the fact that agents in sequential games cannot (irreversibly) choose strategies must be faced and dealt with explicitly. We assume that agents are sophisticated: They understand their future contingent incentives and choose actions in early stages predicting that such incentives determine their actions in later stages, that is, they plan by “folding back” given their subjective conditional beliefs about the strategies of other players. Hence, they execute “unimprovable” strategies. Since unimprovable strategies may be different from best replies in the normal form of the game, the definition of equilibrium due to BCMM cannot be applied to sequential games.

Two questions naturally arise: First, how does SCE defined on the extensive form relate to SCE in the normal form of the same game? Second, does the comparative ambiguity aversion result extend to games with sequential moves? To elaborate on the first question, fix a sequential game Γ represented in extensive form with a set of paths (terminal nodes) Z and a specification of players’ feedback $f = (f_i)_{i \in I}$ that describes what they can observe at the end of each play; formally, each f_i is a function defined on Z . Then, we can derive the normal (or strategic) form $(G, F) = \mathcal{N}(\Gamma, f)$, where G is given by the normal-form payoff functions, and each player i ’s feedback f_i about the path is replaced by a corresponding normal-form feedback F_i defined on the set $S = \times_{i \in I} S_i$ of pure strategy profiles (e.g., each player only observes his monetary payoff, which is a function of $s \in S$). Under subjective expected utility maximization, which is dynamically consistent, SCE in the extensive form is realization equivalent to SCE in the normal form. With ambiguity aversion, instead, we show that there may be different, non-nested sets of equilibrium outcomes; in other words, working with the normal form is neither too permissive nor too restrictive, it is just wrong.

Given that the sets of SCE outcomes in the extensive and normal form do not coincide, we cannot rely on the monotonicity result of BCMM to argue that, in a game with sequential moves, the set of SCE outcomes expands as ambiguity aversion increases. It is still true that, on the equilibrium path, equilibrium actions entail known risks while deviations may be perceived as ambiguous. But, if a deviation is followed by other actions of the same

⁴See the survey by Brandts and Charness (2011) and the references therein.

⁵Unlike *overt* commitment, covert commitment moves are not observed by other players. The strategic advantages of overt commitment are well known at least since Schelling (1960) and do not depend on dynamic inconsistency as traditionally defined in decision theory.

player, folding-back planning may require that plans (i.e., predictions) about these actions change with ambiguity aversion, and dynamic inconsistency may lead to an increase in the value of deviations. Indeed, as we demonstrate in examples, there may be instances of non-monotonicity when the SCE strategies at some baseline level of ambiguity aversion are such that some agents at some reachable information sets would be willing to pay to commit—if they only could—on a different course of action involving multiple sequential deviations from the folding-back plan. We prove a preliminary result, Lemma ??, stating that such willingness to commit is necessary for non-monotonicity.

Lemma ?? allows us to prove the monotonicity result for two special cases: (i) games where, on each path, no player moves more than once, and (ii) pure strategy equilibria of games with no chance moves.⁶ In case (i) there cannot be multiple sequential deviations from folding-back planning. In case (ii) we show that, loosely speaking, in equilibrium ambiguity aversion is not distinguishable from the risk aversion of expected utility maximizers, who are dynamically consistent. Lemma 1 also implies the following noteworthy result: The set of SCE outcomes with ambiguity neutral agents is always included in the set of SCE outcomes with ambiguity averse agents. This means that the standard version of the SCE concept, by ignoring ambiguity aversion, overestimates the predictability of long-run outcomes in recurrent interactions.

Next, we turn to the well known issue of the impossibility of *overt* (i.e., observable) commitment and how this affects strategic reasoning and equilibrium. Rationality in dynamic games requires that agents choose subjective best replies at *all* information sets, including the unexpected ones, given beliefs revised upon observing unexpected moves. We call this condition “full unimprovability.” If SCE is not refined so as to capture strategic reasoning, requiring full unimprovability rather than simple unimprovability does not change the set of equilibrium outcomes. The reason is that in an SCE agents may hold wrong beliefs about the reactions of others to non-equilibrium moves, which are necessarily unexpected; hence, they can expect responses that are irrational given the actual payoff functions of co-players. Assuming instead that (some features of) such payoff functions are common knowledge, one can embed strategic reasoning into the SCE concept: if (a) players are rational, (b) their beliefs are confirmed, and (c) there is common belief of (a) and (b), then their strategies form a “rationalizable SCE.” This concept has been analyzed under the assumption of subjective expected utility maximization, that is, neutral ambiguity attitudes (e.g., Rubinstein and Wolinsky 1994, Dekel et al. 1999).⁷ Here we provide an extension for non-neutral ambiguity attitudes and prove version of the monotonicity result

⁶We call such equilibria “symmetric” because in the population game scenario they represent situations where all agents in the same role play in the same way.

⁷Battigalli (1987) and Battigalli and Guaitoli (1988) consider a weaker concept of SCE in rationalizable strategies justified by the following assumptions: (a) players are rational, (b) their beliefs are confirmed, and (c) there is common belief of (a) only. To the best of our knowledge, unlike plain SCE, there is no learning foundation of rationalizable SCE. On the other hand, we can give a kind of learning foundation of SCE in rationalizable strategies.

for rationalizable SCE. The analysis involves some delicate technical details.

Related literature This is the first paper that analyzes SCE with non-neutral ambiguity attitudes in sequential games. BCMM and Battigalli et al. (2016b) only consider games played in strategic form. The extensive form game from which the strategic form is derived only affects feedback: two strategy profiles that yield the same terminal node necessarily yield the same feedback to every player. Furthermore, Battigalli et al. (2016b) focus on “maxmin SCE” a notion of equilibrium with extreme ambiguity aversion (see Gilboa and Schmeidler 1989), whereas here—like BCMM—we apply the smooth ambiguity criterion of KMM. Hereafter we call “standard SCE” the self-confirming equilibrium concept with ambiguity neutrality, i.e., subjective expected utility maximization.

A version of the standard SCE concept was first put forward in the undergraduate thesis of Battigalli (1987)⁸ and called “conjectural equilibrium.” The first definition in English appears in a working paper by Battigalli and Guaitoli (1988), who also define and apply an extensive-form rationalizability refinement. Fudenberg and Levine (1993) independently put forward a definition of standard SCE under the maintained assumption that agents perfectly observe ex-post the path of play. Unlike Battigalli (1987) and Battigalli and Guaitoli (1988), the analysis of Fudenberg and Levine applies to large population games. This is important because it justifies the assumption that agents maximize their short-run expected utility even if they are patient. The reason is that in a large population game there are no incentives to affect the future behavior of current opponents, who are almost surely different from the future ones.⁹ For further references and details see the discussion in Section IV of BCMM and the survey by Battigalli et al. (1992).

Fudenberg and Levine (1993) note that standard SCE is not strategic-form invariant, arguing that the strategic—or normal—form is therefore insufficient to characterize SCE. But their comment rests on the maintained assumption that players always observe ex-post the path of play, which in a simultaneous-move game is just the actual profile of strategies (actions of the strategic form). Therefore, when they compare standard SCE in the extensive (i.e., sequential) and strategic form, they change what players can observe ex post about the behavior of others: only on-path actions in the extensive form, and complete strategy profiles in the strategic form. If feedback about the behavior of others is instead kept constant—as we do when we look at strategic-form feedback—standard SCE in the extensive form is equivalent to standard SCE in the strategic form (see Remark ??). On the other hand, we show that SCE with ambiguity aversion is *not* strategic-form invariant (see Example 4).

As mentioned above, BCMM proved that in simultaneous-move games the SCE correspondence is monotone with respect to players’ degree of ambiguity aversion. Here we prove partial versions of this result for sequential games. A similar monotonicity result can

⁸Written in Italian.

⁹Furthermore, incentives to experiment vanish in the long run.

be proved for comparative risk aversion in the particular case of *pure* standard SCE (see the lecture notes of Battigalli 2017), and for comparative risk or ambiguity aversion in the case of rationalizable strategies (Battigalli et al. 2016a). Note that the intuition and proof in the case of rationalizability are very different from the SCE case. Interestingly, we rely on Battigalli et al. (2016a) to prove monotonicity of the rationalizable SCE correspondence in some special cases.

Next we compare our contribution to other papers on equilibria with ambiguity aversion in sequential games. Unlike us, these papers analyze versions or refinements of the Nash equilibrium concept. Lo (1999) analyzes a notion of Nash equilibrium whereby players are mini-maximizers given belief sets in the sense of Gilboa and Schmeidler (1989). Assuming Bayesian updating¹⁰ for each measure in each player’s belief set, preferences at different information sets may be dynamically inconsistent. As Lo points out, this implies that the proposed equilibrium concept is not strategic-form invariant. We differ from Lo (1999) in several ways: we adopt the smooth ambiguity criterion, we analyze a version of SCE rather than Nash equilibrium, and we define rational planning as unimprovability (folding back), whereas Lo imposes a form of sequential—or interim—optimality (mini-maximization with respect to continuation strategies at each reachable information set). Hanany et al. (2017) use the smooth ambiguity criterion to analyze sequential games with incomplete information. Their work differs from ours in two important ways. First, they analyze a notion of perfect Bayesian equilibrium, which assumes that players’ conjectures about the (type-dependent) strategies of other players are correct, while beliefs about the distribution of types are not disciplined; we instead allows for incorrect conjectures about strategies, but the confirmed-beliefs requirement disciplines beliefs about types (cf. Dekel et al 2004). Second, they model subjectively rational planning in a different way. We maintain that players’ beliefs satisfy the standard rules of conditional probability, which imply Bayes rule; given this, each player adopts an unimprovable strategy that—by the dynamic inconsistency of ambiguity averse preferences—may be suboptimal according to his initial preferences, or the interim preferences he holds at some information set. Hanany et al. instead assume that each player’s equilibrium strategy is sequentially optimal, that is, optimal according to his preferences at each information set. Ambiguity aversion implies that this cannot always be reconciled with Bayesian updating. However, they show it is without loss of generality to restrict attention to beliefs generated using a particular generalization of Bayesian updating, where more ambiguity aversion implies larger deviations from standard Bayesian updating. Eichberger et al. (2017) analyze a notion of equilibrium with ambiguity aversion in sequential games that combines Choquet-expected utility maximization (Schmeidler 1989) with a generalized form of Bayesian updating. Unlike Lo (1999) and Hanany et al (2017), and like us, they assume folding-back planning.

¹⁰Or maximum likelihood updating.

Outline The rest of the paper is organized as follows. Sections 2 and 3 introduce the setup and the smooth ambiguity criterion; Sections 4 and 7 respectively discuss unimprovability and full unimprovability; Section 5 defines our SCE concept; Section 6 presents comparative results for SCE; Section 8 analyzes rationalizable SCE; all proofs are collected in the Appendix.

2 Framework

We analyze an agent with non-neutral ambiguity attitudes who plays a game with sequential moves. We assume that the commitment technology of this agent is explicitly represented by the rules of the game. Therefore, the agent can control — i.e., irreversibly choose — only his (pure) actions at whatever information set is being reached. We also assume that he is sophisticated and therefore he takes this into account when he plans how to play the game.

We take the point of view of an agent who plays in role $i \in I$ of a *finite extensive-form game* Γ and has *perfect recall*. Let H_i denote the collection of **information sets** of i and let $A_i(h)$ be the set of actions available at $h \in H_i$. We assume for expositional simplicity that $|A_i(h)| \geq 2$ for each $h \in H_i$, where $|X|$ denotes the cardinality of a finite set X . This means that we include in H_i only the information sets where i is active. Let \emptyset denote the root of the game, then $\{\emptyset\} \in H_i$ if and only if i is a first mover.¹¹ We endow H_i with the weak (respectively, strict) precedence relation \preceq (\prec) inherited from the game tree.¹² The set of **strategies** for player i is $S_i = \times_{h \in H_i} A_i(h)$. For every $s_i \in S_i$ and $h \in H_i$, we let $s_{i,h}$ denote the action specified by s_i at h ; thus, $s_i = (s_{i,h})_{h \in H_i} \in \times_{h \in H_i} A_i(h)$.

We model randomization explicitly as the choice of a randomization device. Therefore it is important to allow for **chance moves** as the moves of a special player denoted by $0 \notin I$. With this, H_0 denotes the collection of information sets of the chance player, $A_0(h)$ is set of chance moves available at $h \in H_0$, and $S_0 = \times_{h \in H_0} A_0(h)$ is the set of “strategies” of the pseudo-player 0. Throughout, we maintain for simplicity the assumption that the *probabilities of chance moves are commonly known*. Such probabilities are specified by a “behavioral strategy” $\beta_0 \in \times_{h \in H_0} \Delta(A_0(h))$, with $\beta_0(a_0|h) > 0$ for every $h \in H_0$ and

¹¹Our preferred representation of games in extensive form starts from sequences of action profiles, that correspond to the nodes of the game tree (e.g., Chapters 6 and 11 of Osborne and Rubinstein, 1994) and allows for the representation of players’ information also at nodes where they are not active, such the root for players who are not first-movers. This affects the way we draw pictures and describe examples, but it is otherwise irrelevant for the analysis of the paper.

¹²Perfect recall implies that, for all $h, h' \in H_i$, there are nodes (histories) $x \in h$ and $x' \in h'$ such that x precedes x' if and only if every node of h' is preceded by a node of h . With this, we can stipulate that, for all $h, h' \in H_i$, h strictly **precedes** h' , written $h \prec h'$, if every node of h' is strictly preceded by a node of h . Perfect recall implies that each $h \in H_i$ can have at most one immediate predecessor. The reflexive closure of \prec is \preceq , an antisymmetric, and transitive relation that makes H_i a directed forest. If $\{\emptyset\} \in H_i$, then (H_i, \preceq) is a directed tree.

$a_0 \in A_0(h)$; β_0 induces the “mixed strategy” $\sigma_0 \in \Delta(S_0)$ such that

$$\sigma_0(s_0) = \prod_{h \in H_0} \beta_0(s_{0,h}|h) > 0$$

for every $s_0 \in S_0$. Thus, the outcome distributions respectively induced by β_0 and σ_0 coincide for every strategy profile of the true players (see Kuhn, 1953). Since players are always certain that the “mixed strategy of chance” is σ_0 , we model explicitly only each player i 's beliefs about the true opponents $-i = I \setminus \{i\}$. We let $S = \times_{j \in I} S_j$ denote the set of pure strategy profiles of the true players, whereas $S_{-i} = \times_{j \in I \setminus \{i\}} S_j$ and $S_{0,-i} = S_0 \times S_{-i}$ denote the set of pure strategy profiles of opponents respectively excluding and including chance.

Let Z denote the set of terminal nodes of the game. Every profile (s_0, s) induces a complete path, hence a terminal node, through the **outcome function**

$$\zeta : S_0 \times S \rightarrow Z.$$

Since the definition of ζ is standard, we take it for granted and then define some derived concepts using ζ .

For conceptual clarity, we also include in the description of the extensive form Γ a **consequence function**

$$\gamma : Z \rightarrow C$$

which specifies the material consequence $c = \gamma(z) \in C$ of each terminal node $z \in Z$. For example, we may have $C \subseteq \mathbb{R}^I$ where $c = (c_j)_{j \in I}$ is a consumption allocation or a distribution of monetary payoffs to players. Thus, player i 's risk attitudes (preferences over objective lotteries of consequences) are represented by a **vNM utility function**

$$v_i : C \rightarrow \mathbb{R}.$$

It is convenient to specify the information about strategies implied by any $h \in H_i$. Let $S_{0,I}(h)$ denote the set of strategy profiles (s_0, s) reaching h , formally

$$S_{0,I}(h) = \{(s_0, s) \in S_0 \times S : \exists x \in h, x \prec \zeta(s_0, s)\}.$$

Then, the set of pure strategy profiles of chance and opponents that allow for h is

$$S_{0,-i}(h) = \text{proj}_{S_{0,-i}} S_{0,I}(h) = \{s_{0,-i} \in S_{0,-i} : \exists (x, s_i) \in h \times S_i, x \prec \zeta(s_i, s_{0,-i})\}.$$

Similarly, the sets of i 's and $-i$'s strategies allowing for h are, respectively,

$$S_i(h) = \text{proj}_{S_i} S_{0,I}(h) = \{s_i \in S_i : \exists (x, s_{0,-i}) \in h \times S_{0,-i}, x \prec \zeta(s_i, s_{0,-i})\}$$

and

$$S_{-i}(h) = \text{proj}_{S_{-i}} S_{0,I}(h) = \{s_{-i} \in S_{-i} : \exists (x, s_{0,i}) \in h \times S_0 \times S_i, x \prec \zeta(s_{0,i}, s_{-i})\}.$$

It is useful to keep in mind that perfect recall implies the following factorization:

$$\forall h \in H_i, S_{0,I}(h) = S_i(h) \times S_{0,-i}(h).$$

Furthermore, it also implies that

$$\forall g, h \in H_i, g \prec h \Rightarrow S_{0,-i}(h) \subseteq S_{0,-i}(g).$$

Intuitively, i obtains finer information about strategies as the play unfolds.

Each agent playing in role i knows that his opponents are drawn at random from large populations $j \in I \setminus \{i\}$ of agents, with each agent playing a pure strategy. The **distribution of pure strategies in population j** is some *unknown* measure $\sigma_j \in \Delta(S_j)$, hence, by random matching, the *objective* probability of facing opponents playing pure strategy profile $s_{-i} = (s_j)_{j \in I \setminus \{i\}}$ is the unknown product¹³

$$\sigma_{-i}(s_{-i}) = \prod_{j \in I \setminus \{i\}} \sigma_j(s_j).$$

To ease notation, we identify each profile of distributions $(\sigma_j)_{j \in I \setminus \{i\}}$ with the corresponding product distribution on S_{-i} . We let

$$\Sigma_{-i} = \left\{ \sigma_{-i} \in \Delta(S_{-i}) : \exists (\sigma_j)_{j \in I \setminus \{i\}} \in \times_{j \in I \setminus \{i\}} \Delta(S_j) : \sigma_{-i} = \times_{j \in I \setminus \{i\}} \sigma_j \right\}$$

denote the set of these product distributions.¹⁴ We endow Σ_{-i} with the standard topology inherited from the Euclidean topology on $\mathbb{R}^{S_{-i}}$, which makes it compact, and with the Borel sigma algebra $\mathcal{B}(\Sigma_{-i})$ generated by the standard topology.

At each point of the game, the agent playing in role i has some **belief** $\mu_i \in \Delta(\Sigma_{-i})$. The belief μ_i that i holds at the beginning of the game is i 's **prior**. For each $\mu_i \in \Delta(\Sigma_{-i})$, we let $p_{\mu_i} \in \Delta(S_{-i})$ denote the **predictive probabilities** implied by μ_i : for each $s_{-i} \in S_{-i}$,

$$p_{\mu_i}(s_{-i}) = \int_{\Sigma_{-i}} \sigma_{-i}(s_{-i}) \mu_i(d\sigma_{-i}).$$

We summarize our notation in the following table and illustrate it with an example.

¹³Statistical independence follows from random matching: For each $i \in I$, let P_i denote the of agents playing in role i , and let $\varsigma_i : P_i \rightarrow S_i$ denote the (measurable) strategy map of population i . If agents are drawn at random from their populations, that is, according to a *uniform* distribution on $\times_{i \in I} P_i$, then the induced distribution on S given $(\varsigma_i)_{i \in I}$ is a product measure.

¹⁴We use symbol \times to denote both the Cartesian product of sets and the product of measures.

We will refer to this example repeatedly.

Notation	Terminology
Γ	extensive-form game
$i, j \in I$	players ($j = 0 \notin I$ denotes chance)
$h \in H_i$	information sets of i
\preceq ($<$)	(strict) precedence relation of Γ
(H_i, \preceq)	directed forest of information sets of i
$a_i \in A_i(h)$	i 's actions at $h \in H_i$

$s_i \in S_i = \times_{h \in H_i} A_i(h)$	strategies of i
$s \in S$ ($s_{-i} \in S_{-i}$, $s_{0,-i} \in S_{0,-i}$)	strategy profiles (of $-i = I \setminus \{i\}$, of $\{0\} \cup I \setminus \{i\}$)
$S_{0,I}(h)$	strategy profiles (including 0) reaching h
$S_i(h) = \text{proj}_{S_i} S_{0,I}(h)$	strategies of i allowing for h
$S_{0,-i}(h) = \text{proj}_{S_{0,-i}} S_{0,I}(h)$	strategy profiles of 0 and $-i$ allowing for h
$S_{-i}(h) = \text{proj}_{S_{-i}} S_{0,I}(h)$	strategy profiles of $-i$ allowing for h
$\sigma_j \in \Delta(S_j)$	strategy distributions on S_j
$\sigma_{-i} \in \Sigma_{-i} \subset \Delta(S_{-i})$	product distributions on S_{-i}
$\mu_i \in \Delta(\Sigma_{-i})$	beliefs of i
$p_{\mu_i} \in \Delta(S_{-i})$	predictive probabilities implied by μ_i
$z \in Z$	terminal histories/nodes
$\zeta : S_0 \times S \rightarrow Z$	outcome function
$\gamma : Z \rightarrow C$	consequence function
$v_i : C \rightarrow \mathbb{R}$	vNM utility function
$v_i \circ \gamma : Z \rightarrow \mathbb{R}$	payoff function

Example 1 *The game depicted in Figure 1 is a two-person, common-interest, multistage game where $I = \{1, 2\}$ and 0 is chance. If we identify nodes with histories,¹⁵ information sets, actions sets and terminal histories/nodes are as follows:*

$$\begin{aligned}
H_0 &= \{\text{In}\}, H_1 = \{\{\emptyset\}, \{(\text{In}, \text{G})\}\}, H_2 = \{(\text{In}, \text{G})\}, \\
A_0(\{\text{In}\}) &= \{\text{E}, \text{G}\}, A_2(\{(\text{In}, \text{G})\}) = \{\text{L}, \text{R}\}, \\
A_1(\{\emptyset\}) &= \{\text{In}, \text{Out}\}, A_1(\{(\text{In}, \text{G})\}) = \{\text{T}, \text{M}, \text{B}\}, \\
Z &= \{\text{Out}, (\text{In}, \text{E})\} \cup \{(\text{In}, \text{G})\} \times \{\text{T}, \text{M}, \text{B}\} \times \{\text{L}, \text{R}\}.
\end{aligned}$$

Numbers at terminal histories/nodes, including the boxes in the matrix subgame, give the common payoff of players 1 and 2. The probabilities of chance moves are $\sigma_0(\text{E}) = \sigma_0(\text{G}) = \frac{1}{2}$.

¹⁵See, e.g., Chapter 11 in Osborne and Rubinstein (1994).

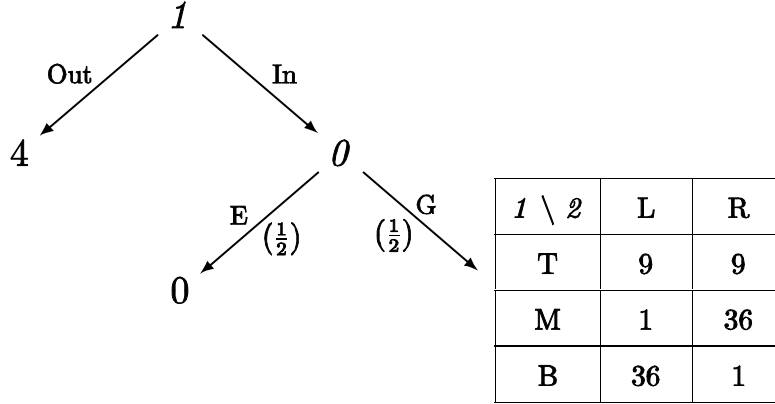


Figure 1: Leading Example

Assume that the agent in role 1 has the following belief:

$$\mu_1(\sigma_2) = \begin{cases} \frac{1}{2} & \text{if } \sigma_2 \in \{\delta_L, \delta_R\}, \\ 0 & \text{otherwise,} \end{cases}$$

where δ_x denotes the Dirac measure supported by x . Intuitively, he thinks that all the agents playing in role 2 “attended the same school” and hence are doing the same thing, but he does not know what. The induced predictive probabilities are $p_{\mu_1}(\text{L}) = p_{\mu_1}(\text{R}) = \frac{1}{2}$. \blacktriangle

3 Smooth-ambiguity preferences over actions

In this section, we take the perspective of an agent, or decision maker, playing in role i , henceforth DM_i , with given beliefs about the behavior of agents in different roles and a well defined plan. Specifically, he has a plan s_i specifying the action $s_{i,h} \in A_i(h)$ he expects to take (but he is not committed to take) at each information set $h \in H_i$; he has no randomization technology beyond what is already explicitly represented in the extensive form of the game (see Section 2), and we assume that he is certain about his contingent behavior, i.e., he has a deterministic contingent plan.

Conditional distributions and conditional objective expected utility Let $\Sigma_{-i}(h)$ denote the set of distributions that assign positive probability to $S_{-i}(h)$, that is,

$$\Sigma_{-i}(h) = \{\sigma_{-i} \in \Sigma_{-i} : \sigma_{-i}(S_{-i}(h)) > 0\}.$$

For every $h \in H_i$ and $\sigma_{-i} \in \Sigma_{-i}(h)$, we can compute the objective conditional distribution on the opponents' strategy profiles consistent with h :

$$\forall (s_0, s_{-i}) \in S_{0,-i}(h), \sigma_{0,-i}(s_0, s_{-i}|h) = \frac{\sigma_0(s_0)\sigma_{-i}(s_{-i})}{(\sigma_0 \times \sigma_{-i})(S_{0,-i}(h))} \quad (1)$$

(note that that $\sigma_{-i} \in \Sigma_{-i}(h)$ if and only if $(\sigma_0 \times \sigma_{-i})(S_{0,-i}(h)) > 0$, because σ_0 is strictly positive). With this, we can define the **strategic-form vNM conditional expected utility function**

$$U_i(\cdot, \cdot|h) : \begin{array}{ll} S_i(h) \times \Sigma_{-i}(h) & \rightarrow \mathbb{R}, \\ (s_i, \sigma_{-i}|h) & \mapsto \sum_{s_{0,-i} \in S_{0,-i}(h)} \sigma_{0,-i}(s_{0,-i}|h) v_i(\gamma(\zeta(s_i, s_{0,-i}))). \end{array}$$

In words, if i is certain that σ_{-i} is the true objective probability model, then upon observing h (he believes that) his conditional objective expected utility from following strategy $s_i \in S_i(h)$ is $U_i(s_i, \sigma_{-i}|h)$.¹⁶

Plans and replacements Plan s_i yields a continuation on the information sets in H_i following any given $\bar{h} \in H_i$ (that is, the projection of s_i onto $\times_{\{h \in H_i: \bar{h} \preceq h\}} A_i(h)$). DM_i expects to continue according to this plan, but he knows (by perfect recall) that he has already chosen the actions leading to \bar{h} , possibly violating s_i , and he considers the consequences of choosing action $a_i \in A_i(\bar{h})$, again possibly violating s_i . With this, it is convenient to define the **replacement** plan $(s_i|\bar{h}, a_i)$ obtained by replacing s_i with the already chosen actions at information sets preceding \bar{h} and with action a_i at \bar{h} :

$$(s_i|\bar{h}, a_i)_h = \begin{cases} a_i & \text{if } h = \bar{h}, \\ \alpha_i(h, \bar{h}) & \text{if } h \prec \bar{h}, \\ s_{i,h} & \text{otherwise,} \end{cases}$$

where $\alpha_i(h, \bar{h})$ is the action chosen at $h \prec \bar{h}$ in order to reach \bar{h} .¹⁷ To ease notation, we let $s_i|h$ denote the replacement plan obtained when $s_{i,h}$ is played at \bar{h} , that is, $(s_i|\bar{h}) = (s_i|\bar{h}, s_{i,h})$.

Action values We assume that DM_i preferences over actions, given his beliefs and plan, satisfy the smooth-ambiguity model of KMM: On top of the **vNM utility function** v_i :

¹⁶One can show that this coincides with the more familiar formula

$$U_i(s_i, \sigma_{-i}|h) = \sum_{x \in h} \mathbb{P}_{s_i, \sigma_{-i}}(x|h) \sum_{z \in Z} \mathbb{P}_{s_i, \sigma_{-i}}(z|x) v_i(\gamma(z))$$

where $\mathbb{P}_{s_i, \sigma_{-i}}(\cdot|\cdot)$ denotes the probability of reaching a node conditional on an information set, or an earlier node, given by σ_{-i} and the known probabilities of chance moves.

¹⁷By perfect recall, α_i is well defined.

$C \rightarrow \mathbb{R}$ specified by game Γ , we assume that there is a continuous and *strictly increasing second-order utility* function

$$\phi_i : \mathbb{V}_i \rightarrow \mathbb{R},$$

where

$$\mathbb{V}_i = \left[\min_z v_i(\gamma(z)), \max_z v_i(\gamma(z)) \right]$$

is the convex hull of the range of v_i . For every given $h \in H_i$, $\mu_i \in \Delta(\Sigma_{-i}(h))$, and $s_i \in S_i$, DM_i assigns values to actions $a_i \in A_i(h)$ as follows:

$$V_i(a_i|h; s_i, \mu_i, \phi_i) = \phi_i^{-1} \left(\int_{\Sigma_{-i}(h)} \phi_i(U_i((s_i|h, a_i), \sigma_{-i}|h)) \mu_i(d\sigma_{-i}) \right), \quad (2)$$

where If μ_i assigns probability 1 to some $\sigma_{-i} \in \Sigma_{-i}(h)$ (that is, $\mu_i = \delta_{\sigma_{-i}} \in \Delta(\Sigma_{-i}(h))$), then

$$V_i(a_i|h; s_i, \mu_i, \phi_i) = \phi_i^{-1}(\phi_i(U_i((s_i|h, a_i), \sigma_{-i}|h))) = U_i((s_i|h, a_i), \sigma_{-i}|h).$$

Thus, ambiguity attitudes are immaterial when DM_i is certain about the true probability model, because in this case he does not perceive any ambiguity. Note also that (2) boils down to the classical subjective expected utility formula if ϕ_i is linear (ambiguity neutrality), hence equivalent to the identity function $\text{Id}_{\mathbb{V}_i}$. On the other hand, ambiguity aversion is characterized by the concavity of ϕ_i . If $0 < \mu_i(\Sigma_{-i}(h)) < 1$, we replace μ_i in the right hand side of (2) with the conditional belief $\mu_i(\cdot|h)$. We explain the details in Section 4. We emphasize in our notation only the dependence of values of i 's actions on parameter ϕ_i , not on the vNM utility function v_i , because we are going to consider different possible shapes of ϕ_i (in particular, linear and concave) with a fixed v_i .

Example 2 *In the game of Figure 1,*

$$\begin{aligned} U_1(\text{In.B}, \delta_L | \{(\text{In}, \text{G})\}) &= 36, \\ U_1(\text{In.B}, \delta_R | \{(\text{In}, \text{G})\}) &= 1, \end{aligned}$$

and

$$\begin{aligned} U_1(\text{In.B}, \delta_L | \{\emptyset\}) &= \frac{1}{2} \cdot 36 = 18, \\ U_1(\text{In.B}, \delta_R | \{\emptyset\}) &= \frac{1}{2} \cdot 1 = \frac{1}{2}. \end{aligned}$$

Assume $\phi_1(u_i) = \sqrt{u_i}$ and let μ_1 be the belief of Example 1,

$$\mu_1(\sigma_2) = \begin{cases} \frac{1}{2} & \text{if } \sigma_2 \in \{\delta_L, \delta_R\}, \\ 0 & \text{otherwise,} \end{cases}$$

then

$$V_1(\text{B} | \{(\text{In}, \text{G})\}; s_1, \mu_1, \phi_1) = \left(\frac{1}{2} \sqrt{36} + \frac{1}{2} \sqrt{1} \right)^2 = (3.5)^2 = 12.25$$

for each $s_1 \in S_1((\text{In}, \text{G}))$. ▲

4 Conditional beliefs and unimprovability

A prior belief $\mu_i \in \Delta(\Sigma_{-i})$ over co-players' strategy distributions induces a joint belief $\pi_{\mu_i} \in \Delta(S_{0,-i} \times \Sigma_{-i})$ determined by the following equation:

$$\forall (s_0, s_{-i}, E_{-i}) \in S_{0,-i} \times \mathcal{B}(\Sigma_{-i}), \pi_{\mu_i}(\{(s_0, s_{-i})\} \times E_{-i}) = \sigma_0(s_0) \int_{E_{-i}} \sigma_{-i}(s_{-i}) \mu_i(d\sigma_{-i}). \quad (3)$$

Note that¹⁸

$$\pi_{\mu_i}(S_{0,-i} \times E_{-i}) = \mu_i(E_{-i}), \quad (4)$$

and

$$\pi_{\mu_i}(S_0 \times \{s_{-i}\} \times \Sigma_{-i}) = p_{\mu_i}(s_{-i}).$$

Each information set $h \in H_i$ corresponds to the conditioning event $S_{0,-i}(h) \times \Sigma_{-i}(h)$. Let

$$H_i(\mu_i) = \{h \in H_i : \pi_{\mu_i}(S_{0,-i}(h) \times \Sigma_{-i}(h)) > 0\} = \{h \in H_i : p_{\mu_i}(S_{-i}(h)) > 0\}$$

denote the subset of information sets of player i that he believes he can reach with positive probability.¹⁹ If $h \in H_i(\mu_i)$, then we can derive the conditional probability of every measurable set $E_{-i} \in \mathcal{B}(\Sigma_{-i})$ of strategy distributions: **[Shorten this?][E: in the version we submit yes, now I would keep it]**

$$\begin{aligned} \mu_i(E_{-i}|h) &= \pi_{\mu_i}(S_{0,-i} \times E_{-i} | S_{0,-i}(h) \times \Sigma_{-i}(h)) \\ &= \frac{\pi_{\mu_i}((S_{0,-i} \times E_{-i}) \cap (S_{0,-i}(h) \times \Sigma_{-i}(h)))}{\pi_{\mu_i}(S_{0,-i}(h) \times \Sigma_{-i}(h))} \\ &= \frac{\pi_{\mu_i}(S_{0,-i}(h) \times (\Sigma_{-i}(h) \cap E_{-i}))}{\sum_{(s_0, s_{-i}) \in S_{0,-i}(h)} \sigma_0(s_0) \int_{\Sigma_{-i}(h)} \sigma_{-i}(s_{-i}) \mu_i(d\sigma_{-i})} \\ &= \frac{\sum_{(s_0, s_{-i}) \in S_{0,-i}(h)} \sigma_0(s_0) \int_{\Sigma_{-i}(h) \cap E_{-i}} \sigma_{-i}(s_{-i}) \mu_i(d\sigma_{-i})}{\sum_{(s_0, s_{-i}) \in S_{0,-i}(h)} \sigma_0(s_0) p_{\mu_i}(s_{-i})} \\ &= \frac{\int_{E_{-i}} (\sigma_0 \times \sigma_{-i})(S_{0,-i}(h)) \mu_i(d\sigma_{-i})}{(\sigma_0 \times p_{\mu_i})(S_{0,-i}(h))}. \end{aligned} \quad (5)$$

¹⁸One may use beliefs over the product space $S_{0,-i} \times \Sigma_{-i}$ as the primitive object. Of course, the structural assumption that agents know the probabilities of chance moves and that they are randomly matched with opponents drawn large populations implies that only beliefs that admit the representation (3) for some $\mu_i \in \Delta(\Sigma_{-i})$ are admissible. Then we can derive the belief over distributions using (4). For a decision theoretic approach akin to the one discussed in this footnote, with models in place of population distributions, see Cerreia-Vioglio, Maccheroni, Marinacci, and Montrucchio (2013).

¹⁹The equality holds because the probabilities of chance moves are strictly positive.

For example, if μ_i has *finite support*:

$$\mu_i(\sigma_{-i}|h) = \frac{\mu_i(\sigma_{-i})(\sigma_0 \times \sigma_{-i})(S_{0,-i}(h))}{\sum_{\sigma'_{-i} \in \text{Supp}\mu_i} \mu_i(\sigma'_{-i})(\sigma_0 \times \sigma'_{-i})(S_{0,-i}(h))}.$$

With this, we consider the profile of conditional beliefs $(\mu_i(\cdot|h))_{h \in H_i}$ derived from prior belief μ_i , we let

$$V_i(a_i|h; s_i, \mu_i, \phi_i) = \phi_i^{-1} \left(\int_{\Sigma_{-i}(h)} \phi_i(U_i((s_i|h, a_i), \sigma_{-i}|h)) \mu_i(d\sigma_{-i}|h) \right) \quad (6)$$

whenever $h \in H_i(\mu_i)$, and we derive well defined preferences over actions (given μ_i and s_i) only for information sets that are possible according to μ_i . This is what we need for our definition of unimprovability, which is instead silent about choices at information sets deemed unreachable.

Definition 1 *A strategy s_i is (μ_i, ϕ_i) -unimprovable if*

$$\forall h \in H_i(\mu_i), s_{i,h} \in \arg \max_{a_i \in A_i(h)} V_i(a_i|h; s_i, \mu_i, \phi_i).$$

Since the game has finite horizon, we can interpret unimprovability as “folding-back optimality”: given (μ_i, ϕ_i) , DM_i derives a contingent plan that prescribes a choice for each information set he deems reachable. He starts from information sets $h \in H_i(\mu_i)$ with no followers in $H_i(\mu_i)$, and to each one of them he assigns an action $s_{i,h}$ that maximizes $V_i(a_i|h; s_i, \mu_i, \phi_i)$. Then he folds back, considering the information sets $h \in H_i(\mu_i)$ such that every follower has no further followers; for every such follower, viz. h' , DM_i predicts that the previously selected maximizing action $s_{i,h'}$ will be chosen; and so on until all the reachable information sets in $H_i(\mu_i)$ have been covered backwards.

Of course, we could define beliefs, and thus impose optimality requirements, also at information sets in $H_i \setminus H(\mu_i)$. For the time being, we are not interested in doing so: the moves of DM_i at each $h \in H_i \setminus H(\mu_i)$ will be immaterial for the equilibrium outcomes (by ex post perfect recall and confirmed beliefs, see Section 5), and impossible to predict for the opponents as long as we do not assume that they know the payoff function of i (cf. Section 7).

From the point of view of an external observer, or of agents in roles different from i , it is impossible to distinguish between two strategies of DM_i that yield the same outcomes independently of the opponents' behavior. This leads to the following notion of equivalence, which will play an important role in comparing self-confirming equilibria for different levels of ambiguity aversion (see Section 6).

Definition 2 (Kuhn, 1953) Two (possibly degenerate) strategy distributions σ_i^* and σ_i are **realization-equivalent** if they induce the same distribution on terminal nodes, that is,

$$\forall (z, s_0, s_{-i}) \in Z \times S_0 \times S_{-i}, \quad \sum_{s_i: \zeta(s_0, s_i, s_{-i})=z} \sigma_i^*(s_i) = \sum_{s_i: \zeta(s_0, s_i, s_{-i})=z} \sigma_i(s_i).$$

The set of strategy distributions realization-equivalent to σ_i^* is denoted by $[\sigma_i^*]$.

Let $H_i(s_i) = \{h \in H_i : s_i \in S_i(h)\}$ denote the subset of information sets of DM_i that are compatible with strategy s_i . Focusing on pure strategies, we obtain the following observation:

Remark 1 (Theorem 1, Kuhn 1953) Fix strategies $s_i, s_i^* \in S_i$; $s_i \in [s_i^*]$ if and only if s_i and s_i^* are behaviorally equivalent, that is, if and only if $H_i(s_i) = H_i(s_i^*)$ and $s_{i,h} = s_{i,h}^*$ for each $h \in H_i(s_i^*)$.

Now, suppose that DM_i is ambiguity neutral: $\phi_i = \text{Id}_{\mathbb{V}_i}$. Then, by a classical dynamic programming result, unimprovability is equivalent to “global” (ex ante) subjective EU-maximization.²⁰

Proposition 1 For every strategy $s_i^* \in S_i$ and prior $\mu_i \in \Delta(\Sigma_{-i})$ the following are equivalent:

- (1) $[s_i^*]$ contains a $(\mu_i, \text{Id}_{\mathbb{V}_i})$ -unimprovable strategy,
- (2) $s_i^* \in \arg \max_{s_i \in S_i} \sum_{(s_0, s_{-i})} \sigma_0(s_0) p_{\mu_i}(s_{-i}) v_i(\gamma(\zeta(s_0, s_i, s_{-i})))$.

We introduce the following strengthening of unimprovability:

Definition 3 A strategy s_i is (μ_i, ϕ_i) -**sequentially optimal** if

$$\forall h \in H_i(\mu_i), s_i \in \arg \max_{s'_i \in S_i} \phi_i^{-1} \left(\int_{\Sigma_{-i}(h)} \phi_i(U_i((s'_i|h), \sigma_{-i}|h)) \mu_i(d\sigma_{-i}|h) \right).$$

If DM_i has dynamically inconsistent preferences over strategies, a (μ_i, ϕ_i) -sequentially optimal strategy may not exist, as illustrated in Example 3 below. However, if DM_i is ambiguity neutral (hence, his preferences are dynamically consistent), unimprovability coincides with sequential optimality:

Proposition 2 A strategy s_i is $(\mu_i, \text{Id}_{\mathbb{V}_i})$ -unimprovable if and only if it is $(\mu_i, \text{Id}_{\mathbb{V}_i})$ -sequentially optimal.

²⁰ All the dynamic programming results of this section can be proved by standard folding-back arguments.

Example 3 Consider the game of Figure 1 and the belief μ_1 of Example 1:

$$\mu_1(\sigma_2) = \begin{cases} \frac{1}{2} & \text{if } \sigma_2 \in \{\delta_L, \delta_R\}, \\ 0 & \text{otherwise,} \end{cases}$$

Then

$$H_1(\mu_1) = H_1 = \{\{\emptyset\}, \{\text{In}, \text{G}\}\}.$$

The induced belief on $S_{0,2} \times \Sigma_2$ is

$$\pi_{\mu_1}(s_0, s_2, \sigma_2) = \begin{cases} \frac{1}{4} & \text{if } (s_0, s_2, \sigma_2) \in \{\text{E}, \text{G}\} \times \{(\text{L}, \delta_L), (\text{R}, \delta_R)\}, \\ 0 & \text{otherwise.} \end{cases}$$

For an ambiguity neutral player 1 with belief μ_1 , the value of M and B at (In, G) is $(\frac{1}{2}36 + \frac{1}{2}1) > 9$, so it is higher than the value of T. Therefore, by folding-back, the (μ_1, Id_{V_i}) -unimprovable strategies are In.M and In.B.

Now suppose instead that the ambiguity attitudes of player 1 are represented by some strictly concave $\bar{\phi}_1$ with:

$$\begin{aligned} \bar{\phi}_1(u) &= \sqrt{u} \text{ if } 1 \leq u < 36, \\ \bar{\phi}_1\left(\frac{1}{2}\right) &= -1. \end{aligned}$$

At (In, G), player 1 still prefers M (or B) over T, because:

$$\begin{aligned} V_1(\text{M} | \{\text{In}, \text{G}\}; s_1, \mu_1; \bar{\phi}_1) &= \bar{\phi}_1^{-1}\left(\frac{1}{2}\bar{\phi}_1(1) + \frac{1}{2}\bar{\phi}_1(36)\right) \\ &= (3.5)^2 > 9 = V_1(\text{T} | \{\text{In}, \text{G}\}; s_1, \mu_1; \bar{\phi}_1). \end{aligned}$$

Hence, a $(\mu_1, \bar{\phi}_1)$ -unimprovable strategy must prescribe action M or B at (In, G). But then, it must also prescribe action Out at $\{\emptyset\}$. Indeed, for every strategy s_1 such that $s_{1,(\text{In}, \text{G})} \in \{\text{M}, \text{B}\}$ we have:

$$\begin{aligned} V_1(\text{In} | \{\emptyset\}; s_1, \mu_1, \bar{\phi}_1) &= \bar{\phi}_1^{-1}\left(\frac{1}{2} \cdot \bar{\phi}_1\left(\frac{1}{2} \cdot 36\right) + \frac{1}{2} \cdot \bar{\phi}_1\left(\frac{1}{2} \cdot 1\right)\right) \\ &= \left(\frac{1}{2} \cdot \sqrt{18} - \frac{1}{2}\right)^2 < 4 = V_1(\text{Out} | \{\emptyset\}; s_1, \mu_1, \bar{\phi}_1). \end{aligned}$$

So, on the one hand, the only $(\mu_1, \bar{\phi}_1)$ -unimprovable strategies are Out.M and Out.B. On the other hand, from the perspective of the agent at the root of Γ , the value of committing to strategy In.T is

$$V_1(\text{In} | \{\emptyset\}; \text{In.T}, \mu_1, \bar{\phi}_1) = \left(\sqrt{\frac{1}{2} \cdot 9}\right)^2 > 4 = V_1(\text{Out} | \{\emptyset\}; \text{Out}.a_1, \mu_1, \bar{\phi}_1), \quad (7)$$

for all $a_1 \in \{\text{T}, \text{M}, \text{B}\}$. So, player 1 would commit to In.T if he only could. \blacktriangle

This example illustrates the well-known dynamic inconsistency of preferences of decision makers with non-neutral attitudes towards ambiguity.^{21,22} To address this problem, we assume that agents are sufficiently sophisticated to understand the incentives they would face at each information set deemed possible, and plan/predict their contingent behavior by folding back. The resulting (μ_i, ϕ_i) -unimprovable strategy (or strategies) represents how agents in role i with belief μ_i predict they would choose at future information sets; such strategy and μ_i yield a value for each action available at the current information set.

Of course, DM_i may be indifferent at some information sets. A further refinement can be obtained for ambiguity averse agents by imposing a consistent-planning condition: whenever DM_i is indifferent at h then he breaks ties according to the preferences at the immediate predecessor of h in H_i . If this does not solve all the indifferences, ties are broken according to the preferences of the twice-removed predecessor, and so on. We omit this refinement for simplicity, and also because we find it arbitrary.

5 Selfconfirming equilibrium

BCMM analyze a notion of smooth self-confirming equilibrium under the assumption that agents play the strategic form of a game with feedback and ambiguity attitudes, as with the strategy method in lab experiments. Specifically, consider a triple (Γ, f, ϕ) , where Γ is a standard extensive-form game, $f = (f_i : Z \rightarrow M)_{i \in I}$ is a profile of feedback functions such that every f_i describes the message $m \in M$ that player i observes ex post as a function of the terminal node, and $\phi = (\phi_i : \mathbb{V}_i \rightarrow \mathbb{R})_{i \in I}$ is a profile of strictly increasing functions capturing players attitudes toward ambiguity. The **structural strategic feedback function** (as if chance were a player) is then $F_i = f_i \circ \zeta : S_0 \times S \rightarrow M$. We let $\hat{F}_i(s_i, \sigma_{-i}) \in \Delta(M)$ denote the **pushforward** distribution of messages induced by strategy s_i and the profile of strategy distributions σ_{-i} , given σ_0 . Specifically:

$$\forall (s_i, \sigma_{-i}, m) \in S_i \times \Sigma_{-i} \times M, \hat{F}_i(s_i, \sigma_{-i})(m) = \sum_{(s_0, s_{-i}) : F_i(s_0, s_i, s_{-i}) = m} \sigma_0(s_0) \sigma_{-i}(s_{-i}).$$

To relate to BCMM it is convenient to define the strategic, or normal form of a game with feedback (Γ, f) . The **normal-form** (expected) **payoff function** of player i is $U_i : S \rightarrow \mathbb{R}$ with

$$\forall s \in S, U_i(s) = \sum_{s_0 \in S_0} \sigma_0(s_0) v_i(\gamma(\zeta(s_0, s)))$$

²¹ See Siniscalchi (2011), for illustrative examples and an in-depth analysis of this issue.

²² Note that these dynamic inconsistencies arise as a consequence of the combination of Bayesian updating and non-neutral ambiguity attitudes. Indeed, it is not even obvious from the decision theoretic literature that ambiguity averse players are supposed to update beliefs according to the standard rules of conditional probabilities (see Epstein and Schneider 2007, Hanany and Klibanoff 2009, and Hanany et al. 2017). Instead we take the position that these rules are part of rational cognition, and we stick to them. This position is supported also by works that justify Bayesian updating in an evolutionary perspective, see Blume and Easley (2006) and the references therein.

Similarly, we define **normal-form feedback function** $\bar{F}_i : S \rightarrow \bar{M}$ as follows: If strategy profile s is played in the long run, then i observes the *distribution* of messages determined by s and chance probabilities. Therefore, we let $\bar{M} = \Delta(M)$ and

$$\forall (s, m) \in S \times M, \bar{F}_i(s)(m) = \hat{F}_i(s_i, s_{-i})(m) = \sum_{s_0: F_i(s_0, s) = m} \sigma_0(s_0).$$

With this, the normal form of (Γ, f) is $\mathcal{N}(\Gamma, f) = (S_i, U_i, \bar{F}_i)_{i \in I}$. The equilibrium concept of BCMM applies to $(\mathcal{N}(\Gamma, f), \phi) = (S_i, U_i, \bar{F}_i, \phi_i)_{i \in I}$ under the assumption that each agent in role i covertly commits in advance to a strategy s_i . Here, instead, we analyze an equilibrium concept that is appropriate when agents play (Γ, f, ϕ) with the “direct method” making choices as the play unfolds, and we compare it with the strategic-form concept of BCMM.

Given that the information structure of Γ is assumed to satisfy perfect recall, we maintain the assumption that (Γ, f) satisfies “**ex post perfect recall**”:²³

Assumption (*Ex post perfect recall*) For every player $i \in I$, the augmented collection of information sets that includes the partition of Z induced by f_i , $H_i \cup \{f_i^{-1}(m) : m \in f_i(Z)\}$, satisfies the perfect recall assumption. In particular, for all terminal histories $z, z' \in Z$, if there are an information set $h \in H_i$ and a node $x \in h$ such that $x \prec z$ and either z' has no predecessor in h , or $\alpha_i(x, z) \neq \alpha_i(x', z')$ for the predecessor x' of z' in h , then $f_i(z) \neq f_i(z')$.

Furthermore, we also consider (but we do not always assume) the following property of feedback:²⁴

Definition 4 An extensive-form game with feedback (Γ, f) satisfies **observable payoffs** whenever the payoff of every player only depends on his ex post information:

$$\forall (i, z, z') \in I \times Z^2, f_i(z) = f_i(z') \Rightarrow v_i(\gamma(z)) = v_i(\gamma(z')).$$

In other words, the payoff function is constant on each element of the ex post information partition $\{f_i^{-1}(m) : m \in f_i(Z)\}$. We say that (Γ, f, ϕ) satisfies observable payoffs if (Γ, f) does.

In some examples, we will assume that agents observe the terminal node they reach, that is $f_i = \text{Id}_Z$. We call this hypothesis **perfect feedback**.

²³See Battigalli et al. (2016b). There the statement of the ex post perfect recall property is slightly different. The two versions are equivalent for extensive-form representations that specify the information of each player i at each nonterminal node, not only those where i is active (cf. Battigalli and Bonanno, 1999). Otherwise, one should use the statement of this paper.

²⁴See Battigalli et al. (2016b) for an in-depth analysis of the properties of feedback and how they affect the SCE set.

Definition 5 A *self-confirming equilibrium (SCE)* of (Γ, f, ϕ) is a profile of strategy distributions $\bar{\sigma} = (\bar{\sigma}_i)_{i \in I}$ with the following property: For each $i \in I$ and $\bar{s}_i \in \text{Supp} \bar{\sigma}_i$ there is a belief $\mu_{\bar{s}_i} \in \Delta(\Sigma_{-i})$ such that

(rationality) \bar{s}_i is $(\mu_{\bar{s}_i}, \phi_i)$ -unimprovable,

(confirmed beliefs) $\mu_{\bar{s}_i} \left(\left\{ \sigma_{-i} \in \Sigma_{-i} : \hat{F}_i(\bar{s}_i, \sigma_{-i}) = \hat{F}_i(\bar{s}_i, \bar{\sigma}_{-i}) \right\} \right) = 1$.

An SCE $\bar{\sigma}$ is a **symmetric SCE (symSCE)** if, for each $i \in I$, there is a pure strategy \bar{s}_i with $\bar{\sigma}_i(\bar{s}_i) = 1$, that is, if all agents in the same population i play the same pure strategy.

The confirmed beliefs condition requires that the belief $\mu_{\bar{s}_i}$ justifying \bar{s}_i exclude all the distributions that are not observationally equivalent to the true one, $\bar{\sigma}_{-i}$. When profiles $\bar{\sigma}$ and $(\mu_{\bar{s}_i})_{i \in I, \bar{s}_i \in \text{Supp} \bar{\sigma}_i}$ satisfy the foregoing SCE conditions, we say that $\bar{\sigma}$ is **justified by confirmed beliefs** $(\mu_{\bar{s}_i})_{i \in I, \bar{s}_i \in \text{Supp} \bar{\sigma}_i}$. The set of (symmetric) self-confirming equilibria of (Γ, f, ϕ) is denoted by $SCE(\Gamma, f, \phi)$ ($\text{symSCE}(\Gamma, f, \phi)$).

As in BCMM, the confirmed beliefs condition says that an agent rules out opponents' strategy distributions inconsistent with his "empirical distribution" of observations. More specifically, we consider stability conditions for a profile of strategy distributions in a scenario where agents drawn at random from large populations (corresponding to game roles) play the given game recurrently and learn from their personal experience. Suppose each agent keeps playing the same (pure) strategy for a very long time, and consider an agent in role i who has been playing \bar{s}_i and accumulated a large dataset of personal observations. With probability 1, and in the limit, this dataset is summarized by the frequency distribution of observations generated by his strategy \bar{s}_i and by the actual strategy distributions for the opponents' populations, $\bar{\sigma}_{-i}$, that is, $\hat{F}_i(\bar{s}_i, \bar{\sigma}_{-i})$.²⁵ Every profile of distributions σ_{-i} that yields the same distribution of observations is empirically indistinguishable from the true one, $\bar{\sigma}_{-i}$, and hence it cannot be objectively ruled out.

We first observe that every game with feedback and ambiguity attitudes has an SCE:

Proposition 3 For every (Γ, f, ϕ) , $SCE(\Gamma, f, \phi) \neq \emptyset$.

Intuitively, every finite game Γ has a sequential equilibrium $(\bar{\beta}_i)_{i \in I}$ in behavioral strategies. Consider the corresponding mixed strategy profile $(\bar{\sigma}_i)_{i \in I}$ and let $\mu_{\bar{s}_i} = \delta_{\bar{\sigma}_{-i}}$ for every i and $\bar{s}_i \in \text{Supp} \bar{\sigma}_i$. Since no ambiguity is perceived, agents with these beliefs behave as expected utility maximizers. With this, it can be shown that each $\bar{s}_i \in \text{Supp} \bar{\sigma}_i$ is $(\mu_{\bar{s}_i}, \phi_i)$ -unimprovable, because $\bar{\sigma}$ corresponds to a sequential equilibrium.

Observe that our definition does not coincide with the one proposed in BCMM, because the rationality assumption of BCMM is given by the ex ante KMM criterion:

$$\bar{s}_i \in \arg \max_{s_i \in \bar{S}_i} \int_{\Sigma_{-i}} \phi_i(U_i(s_i, \sigma_{-i})) \mu_i(d\sigma_{-i}).$$

²⁵See Battigalli et al. (2016c) for a learning foundation of selfconfirming equilibrium with non-neutral ambiguity attitudes.

This best reply condition is appropriate only in simultaneous moves games, possibly obtained by having agents play the strategic form of a sequential game (cf. BCMM, pp 665-667). Here, instead, we require agents to maximize the KMM value over actions at every information set they deem reachable. Therefore, the set of self-confirming equilibria à la BCMM of a sequential game (Γ, f, ϕ) is $SCE(\mathcal{N}(\Gamma, f), \phi) = SCE\left((S_i, U_i, \bar{F}_i, \phi_i)_{i \in I}\right)$.

Proposition 1 implies that our definition is (realization) equivalent to the one of BCMM when agents are *ambiguity neutral*: Given $\bar{\sigma} \in SCE(\mathcal{N}(\Gamma, f), \text{Id}_{V_i})$ with associated beliefs $(\mu_{s_i})_{i \in I, s_i \in \text{Supp}\bar{\sigma}_i}$, replace each $s_i \in \text{Supp}\bar{\sigma}_i$ with a realization-equivalent $(\mu_{s_i}, \text{Id}_{V_i})$ -unimprovable strategy $\hat{s}_i(s_i)$ (see Proposition 1), then define σ_i as the pushforward of $\bar{\sigma}_i$ under map $\hat{s}_i(\cdot)$, that is, $\sigma_i = \bar{\sigma}_i \circ \hat{s}_i^{-1}$; the resulting profile σ with associated beliefs $(\mu_{\hat{s}_i(s_i)})_{i \in I, s_i \in \text{Supp}\bar{\sigma}_i}$ satisfies the SCE conditions. To sum up:

Remark 2 *Suppose that ϕ_i is linear for each $i \in I$. Then, for every $\bar{\sigma} \in SCE(\mathcal{N}(\Gamma, f), \phi)$ there is some $\sigma \in SCE(\Gamma, f, \phi)$ such that, for each $i \in I$, $\bar{\sigma}_i$ and σ_i are realization-equivalent.*

If agents instead are ambiguity averse, the SCE's of a game are not realization-equivalent to SCE's of its strategic-form representation. Indeed, the two sets of SCE outcomes may be non-nested.

Example 4 *Consider the game of Figure 1. In Example ?? we considered a belief μ_1 and ambiguity attitudes $\bar{\phi}_1$ of player 1 such that Out.M and Out.B are $(\mu_1, \bar{\phi}_1)$ -unimprovable strategies. It follows that (Out.M, L) (for instance) is a symmetric SCE of the game, where Out.M is justified by belief μ_1 (trivially confirmed for any feedback function) and L is (vacuously) justified by any confirmed belief (for any $\bar{\phi}_2$). However, note that inequality (??) implies that (Out.M, L) does not belong to $SCE(\mathcal{N}(\Gamma, f), \bar{\phi})$. Specifically, (??) implies that, for every belief μ_1 and action $a_1 \in \{\text{T}, \text{M}, \text{B}\}$, strategy Out. a_1 is not ex-ante optimal, and thus it does not satisfy the best reply condition of BCMM. Hence, for every feedback f , Out is an SCE outcome of (Γ, f) , but it is not an SCE outcome of $\mathcal{N}(\Gamma, f)$. Now suppose that players just observe their monetary payoff (that is, $f = u$). Then it can be checked that, for every α_2 , (In.T, α_2) is an SCE of $\mathcal{N}(\Gamma, f)$ supported by the same belief μ_1 as before,²⁶ because player 1 prefers to commit to T and this prevents him from observing the long-run frequencies of L and R. Yet, In.T is not an SCE strategy of (Γ, f) because—as argued in Example ??—player 1 with prior μ_1 would deviate from T in the subgame despite the fact that μ_1 is the belief that minimizes the maximum value of deviating from T.²⁷ ▲*

Comment on knowledge of the game The definition of self-confirming equilibrium relies on very weak interpretive assumptions about agents' knowledge of the game: each agent playing in role i has to know only his preferences (v_i, ϕ_i) , the extensive-game form

²⁶The belief of player 2 is immaterial.

²⁷See Lemma ??, or Lemma 6 in BCMM.

(hence, also $\zeta : S \rightarrow Z$ and $\gamma : Z \rightarrow C$), and his feedback function $f_i : Z \rightarrow M$. Therefore, in a self-confirming equilibrium an agent in population i may believe that a positive fraction of agents in population j are implementing strategies that cannot be justified by any belief, given their true preferences and feedback (v_j, ϕ_j, f_j) (possibly unknown to agents in population i). Essentially, self-confirming equilibrium is a solution concept for incomplete information games with private values. In Section 8, we analyze a notion of rationalizable self-confirming equilibrium that is appropriate when there is common knowledge of (Γ, f, ϕ) . Here, we just note that our definition of SCE is realization equivalent to one that replaces unimprovability with respect to a prior belief μ_i with full unimprovability with respect to a system of conditional beliefs $(\mu_i(\cdot|h))_{h \in H_i}$ (see Proposition 9 in Section 7).

6 Comparative statics

In this section we analyze changes in the set of equilibria when feedback and ambiguity attitudes are modified with respect to some baseline \bar{f} and $\bar{\phi}$ respectively. Say that the feedback profile f is *coarser* than \bar{f} if f_i is \bar{f}_i -measurable²⁸ for each $i \in I$; in other words, for each player i , the partition of Z induced by f_i is coarser than the partition of Z induced by \bar{f}_i . It is quite straightforward to show that *if f is coarser than \bar{f} then $SCE(\Gamma, \bar{f}, \phi) \subseteq SCE(\Gamma, f, \phi)$* , because coarser ex post information makes it easier to satisfy the confirmed beliefs condition (see BCMM).

We would also like to prove an extension for sequential games of the following monotonicity theorem of BCMM: under observable payoffs, as ambiguity aversion increases, the set of SCE expands.

Say that (Γ, f, ϕ) features **more ambiguity aversion** than $(\Gamma, \bar{f}, \bar{\phi})$ if, for each $i \in I$, $\phi_i = \varphi_i \circ \bar{\phi}_i$ for some concave and strictly increasing function φ_i . Say that (Γ, f, ϕ) features **ambiguity aversion** if each ϕ_i is concave. BCMM proved that if (Γ, f) has observable payoffs and (Γ, f, ϕ) features more ambiguity aversion than $(\Gamma, \bar{f}, \bar{\phi})$, then

$$SCE(\mathcal{N}(\Gamma, f), \phi) \supseteq SCE(\mathcal{N}(\Gamma, \bar{f}), \bar{\phi}).$$

As already observed, self-confirming equilibria in the strategic and extensive form of the game are not realization equivalent. Hence, the monotonicity result of BCMM cannot be invoked to obtain an equivalent result for SCE in sequential games, not even in terms of induced outcome distributions. Yet, the core of the argument of BCMM can be adapted to sequential games when all the strategies in the support of an SCE with baseline ambiguity attitudes $\bar{\phi}$ are *sequentially optimal* under the confirmed beliefs that justify them, that is, at every reachable information set h the prescribed continuation strategy is the one that maximizes the value at h (see Definition 3).

²⁸That is, $f_i = \varphi_i \circ \bar{f}_i$ for some $\varphi_i: \bar{f}_i(Z) \rightarrow M$.

Lemma 1 *Suppose that (Γ, f, ϕ) features more ambiguity aversion than $(\Gamma, f, \bar{\phi})$ and fix an SCE $\bar{\sigma}$ of $(\Gamma, f, \bar{\phi})$ justified by the confirmed beliefs $(\mu_{s_i})_{i \in I, s_i \in \text{Supp} \bar{\sigma}_i}$. Suppose that for each $i \in I$, every $s_i \in \text{Supp} \bar{\sigma}_i$ is $(\mu_{s_i}, \bar{\phi}_i)$ -sequentially optimal. Then, there exists an SCE σ of (Γ, f, ϕ) that is realization equivalent to $\bar{\sigma}$ and is justified by the same confirmed beliefs.*

We provide a sketch of proof of Lemma 1 because it helps to understand how dynamic (in)consistency matter for SCE analysis.²⁹ Fix $\bar{\sigma} \in \text{SCE}(\Gamma, f, \bar{\phi})$ and consider a strategy $\bar{s}_i \in \text{Supp} \bar{\sigma}_i$ that is sequentially optimal given $\mu_{\bar{s}_i}$. Since the confirmed-beliefs condition does not depend on ambiguity attitudes, we only have to argue that some $s_i \in [\bar{s}_i]$ is $(\mu_{\bar{s}_i}, \phi_i)$ -unimprovable. Fix any h consistent with \bar{s}_i ($h \in H_i(\bar{s}_i)$). By ex post perfect recall, $\mu_{\bar{s}_i}$ assigns probability 1 to the set of distributions σ_{-i} such that $\sigma_{-i}(S_{-i}(h)) = \bar{\sigma}_{-i}(S_{-i}(h))$, because i observes the frequency of h . There are two cases. (1) If $\bar{\sigma}_{-i}(S_{-i}(h)) > 0$, then $\mu_{\bar{s}_i}(\Sigma_{-i}(h)) > 0$ and conditional belief $\mu_{\bar{s}_i}(\cdot|h)$ is determined by Bayes rule. Since payoffs are observable, according to conditional belief $\mu_{\bar{s}_i}(\cdot|h)$ the equilibrium action $\bar{s}_{i,h}$ is unambiguous (that is, it involves known risks), whereas deviations are untested and can be perceived as ambiguous. Thus, keeping continuation plan and beliefs fixed, an increase in ambiguity aversion from the baseline $\bar{\phi}_i$ to the more concave ϕ_i decreases the value of deviations without affecting the value of $\bar{s}_{i,h}$. Moreover, by sequential optimality, after each deviation the original continuation plan described by \bar{s}_i is optimal under $\bar{\phi}_i$ given $\mu_{\bar{s}_i}(\cdot|h)$. This implies that any $(\mu_{\bar{s}_i}, \phi_i)$ -unimprovable continuation plan makes deviations less attractive under $\bar{\phi}_i$ (if it involves further deviations from \bar{s}_i down the road) and hence less attractive than $\bar{s}_{i,h}$ under the higher ambiguity aversion represented by ϕ_i . (2) If $\bar{\sigma}_{-i}(S_{-i}(h)) = 0$, also $\mu_{\bar{s}_i}$ assigns probability 0 to h (that is, $h \notin H_i(\mu_{\bar{s}_i})$) and unimprovability does not impose any optimality requirement on $\bar{s}_{i,h}$ (see Definition 1). Hence, one can find a strategy s_i that is realization equivalent to \bar{s}_i and $(\mu_{\bar{s}_i}, \phi_i)$ -unimprovable. The following example illustrates this intuition.

Example 5 *Let Γ' be the game of Figure 1, but with payoff 5 instead of 4 at outcome Out. The symmetric SCE (Out.M, L) of $(\Gamma, f, \bar{\phi})$ of Example ?? is also a symmetric SCE of $(\Gamma', f, \bar{\phi})$ justified by the same confirmed beliefs as in Examples 1 and ??:*

$$\mu_1(\sigma_2) = \begin{cases} \frac{1}{2} & \text{if } \sigma_2 \in \{\delta_L, \delta_R\}, \\ 0 & \text{otherwise.} \end{cases}$$

Now, Out.M is not only a $(\mu_1, \bar{\phi}_1)$ -unimprovable strategy, but it is also $(\mu_1, \bar{\phi}_1)$ -sequentially optimal, because at the root the best alternative strategy In.T (see Example 3) yields an unambiguous expected payoff of $4.5 < 5$. Consider now any strictly increasing and concave transformation $\phi_1 = \varphi_1 \circ \bar{\phi}_1$ such that:

$$\begin{aligned} \phi_1(u) &= \bar{\phi}_1(u) = \sqrt{u} & \text{if } u \in [0, 9], \\ \bar{\phi}_1(u) &= 6 > \phi_1(36) > \phi_1(9) = 3. \end{aligned}$$

²⁹Note that a sequentially optimal strategy may be suboptimal in the normal form of the game. For this reason, we cannot directly apply the monotonicity result of BCMM to prove the lemma.

Specifically, ϕ_1 (36) is close to 3. In the transformation from $\bar{\phi}_1$ to ϕ_1 , the values of In.M and In.B at the root (given μ_1) decrease, whereas the value of In.T is constant. Then, at the root, In. a_1^ϕ is worse than Out, where a_1^ϕ denotes any action that maximizes player 1's (μ_1, ϕ_1) -value conditional on $\{(In, G)\}$. This implies that Out. a_1^ϕ is (μ_1, ϕ_1) -unimprovable and $(Out.a_1^\phi, L)$ is a symmetric SCE of (Γ', f, ϕ) realization equivalent to $(Out.M, L)$. \blacktriangle

What can go wrong when the SCE strategies are *not* sequentially optimal under the confirmed beliefs that justify them? Take the viewpoint of an agent in population i at some information set $h \in H_i(\mu_i) \cap H(s_i)$, where s_i is the agent's strategy in an SCE of $(\Gamma, \bar{\phi}, f)$ and μ_i is the confirmed belief that justifies it. At h , the agent evaluates a deviation from the equilibrium action $s_{i,h}$ to an alternative action a_i , after which he might play once more at an information set $h' \in H_i(\mu_i)$. Suppose that, given his belief and some action a'_i at h' , the deviation at h has a higher value than the SCE expected payoff. Yet, s_i is not $(\mu_i, \bar{\phi}_i)$ -sequentially optimal, the agent also realizes that his "future self" at h' will play action $s_{i,h'}$ different from a'_i , and this makes the agent prefer $s_{i,h}$ to a_i at h . But, as his ambiguity aversion increases from $\bar{\phi}_i$ to ϕ_i , the future self of the agent may switch from $s_{i,h'}$ to a'_i at h' for all the confirmed beliefs that justify $s_{i,h}$ under $\bar{\phi}_i$. Then, although for any fixed belief and action at h' the value of a_i compared to $s_{i,h}$ at h decreases (because a_i exposes the agent to ambiguity while $s_{i,h}$ does not, and the agent has become more ambiguity averse), the value of a_i at h under the predicted choice at h' can increase when moving from $\bar{\phi}_i$ to ϕ_i . The following example demonstrates this possibility.

Example 6 Consider the game of Figure 1. Let $\bar{\phi}_1$ and ϕ_1 be the ambiguity attitudes described in Examples ?? and 5. In Appendix 6 we show that, if ϕ (36) is sufficiently small, Out does not belong to the set of SCE outcomes of (Γ, ϕ, f) (although it does for $(\Gamma, f, \bar{\phi})$). The intuition is as follows. For an intermediate level of ambiguity aversion, captured by the baseline second-order utility $\bar{\phi}_1$, upon reaching (In, G) player 1 is tempted by actions M and B even under the most pessimistic belief $(\mu_1(\cdot | \{(In, G)\}) = \frac{1}{2}\delta_L + \frac{1}{2}\delta_R)$, cf. Lemma 9). At the root, he anticipates this and, scared by the implied ex-ante objective expected reward of $\frac{1}{2}$ under the "bad model" (δ_L if he plans M and δ_R if he plans B), he chooses Out. For a higher level of ambiguity aversion, captured by ϕ_1 , at (In, G) player 1 is tempted by M and B only for sufficiently optimistic beliefs. As a consequence, at the root, he is less worried by the "bad model," either because now he plans the unambiguous action T for the subgame, or because he plans an ambiguous action and he deems the "bad model" sufficiently unlikely. Therefore, he chooses In at the root. \blacktriangle

Thus, the monotonicity result of BCMM does not extend to all sequential games, even if we restrict our attention to distributions of outcomes.³⁰ Yet, we can use Lemma 1 to show that the monotonicity result holds for classes of games and equilibria of interest.

³⁰This shows that the conjecture informally stated by BCMM (p. 667) is false.

6.1 No player moves more than once

We say that **no player moves more than once**³¹ in Γ if, for every $i \in I$ and $z \in Z$, there is at most one information set $h \in H_i$ that contains a predecessor of z . In this class of games, at any information set $h \in H_i$, the agent does not move again after h ; thus, no information set of i is prevented by any strategy of i , and classes of realization equivalent strategies are singletons. Therefore, the value of an action at $h \in H_i$ does not depend on i 's strategy, unimprovability coincides with sequential optimality,³² and Lemma 1 implies the following result:

Corollary 4 *Fix two games with observable payoffs where no player moves more than once, (Γ, f, ϕ) and $(\Gamma, f, \bar{\phi})$, so that (Γ, f, ϕ) features more ambiguity aversion than $(\Gamma, f, \bar{\phi})$. Then*

$$SCE(\Gamma, f, \phi) \supseteq SCE(\Gamma, f, \bar{\phi}).$$

6.2 Symmetric self-confirming equilibria

Consider now a *symmetric* SCE \bar{s} in a game with observable payoffs and *without chance moves*. We can show that there is a symmetric SCE equilibrium s with the same path.³³

Theorem 5 *Fix two games with observable payoffs, ambiguity aversion, and without chance moves, (Γ, f, ϕ) and $(\Gamma, f, \bar{\phi})$, so that (Γ, f, ϕ) features more ambiguity aversion than $(\Gamma, f, \bar{\phi})$. Then:*

$$\zeta(\text{symSCE}(\Gamma, f, \phi)) \supseteq \zeta(\text{symSCE}(\Gamma, f, \bar{\phi})).$$

We prove Theorem 5 in the Appendix, here we provide an intuitive argument. The SCE strategies may not be sequentially optimal under the confirmed beliefs that justify them. However, consider alternative beliefs that are supported by Dirac models and give the same predictive beliefs as the original ones. By construction, these beliefs are confirmed by the equilibrium play, and they feature two additional properties. First, they are the most pessimistic beliefs among those that give rise to the same predictive probabilities; thus, by certainty of the equilibrium payoff, they justify a realization equivalent symmetric SCE. This is shown by Lemma 9 in the Appendix and it is based on the following intuition: A belief supported by Dirac models polarizes the objective expected payoffs that the agent deems possible, and the concavity of ϕ_i implies that player i is averse to such polarization. In other words, the transformation of a belief into a “prediction-equivalent” belief supported by Dirac models has the same effect as a mean-preserving spread in a

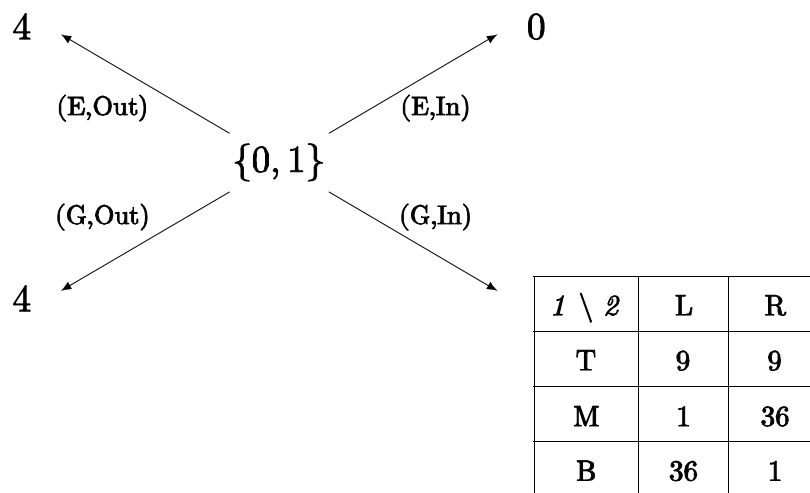
³¹Perfect information games with this property have been called “simple” (see, e.g., Fudenberg and Levine, 2006). We prefer the more explicative phrase because (i) we are not assuming perfect information and (ii) “simple” is vague and had been used earlier with a different meaning (see Battigalli, 2003).

³²Indeed, Proposition 2 can be extended to any ϕ_i in one-move games; yet, Proposition 1 cannot!

³³Note that we can identify $\text{symSCE}(\Gamma, f_i, \bar{\phi})$ with a subset of S . Hence, it makes sense to write $\zeta(\text{symSCE}(\Gamma, f_i, \bar{\phi}))$.

decision problem under risk aversion. Second, absent chance moves, beliefs supported by Dirac models cannot entail dynamic inconsistencies of preferences over strategies; thus, the symmetric SCE's they justify features sequentially optimal strategies. This is shown by Lemma 8 in the Appendix. The main idea behind this result is that dynamic inconsistency is due to the perception of hedging opportunities that may be optimal ex-ante, but would not be implemented ex-post after (partial) resolution of the uncertainty. But, absent chance moves, an agent whose belief is supported by Dirac models does not perceive such hedging opportunities. With this, we can use Lemma 1 to prove the monotonicity result for equilibrium outcomes.

Absence of chance moves and symmetry are tight conditions. In Example 6, outcome (Out) is induced by a symSCE of the game with chance moves $(\Gamma, f, \bar{\phi})$, but it is not induced by any symSCE of the game (Γ, f, ϕ) , which features more ambiguity aversion. As for the role of symmetry (pure equilibrium), consider the following example.



A 3-person common interest game: $I = \{0, 1, 2\}$

Example 7 Let Γ' be the following modification of the game of Figure 1: deviating from our standard notation, here 0 is not chance but rather an actual player, choosing between E and G simultaneously with player 1 at the root. Assuming common interests, all players obtain payoff 4 at both outcomes (E, Out) and (G, Out). See Figure 6.2. Assume perfect feedback and consider the same $\bar{\phi}_1$ and ϕ_1 of Example 6. The SCE $\bar{\sigma} = (\sigma_0, \text{Out.M, L})$ of $(\Gamma', f, \bar{\phi})$ yields outcomes (E, Out) and (G, Out) with probability $\frac{1}{2}$ (cf. Example ??). By perfect feedback and confirmed beliefs, in any realization-equivalent SCE player 1 is certain of σ_0 . Yet, for the same argument as for Example 6, no (μ_1, ϕ_1) -unimprovable strategy of player 1 prescribes action Out when the marginal of μ_1 on Σ_0 is σ_0 . Thus, no SCE of (Γ', f, ϕ) yields the same outcome distribution as $\bar{\sigma}$. \blacktriangle

6.3 Ambiguity aversion versus ambiguity neutrality

Fix a sequential game with feedback (Γ, f) and let $\text{Id}_{\mathbb{V}} = (\text{Id}_{\mathbb{V}_i})_{i \in I}$ denote the profile of players' identity functions characterizing their neutrality toward ambiguity. We relate the set $SCE(\Gamma, f, \text{Id}_{\mathbb{V}})$ of SCE's of Γ given feedback functions f and neutral ambiguity attitudes with the set $SCE(\Gamma, f, \phi)$ with non-neutral ambiguity attitudes ϕ . We start with a preliminary observation:³⁴

Remark 3 *For every two-person game with feedback (Γ, f) and every profile ϕ of second-order utility functions, $SCE(\Gamma, f, \text{Id}_{\mathbb{V}}) \subseteq SCE(\Gamma, f, \phi)$.*

To see this, note that ambiguity neutrality and the convexity of $\Sigma_{-i} = \Delta(S_{-i})$ in two-person games allow to replace the justified beliefs supporting $\bar{\sigma}$ as an SCE of $(\Gamma, f, \text{Id}_{\mathbb{V}})$ with the corresponding Dirac beliefs supported by their predictive measure. Since ambiguity attitudes are immaterial for agents with Dirac beliefs, $\bar{\sigma}$ is also an SCE of (Γ, f, ϕ) . This argument does not hold with $n > 2$ players, because in this case Σ_{-i} is not convex, hence, the Dirac measure supported by the predictive of $\mu_{\bar{s}_i}$ may belong to $\Delta(\Delta(S_{-i}) \setminus \Delta(\Sigma_{-i}))$. Nonetheless, we can relate $SCE(\Gamma, f, \text{Id}_{\mathbb{V}})$ and $SCE(\Gamma, f, \phi)$ for a large class of games.

As a corollary of their main monotonicity result, BCMM show that, under observable payoffs, $SCE(\mathcal{N}(\Gamma, f), \text{Id}_{\mathbb{V}})$ is contained in $SCE(\mathcal{N}(\Gamma, f), \phi)$ if each ϕ_i is concave and there are observable payoffs. Even if the main monotonicity result of BCMM does not extend to SCE of sequential games for the entire spectrum of ambiguity attitudes, we are still able to obtain a sequential version of this corollary in terms of induced outcome distributions. By dynamic consistency under ambiguity neutrality (i.e., by Proposition 2), this result is a corollary of Lemma 1. Define the function

$$\begin{aligned} \hat{\zeta} : \Sigma &\rightarrow \Delta(Z) \\ \bar{\sigma} &\mapsto \hat{\zeta}(\bar{\sigma})(z) = \sum_{(s_0, s) : \zeta(s_0, s) = z} \sigma_0(s_0) \cdot \bar{\sigma}(s), \end{aligned}$$

where Σ is the set of product measures on S . This is the pushforward map that gives for every (product) distribution over strategy profiles $\bar{\sigma}$ the corresponding probability distribution $\hat{\zeta}(\bar{\sigma})$ on terminal nodes.

Corollary 6 *Suppose that (Γ, f) has observable payoffs and (Γ, f, ϕ) features ambiguity aversion. Then, the set of SCE distributions over terminal nodes of (Γ, f, ϕ) contains the set of SCE distributions over terminal nodes with ambiguity neutrality:*

$$\hat{\zeta}(SCE(\Gamma, f, \phi)) \supseteq \hat{\zeta}(SCE(\Gamma, f, \text{Id}_{\mathbb{V}})).$$

³⁴Cf. footnote 23 of BCMM.

Given the large body of empirical evidence supporting the ambiguity aversion hypothesis, we conclude that the standard SCE concept, which implicitly assumes neutral ambiguity attitudes, overestimates the predictability of long-run outcomes of learning dynamics.

Finally, note that Corollary 4 and Corollary 6, via Lemma 1, imply that all beliefs that justify $SCE(\Gamma, f, \bar{\phi})$ in one-move games or under ambiguity neutrality (linear $\bar{\phi}$) also justify the corresponding equilibria under ϕ . This is not true for symmetric equilibria in games without chance moves (see the proof of Theorem 5).

7 Conditional probability systems and full unimprovability

So far we studied how ambiguity aversion and the ensuing possibility of dynamic inconsistency (the incentive to covertly commit, if possible) affect SCE analysis. However, we neglected strategic reasoning based on common knowledge of (some features of) the game. To analyze strategic reasoning in sequential games we have to address an additional and more traditional issue: We need to model how a player thinks that the other agents would react to unexpected moves. Even if players are ambiguity neutral, the analysis of Sections 4-6 is insufficient to address this issue: As a preliminary step, we need to assume that agents have well defined conditional beliefs at all information sets, including the unexpected ones. The conditional beliefs of any given agent at different information sets have to be mutually consistent if we want to preserve/extend the unimprovability principle. Moreover, we want to model players who reason strategically about the game before playing it. For this purpose, we have to consider also their beliefs at the root, even if they are not first movers. Therefore, for every $i \in I$, we will consider the expanded collection of information sets $\bar{H}_i = H_i \cup \{\{\emptyset\}\}$.

To simplify the analysis of this section and the following one we focus on games *without chance moves*. Thus, the outcome function is $\zeta : S \rightarrow Z$, the strategic-form payoff feedback and payoff functions of player i are $F_i = f_i \circ \zeta : S \rightarrow M$ and $U_i = v_i \circ \gamma \circ \zeta : S \rightarrow \mathbb{R}$.³⁵

To understand the following definition, consider beliefs $\mu_i(\cdot|g), \mu_i(\cdot|h) \in \Delta(\Sigma_{-i})$ at two information sets g and h , so that g precedes h ($g \prec h$), and suppose that h is possible according to $\mu_i(\cdot|g)$, that is, $p_{\mu_i}(S_{-i}(h)|g) > 0$. Then, the conditional belief $\mu_i(\cdot|h)$ can be derived from $\mu_i(\cdot|g)$ in the way prescribed by (5): for all $E_{-i} \in \mathcal{B}(\Sigma_{-i})$,

$$\mu_i(E_{-i}|h) = \frac{\int_{E_{-i} \cap \Sigma_{-i}(g)} \sigma_{-i}(S_{-i}(h)|g) \mu_i(d\sigma_{-i}|g)}{p_{\mu_i}(S_{-i}(h)|g)}.$$

Next note that we can write the required relation between $\mu_i(\cdot|g)$ and $\mu_i(\cdot|h)$ without explicitly stating condition $p_{\mu_i}(S_{-i}(h)|g) > 0$: for all $g, h \in \bar{H}_i$ with $g \prec h$, and $E_{-i} \in$

³⁵The analysis can be extended to games with chance moves, but some the notation and some of the proofs would be more complex.

$\mathcal{B}(\Sigma_{-i})$

$$\mu_i(E_{-i}|h)p_{\mu_i}(S_{-i}(h)|g) = \int_{E_{-i} \cap \Sigma_{-i}(g)} \sigma_{-i}(S_{-i}(h)|g)\mu_i(d\sigma_{-i}|g). \quad (8)$$

Definition 6 A *conditional probability system (CPS)* on (Σ_{-i}, \bar{H}_i) is an array of probability measures

$$\mu_i(\cdot|\cdot) = (\mu_i(\cdot|h))_{h \in \bar{H}_i} \in [\Delta(\Sigma_{-i})]^{\bar{H}_i}$$

such that

- (1) for all $h \in \bar{H}_i$, $\mu_i(\Sigma_{-i}(h)|h) = 1$,
- (2) for all $g, h \in \bar{H}_i$ with $g \prec h$ eq. (8) holds.

Definition 7 A *CPS* on (S_{-i}, \bar{H}_i) is an array of probability measures

$$p_i(\cdot|\cdot) = (p_i(\cdot|h))_{h \in \bar{H}_i} \in [\Delta(S_{-i})]^{\bar{H}_i}$$

such that $(p_i(\cdot|h))_{h \in \bar{H}_i} = (p_{\mu_i(\cdot|h)})_{h \in \bar{H}_i}$ for some CPS $\mu_i(\cdot|\cdot) = (\mu_i(\cdot|h))_{h \in \bar{H}_i}$ on (Σ_{-i}, \bar{H}_i) .

Comments In Definition 7, we define a CPS on (S_{-i}, \bar{H}_i) indirectly as the “predictive” of some CPS on (Σ_{-i}, \bar{H}_i) . We could have given a direct definition: Indeed, $p_i \in [\Delta(S_{-i})]^{\bar{H}_i}$ is a CPS on (S_{-i}, \bar{H}_i) if and only if (1) for every $h \in \bar{H}_i$, $p_i(S_{-i}(h)|h) = 1$ and (2) for all $g, h \in \bar{H}_i$ with $g \prec h$ and all $s_{-i} \in S_{-i}(h)$,

$$p_i(s_{-i}|h)p_i(S_{-i}(h)|g) = p_i(s_{-i}|g). \quad (9)$$

Notation 7 The set of CPS’s on (Σ_{-i}, \bar{H}_i) $[(S_{-i}, \bar{H}_i)]$ is denoted by $\Delta^{\bar{H}_i}(\Sigma_{-i})$ $[\Delta^{\bar{H}_i}(S_{-i})]$.

With this, we can give a stronger definition of unimprovability:

Definition 8 A strategy s_i is **fully** $(\mu_i(\cdot|\cdot), \phi_i)$ -*unimprovable* (where $\mu_i(\cdot|\cdot) \in \Delta^{\bar{H}_i}(\Sigma_{-i})$) if

$$\forall h \in H_i, s_{i,h} \in \arg \max_{a_i \in A_i(h)} V_i(a_i|h; s_i, \mu_i, \phi_i).$$

It is well known that, for ambiguity neutral agents, we have a refined dynamic programming result:³⁶

³⁶Definition 8 requires maximization over actions; hence, it considers only the information sets where i is active. The following propositions relate to maximization over strategies; hence, they also include (ex ante) maximization at $\{\emptyset\}$ even if i is not a first mover. Also, recall that, according to our notation, $(s_i^*|h, s_{i,h}^*)$ is the minimal modification of s_i^* that makes h reachable and plays like s_i^* at all information sets h' that do not strictly precede h .

Proposition 8 For every $s_i^* \in S_i$ and $\mu_i(\cdot|\cdot) \in \Delta^{\bar{H}_i}(\Sigma_{-i})$, the following are true:
(1*) $[s_i^*]$ contains a fully $(\mu_i(\cdot|\cdot), \text{Id}_{V_i})$ -unimprovable strategy if and only if

$$\forall h \in \bar{H}_i(s_i^*), s_i^* \in \arg \max_{s_i \in S_i(h)} \sum_{s_{-i} \in S_{-i}(h)} U_i(s_i, s_{-i}) p_{\mu_i}(s_{-i}|h);$$

(2*) s_i^* is a fully $(\mu_i(\cdot|\cdot), \text{Id}_{V_i})$ -unimprovable strategy if and only if

$$\forall h \in \bar{H}_i, (s_i^*|h, s_{i,h}^*) \in \arg \max_{s_i \in S_i(h)} \sum_{s_{-i} \in S_{-i}(h)} U_i(s_i, s_{-i}) p_{\mu_i}(s_{-i}|h).$$

Full-improvability solves the difficulties described above. Intuitively, we can interpret a $(\mu_i(\cdot|\cdot), \phi_i)$ -fully unimprovable strategy s_i^* as the **plan** of an agent in role i , which can be obtained with a “folding back” dynamic programming procedure on the subjective decision tree implied by CPS $\mu_i(\cdot|\cdot)$. In our perspective, the fact that the beliefs of an agent with perfect recall are given by a CPS reflects the **epistemic unity of the agent’s self**: the agent always incorporates new information into his system of knowledge and beliefs in a way that is consistent with his previous beliefs and with the rules of conditional probabilities, even when new information follows previously unexpected information.

Before we move on to model strategic reasoning, we verify that even if SCE is strengthened by requiring full unimprovability, the set of possible outcomes does not change.

Definition 9 A **fully unimprovable self-confirming equilibrium** of (Γ, f, ϕ) is a profile of strategy distributions $(\sigma_i^*)_{i \in I}$ with following property: For every $i \in I$ and $s_i^* \in \text{Supp} \sigma_i^*$ there is a CPS $\mu_{s_i^*}(\cdot|\cdot) \in \Delta^{\bar{H}_i}(\Sigma_{-i})$ such that
(rationality) s_i^* is fully $(\mu_{s_i^*}(\cdot|\cdot), \phi_i)$ -unimprovable,
(confirmed beliefs) $\mu_{s_i^*} \left(\left\{ \sigma_{-i} \in \Sigma_{-i} : \hat{F}_i(s_i^*, \sigma_{-i}) = \hat{F}_i(s_i^*, \sigma_{-i}^*) \right\} \mid \{\emptyset\} \right) = 1$.

Note that the confirmed beliefs condition refers only to the initial beliefs $\mu_{s_i^*}(\cdot|\{\emptyset\})$ because it implies that conditional beliefs are confirmed by observed conditional frequencies at every history that is reached with positive probability in equilibrium (see the proof of Lemma 4). By inspection of Definitions 5 and 9 it is also clear that every fully unimprovable SCE is also an SCE. The following proposition implies that SCE and fully unimprovable SCE are realization equivalent.

Proposition 9 For every SCE there is a corresponding fully unimprovable SCE that yields the same probability distribution on terminal nodes.

Intuitively, since SCE does not rely on complete information and does not model strategic thinking, requiring rational reactions to unexpected moves adds little to the analysis: Agents are not assumed to know the preferences of others, hence they are not assumed to rule out irrational reactions to deviations.

8 Knowledge of the game and strategic reasoning

What if (Γ, f, ϕ) (or a part of it) is common knowledge? Then it makes sense to explore a notion of “rationalizable SCE” according to which agents reason strategically about their opponents taking into account their preferences and feedback functions (cf. Rubinstein and Wolinsky 1994, Battigalli 1999, Dekel *et al.* 1999, Esponda 2013, Fudenberg and Kamada 2015). We illustrate this with two simple examples.

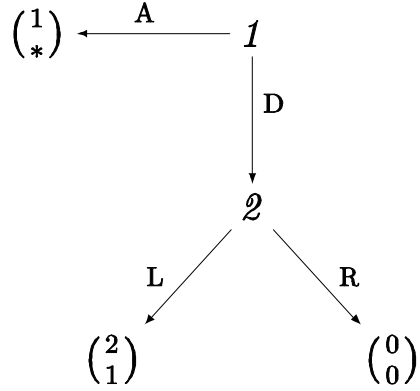


Figure 2: Entry game

Example 8 (Cf. Fudenberg and Levine, 1993). Figure 2 depicts the so called “Entry game” often used to illustrate the shortcomings of the Nash equilibrium concept. To complete the specification of (Γ, f, ϕ) assume observable payoffs (hence, in this case, perfect feedback) and let ϕ be an arbitrary pair of strictly increasing functions.³⁷ It can be checked that the set of SCE’s of (Γ, f, ϕ) is

$$\{\sigma^* \in \Sigma : \sigma_1^*(A) = 1, \text{ or } (\sigma_1^*(A) < 1, \sigma_2^*(L) = 1)\}.$$

First-movers who choose A (go Across) get no feedback and can thus hold trivially confirmed beliefs that make them play A. As for second-movers, their plan becomes relevant only if a positive fraction of first-movers choose D (go Down). In that case, the only rational choice is L (Left). In equilibrium, first-movers going Down correctly predict Left, which makes Down a strict best response. How can a first mover expect that the second mover goes Right with high probability (i.e., that a large fraction of second movers play Right)? Informally, this is possible if either the first mover gives a high probability to the second mover being irrational, or—more reasonably—if the first mover does not know v_2 . If instead the first mover knows v_2 and believes in the rationality of the second mover, then he predicts that

³⁷One can show that ambiguity aversion does not matter when players have only two strategies.

Down would be followed by Left, and he would go Down. Intuitively, only one SCE is consistent with belief in rationality and knowledge of the game: (D,L). However, this is not formally captured by the confirmed beliefs condition of Definition 5. According to such definition, if an agent in population 1 does not play Down, he can keep the belief that the second mover would play Right after Down: If v_2 is unknown to first movers, they have no way to understand that only Left is rational. Relatedly, the set of equilibrium distributions on terminal nodes is unchanged if we adopt the stronger definition of unimprovability: We must have $\sigma_2^*(L) = 1$ in this case, but those first movers who go Across because they have wrong beliefs have no way to find out they are wrong. ▲

Example 9 (Rubinstein and Wolinsky 1994, Battigalli 1999). Two players must simultaneously choose a location among the equally spaced points 0, 1, 2, 3 on the real line. Their payoff is the negative of the distance between their chosen locations and each player, that is, $U_i(s_1, s_2) = -|s_1 - s_2|$; furthermore, each player observes this distance and, of course, remembers his action: $F_i(s_1, s_2) = (s_i, |s_1 - s_2|)$. Again, let us fix ϕ arbitrarily and—for simplicity—let us focus on the symmetric equilibria. The set of symmetric SCE’s is

$$\text{symSCE} = \{(0, 0), (1, 1), (2, 2), (3, 3), (1, 2), (2, 1)\}.$$

Of course, (1,2) and (2,1) are not Nash equilibria. They are symmetric SCE’s supported by confirmed beliefs according to which each agent assigns the same probability to the co-player being either on his right, or on his left. However, if the game is common knowledge, an agent with such beliefs infers that with probability .5 his co-player is “cornered” at an extreme point and cannot be best responding to confirmed beliefs. This is unlikely to be a stable situation under complete information if players reason strategically, because no player has reasons to believe that the co-player would keep playing in the same way. ▲

Symmetrically rationalizable SCE We capture with an inductive definition the behavioral consequences of the following assumptions on rationality and interactive beliefs: (1) agents are rational (in the sense of full unimprovability) and have confirmed beliefs, and (2) there is common belief at the beginning of the game that (1) holds. This is easier to do in a particular case, i.e., when all agents in the same population follow the same plan (symmetric SCE) and there is common full belief of this.³⁸ In this case, the CPS $\mu_i(\cdot|\cdot)$ on (Σ_{-i}, \bar{H}_i) is supported by Dirac models, and therefore it is isomorphic to a CPS p_{μ_i} on (S_{-i}, \bar{H}_i) . A variation of the algorithm defined by Battigalli (1999) (see also Esponda 2013) precisely captures the foregoing epistemic assumptions. As in Section 7, also in this section we assume for simplicity that there are no chance moves.

³⁸Note, we did not say “belief at the beginning of the game,” because now we mean “probability 1 belief conditional on every information set.” This is called “full belief” in epistemic game theory (e.g., Battigalli and Prestipino, 2013).

Given a CPS $\mu_i(\cdot|\cdot) \in \Delta(\Sigma_{-i})^{\bar{H}_i}$, we let

$$r_i(\mu_i(\cdot|\cdot), \phi_i) = \left\{ s_i \in S_i : \forall h \in H_i, s_{i,h} \in \arg \max_{a_i \in A_i(h)} V_i(a_i|h; s_i, \mu_i(\cdot|\cdot), \phi_i) \right\}$$

denote the set of fully (μ_i, ϕ_i) -unimprovable strategies of i . When $\mu_i(\cdot|\cdot)$ is isomorphic to a predictive CPS $p_i(\cdot|\cdot)$, as in the case we are considering right now, it makes sense to write $r_i(p_i(\cdot|\cdot), \phi_i)$.

Definition 10 For each $i \in I$, let $B_i^0 = S_i \times M$, and

$$B_i^{k+1} = \left\{ (\bar{s}_i, \bar{m}_i) \in B_i^k : \begin{array}{l} \exists p_i(\cdot|\cdot) \in \Delta^{\bar{H}_i}(S_{-i}), \bar{s}_i \in r_i(p_i(\cdot|\cdot), \phi_i), \\ p_i(F_{\bar{s}_i}^{-1}(\bar{m}_i) \cap \{s_{-i} : (s_j, F_j(\bar{s}_i, s_{-i}))_{j \in I \setminus \{i\}} \in B_{-i}^k\} | \{\emptyset\}) = 1 \end{array} \right\}$$

for each $k \in \mathbb{N}_0$. A strategy profile \bar{s} is a **symmetrically rationalizable SCE** for game (Γ, f, ϕ) (without chance moves) if $(\bar{s}_i, F_i(\bar{s}))_{i \in I} \in \times_{i \in I} \bigcap_{k \in \mathbb{N}} B_i^k$. The set of symmetrically rationalizable self-confirming equilibria of (Γ, f, ϕ) is denoted by $\text{symRSCE}(\Gamma, f, \phi)$.

Intuitively, $(\bar{s}_i, \bar{m}_i) \in B_i^1$ if \bar{s}_i is justified by some CPS such that i is initially certain to get message \bar{m}_i if he plays \bar{s}_i . Thus, $(\bar{s}_i, F_i(\bar{s}))_{i \in I} \in \times_{i \in I} \bigcap_{k \in \mathbb{N}} B_i^1$ if \bar{s} is an unimprovable SCE (with beliefs supported by Dirac models). Then, $(\bar{s}_i, \bar{m}_i) \in B_i^2$ if \bar{s}_i is justified by some CPS such that i is initially certain to get message \bar{m}_i if he plays \bar{s}_i , and furthermore he is initially certain that everybody else is playing strategies justified by confirmed beliefs. The iterations capture higher levels of (initial) mutual belief in rationality and in confirmation of beliefs.

Remark 4 A strategy profile \bar{s} is a **symmetrically rationalizable SCE** for (Γ, f, ϕ) if and only if there is a profile of finite subsets $(\bar{B}_i)_{i \in I} \in \times_{i \in I} 2^{S_i \times M}$ such that, for every $i \in I$, there is $\bar{m}_i \in M$ with $(\bar{s}_i, \bar{m}_i) \in \bar{B}_i$, and for every $(\hat{s}_i, \hat{m}_i) \in \bar{B}_i$, there is $p_i(\cdot|\cdot) \in \Delta^{\bar{H}_i}(S_{-i})$ with $\hat{s}_i \in r_i(p_i(\cdot|\cdot), \phi_i)$ and

$$p_i(F_{\hat{s}_i}^{-1}(\hat{m}_i) \cap \{s_{-i} : (s_j, F_j(\hat{s}_i, s_{-i}))_{j \neq i} \in \bar{B}_{-i}\} | \{\emptyset\}) = 1. \quad (10)$$

The Remark above shows that our inductive definition is an extensive-form version with ambiguity attitudes of Rubinstein's and Wolinsky's (1994) rationalizable conjectural equilibrium. Next we offer a characterization with sets of strategy profiles. Let $\text{symSCE}^0(\Gamma, f, \phi) = S$ and

$$\text{symSCE}^{k+1}(\Gamma, f, \phi) = \left\{ \bar{s} \in \text{symSCE}^k(\Gamma, f, \phi) : \begin{array}{l} \forall i \in I, \exists p_i(\cdot|\cdot) \in \Delta^{\bar{H}_i}(S_{-i}), \bar{s}_i \in r_i(p_i(\cdot|\cdot), \phi_i), \\ p_i(F_{\bar{s}_i}^{-1}(F_i(\bar{s})) \cap \text{symSCE}^k(\Gamma, f, \phi)_{\bar{s}_i} | \{\emptyset\}) = 1 \end{array} \right\},$$

where, for any subset $X \subseteq S$ and strategy $\bar{s}_i \in S_i$,

$$X_{\bar{s}_i} = \{s_{-i} \in S_{-i} : (\bar{s}_i, s_{-i}) \in X\}$$

is the section of X at \bar{s}_i (thus, $\text{symSCE}^k(\Gamma, f, \phi)_{\bar{s}_i}$ is the section of $\text{symSCE}^k(\Gamma, f, \phi)$ at \bar{s}_i). Note that $\text{symSCE}^1(\Gamma, f, \phi)$ coincides with the set of fully unimprovable symmetric SCE's of (Γ, f, ϕ) justified by confirmed beliefs supported by Dirac models. Thus, $\bar{s} \in \text{symSCE}^2(\Gamma, f, \phi)$ if each \bar{s}_i is a best reply to a confirmed belief that assigns probability 1 to other players choosing best replies to confirmed beliefs supported by Dirac models, given \bar{s}_i and message $m_i = F_i(\bar{s})$.

The following result shows that symmetric RSCE is characterized by the iterated deletion of strategy profiles $(\text{symSCE}^k(\Gamma, f, \phi))_{k \in \mathbb{N}}$:

Lemma 2 *For every $k \in \mathbb{N}$, and $\bar{s} \in S$, $\bar{s} \in \text{symSCE}^k(\Gamma, f, \phi)$ if and only if $(\bar{s}_i, F_i(\bar{s}))_{i \in I} \in B^k$. Therefore,*

$$\bigcap_k \text{symSCE}^k(\Gamma, f, \phi) = \text{symRSCE}(\Gamma, f, \phi).$$

Example 10 *One can easily check that the only rationalizable SCE for the game of Figure C is (D,L), independently of ϕ : Only L is fully unimprovable for player 2, thus $\text{symSCE}^1(\Gamma, f, \phi) = \{(A, L), (D, L)\}$ and $\text{symSCE}^1(\Gamma, f, \phi)_A = \{L\}$. Therefore,*

$$\text{symSCE}^2(\Gamma, f, \phi) = \{(D, L)\}.$$

▲

Example 11 *Now consider the rationalizable SCE's set for the distance game of Example 9:*

$$\text{symSCE}^1(\Gamma, f, \phi) = \text{symSCE}(\Gamma, f, \phi) = \{(0, 0), (1, 1), (2, 2), (3, 3), (1, 2), (2, 1)\}.$$

Let $\bar{s} = (1, 2)$. We show that $\bar{s} \notin \text{symSCE}^2(\Gamma, f, \phi)$. Consider player 1. To have $\bar{s} \in \text{symSCE}^2(\Gamma, f, \phi)$ we have to justify $\bar{s}_1 = 1$ as a best reply to a belief p_1 such that $p_1(\{0, 2\} \cap \text{symSCE}^1(\Gamma, f, \phi)_{\bar{s}_1}) = 1$. But $\text{symSCE}^1(\Gamma, f, \phi)_{\bar{s}_1} = \{1, 2\}$. Thus $\{0, 2\} \cap \text{symSCE}^1(\Gamma, f, \phi)_{\bar{s}_1} = \{2\}$, and the best reply to $s_2 = 2$ is $s_1 = 2$. A similar argument shows that $(2, 1) \notin \text{symSCE}^1(\Gamma, f, \phi)$. Hence only the pure Nash equilibria are symmetrically rationalizable self-confirming. ▲

We can prove two monotonicity results for the symRSCE correspondence:

Proposition 10 Theorem 11 *Fix two games with observable payoffs and no chance moves where no player moves more than once, (Γ, f, ϕ) and $(\Gamma, f, \bar{\phi})$, so that (Γ, f, ϕ) features more ambiguity aversion than $(\Gamma, f, \bar{\phi})$. Then,*

$$\text{symRSCE}(\Gamma, f, \bar{\phi}) \subseteq \text{symRSCE}(\Gamma, f, \phi).$$

Theorem 12 Fix two games with ambiguity aversion, observable payoffs, and no chance moves, (Γ, f, ϕ) and $(\Gamma, f, \bar{\phi})$, so that (Γ, f, ϕ) features more ambiguity aversion than $(\Gamma, f, \bar{\phi})$. Then, for every $\bar{s} \in \text{symRSCE}(\Gamma, f, \bar{\phi})$ there is some $s \in \text{symRSCE}(\Gamma, f, \phi)$ such that, for each player i , s_i is realization equivalent to \bar{s}_i ; therefore,

$$\zeta(\text{symRSCE}(\Gamma, f, \bar{\phi})) \subseteq \zeta(\text{symRSCE}(\Gamma, f, \phi)).$$

Intuitively these results rely on the following intermediate step:³⁹

Monotonicity of the justifiability correspondence: For every information set $h \in H_i$, if an action a_i is justified as a best reply to some conditional belief $\bar{p}_{i,h} \in \Delta(S_{-i}(h))$ given baseline ambiguity attitudes $\bar{\phi}_i$, then there is some $p_{i,h} \in \Delta(S_{-i}(h))$ that justifies a_i as a best reply given the more ambiguity averse attitudes ϕ_i .

Then we can show that if a strategy is fully $(\bar{p}_i, \bar{\phi}_i)$ -unimprovable, then it is also fully (p_i, ϕ_i) -unimprovable for some suitably chosen belief system p_i that coincides with \bar{p}_i on the equilibrium path, and can be chosen off the equilibrium path invoking monotonicity of the justifiability correspondence.

Rationalizable SCE We now define rationalizable SCE for population games, that is, we consider the general non-symmetric version of rationalizable SCE. By analogy with Lemma 2, we perform an iterated deletion of distributions of strategy profiles.

To ease notation, let

$$\hat{\Sigma}_{-i}(s_i, \bar{\sigma}_{-i}) = \left\{ \sigma_{-i} \in \Sigma_{-i} : \hat{F}_i(s_i, \sigma_{-i}) = \hat{F}_i(s_i, \bar{\sigma}_{-i}) \right\}$$

denote the partially identified set of co-players strategy distributions observationally equivalent for i to $\bar{\sigma}_{-i}$ given s_i . With this, let $SCE^0(\Gamma, f, \phi) = \times_{i \in I} \Delta(S_i)$, and

$$SCE^{k+1}(\Gamma, f, \phi) = \left\{ \bar{\sigma} \in SCE^k(\Gamma, f, \phi) : \begin{array}{l} \forall i \in I, \forall s_i \in \text{Supp} \bar{\sigma}_i, \exists \mu_{s_i}(\cdot | \cdot) \in \Delta^{\bar{H}_i}(\Sigma_{-i}), s_i \in r_i(\mu_{s_i}(\cdot | \cdot), \phi_i), \\ \mu_{s_i} \left(\hat{\Sigma}_{-i}(s_i, \bar{\sigma}_{-i}) \cap \text{proj}_{\Sigma_{-i}} SCE^k(\Gamma, f, \phi) | \{\emptyset\} \right) = 1 \end{array} \right\}$$

for every $k \in \mathbb{N}_0$.⁴⁰ Note that $SCE^1(\Gamma, f, \phi)$ is the set of fully unimprovable SCE's of (Γ, f, ϕ) .

³⁹See Lemma 10 and its corollary in the Appendix; cf. Battigalli et al. (2016a) and Weinstein (2016).

⁴⁰Note the following slight abuse of notation. If, for some $k \in \mathbb{N}_0$, $s_i \in S_i$, $\bar{\sigma}_{-i} \in \Sigma_{-i}$ the set $\hat{\Sigma}_{-i}(s_i, \bar{\sigma}_{-i}) \cap \text{proj}_{\Sigma_{-i}} SCE^k(\phi)$ is not measurable, then the requirement

$$\mu_{s_i} \left(\hat{\Sigma}_{-i}(s_i, \bar{\sigma}_{-i}) \cap \text{proj}_{\Sigma_{-i}} SCE^k(\phi) | \{\emptyset\} \right) = 1$$

means that there is a measurable set $\Sigma' \subset \left(\hat{\Sigma}_{-i}(s_i, \bar{\sigma}_{-i}) \cap \text{proj}_{\Sigma_{-i}} SCE^k(\phi) \right)$ with $\mu_{s_i}(\Sigma' | \{\emptyset\}) = 1$.

Definition 11 A profile of strategy distributions $\bar{\sigma} \in \Sigma$ is a rationalizable self-confirming equilibrium for (Γ, f, ϕ) if

$$\bar{\sigma} \in \bigcap_{k \in \mathbb{N}} SCE^k(\Gamma, f, \phi).$$

The set of rationalizable self-confirming equilibria of (Γ, f, ϕ) is denoted by $RSCE(\Gamma, f, \phi)$.

The most important difference with respect to the definition of rationalizable symmetric self-confirming equilibrium is best understood by looking at the second step, where agents check if it is possible that others are best responding to confirmed beliefs. If there is common belief of symmetry, an agent playing strategy s_i is certain that all the other agents in population i also use s_i , and this is taken into account when he checks whether the agents in co-players' populations are best responding to correct beliefs. If instead it is understood that different agents in population i may use different strategies, an agent playing s_i may think that only a negligible fraction of other agents in i is playing s_i , therefore, the fact that he is playing s_i does not enter this calculation. Essentially, we are assuming that each agent has a belief about the whole profile of distributions, that is, a belief over $\Sigma = \times_{j \in I} \Delta(\mathcal{S}_j)$. For the purpose of computing best replies, only the marginal over Σ_{-i} matters. But in order to check whether everybody is best replying to confirmed beliefs, the distribution of strategies in population i is crucial (see Fudenberg and Kamada, 2015, 2018).

Say that Γ (hence also (Γ, f, ϕ)) has **observable deviators** if the factorization

$$S(h) = \times_{j \in I} S_j(h)$$

holds for each information set h of each player. Intuitively, this means that if an information set is reached unexpectedly the active player is able to understand who deviated from the expected path (cf. Fudenberg and Levine 1993, and Battigalli and Guaitoli 1998). We can prove the following monotonicity result for the RSCE correspondence:

Theorem 13 Fix two games with observable payoffs and observable deviators where no player moves more than once, (Γ, f, ϕ) and $(\Gamma, f, \bar{\phi})$, so that (Γ, f, ϕ) features more ambiguity aversion than $(\Gamma, f, \bar{\phi})$. Then $RSCE(\Gamma, f, \bar{\phi}) \subseteq RSCE(\Gamma, f, \phi)$.

We assume observable deviators because in the proof we use a version of the previous result about monotonicity of the justifiability correspondence where $S_{-i}(h)$ is replaced by the space of product distributions on $S_{-i}(h)$, which makes sense if $S_{-i}(h)$ is a product set.

9 Discussion

- heterogeneous pop.: $v_{\theta_i}, \phi_{\theta_i}$
- lack of knowledge of parameters, including chance moves

- initially and strongly rationalizable SCE
- relevance of symmetric equilibrium: pure equilibrium and finite populations; repeated games played by impatient players

10 Appendix

10.1 Proof of Proposition 3

Every finite game Γ has a sequential equilibrium in behavioral strategies $\bar{\beta} = (\bar{\beta}_i)_{i \in I}$. For each $i \in I$, let $\bar{\sigma}_i$ denote the mixed strategy associated with $\bar{\beta}_i$ according to Kuhn's (1953) transformation:

$$\forall s_i \in S_i, \bar{\sigma}_i(s_i) = \prod_{h \in H_i} \bar{\beta}_i(s_{i,h}|h).$$

Let $\mu_{\bar{s}_i} = \delta_{\bar{\sigma}_i}$ for every i and $\bar{s}_i \in \text{Supp} \bar{\sigma}_i$. By construction, these beliefs are correct, hence confirmed. To show that $\bar{\sigma}$ as an SCE justified by these confirmed beliefs, we must prove that, for each i , each $\bar{s}_i \in \text{Supp} \bar{\sigma}_i$ is $(\mu_{\bar{s}_i}, \phi_i)$ -unimprovable.

It is well known that every pure strategy in the support of a sequential equilibrium is a sequential best reply to the equilibrium beliefs. Therefore, for every $i \in I$, $h \in H_i$ and \bar{s}_i such that $\prod_{h' \in H_i} \bar{\beta}_i(s_{i,h'}|h') > 0$, that is, for every $\bar{s}_i \in \text{Supp} \bar{\sigma}_i$,

$$\bar{s}_{i,h} \in \arg \max_{a_i \in A_i(h)} \sum_{x \in h} \mathbb{P}_{(\bar{s}_i|h, a_i), \bar{\beta}_{-i}}(x|h) \sum_{z \in Z} \mathbb{P}_{(\bar{s}_i|h, a_i), \bar{\beta}_{-i}, \bar{\beta}_{-i}}(z|x) v_i(\gamma(z)),$$

where $\mathbb{P}_{s_i, \bar{\beta}_{-i}}(\cdot|\cdot)$ denotes the probability of a node conditional on an information set, or an earlier node, determined by the behavioral strategy profile $(s_i, \bar{\beta}_{-i})$ and the probabilities of chance moves. Since $\bar{\sigma}_{-i}$ is by construction realization-equivalent to $\bar{\beta}_{-i}$, $\mathbb{P}_{s_i, \bar{\beta}_{-i}}(\cdot|\cdot) = \mathbb{P}_{s_i, \bar{\sigma}_{-i}}(\cdot|\cdot)$; hence,

$$\forall s_i \in S_i(h), U_i(s_i, \bar{\sigma}_{-i}|h) = \sum_{x \in h} \mathbb{P}_{s_i, \bar{\beta}_{-i}}(x|h) \sum_{z \in Z} \mathbb{P}_{s_i, \bar{\beta}_{-i}}(z|x) v_i(\gamma(z)).$$

Since $\mu_{\bar{s}_i} = \delta_{\bar{\sigma}_{-i}}$,

$$\forall a_i \in A_i(h), V_i(a_i|h; \mu_{\bar{s}_i}, \phi_i) = \phi_i^{-1}(\phi_i(U_i((\bar{s}_i|h, a_i), \bar{\sigma}_{-i}|h))) = U_i((\bar{s}_i|h, a_i), \bar{\sigma}_{-i}|h).$$

Therefore \bar{s}_i is $(\mu_{\bar{s}_i}, \phi_i)$ -unimprovable. ■

10.2 Example 6

We prove that in Example 6, *there exists a concave and strictly increasing transformation $\phi_1 = \varphi_1 \circ \bar{\phi}_1$ such that for every belief μ_1 , Out is not prescribed by any (ϕ_1, μ_1) -unimprovable*

strategy of player 1. As anticipated in Example 5, we look for a ϕ_1 such that

$$\begin{aligned}\phi_1(u) &= \bar{\phi}_1(u) = \sqrt{u} \quad \text{if } u \in [0, 9], \\ \bar{\phi}_1(36) &= 6 > k = \phi_1(36) > \phi_1(9) = 3,\end{aligned}$$

with $k = \phi_1(36)$ close to 3.

For every σ_2 in $\Delta(A_2(\{\text{In}, \text{G}\})) = \Delta(\{\text{L}, \text{R}\})$, let

$$E(\sigma_2) = 1 \cdot \sigma_2(\text{L}) + 36 \cdot \sigma_2(\text{R}).$$

For every probability measure ν on $\Delta(\{\text{L}, \text{R}\})$, define $G_\nu : \Delta(\{\text{L}, \text{R}\}) \rightarrow [0, 1]$ as

$$G_\nu(\sigma_2) = \nu(\{\sigma'_2 \in \Delta(\{\text{L}, \text{R}\}) : E(\sigma'_2) \leq E(\sigma_2)\}).$$

For Player 1 to choose M over T, we must have

$$V_1(\text{M} | \{\{\text{In}, \text{G}\}\}; s_1, \mu_1, \phi_1) \geq 9 = V_1(\text{T} | \{\{\text{In}, \text{G}\}\}; s_1, \mu_1, \phi_1).$$

Let $\nu = \mu_1(\cdot | \{\{\text{In}, \text{G}\}\})$. Then, for every $\sigma_2 \in \Delta(\{\text{L}, \text{R}\})$,⁴¹

$$\phi_1^{-1}(G_\nu(\sigma_2) \cdot \phi_1(E(\sigma_2)) + (1 - G_\nu(\sigma_2)) \cdot \phi_1(36)) \geq V_1(\text{M} | \{\{\text{In}, \text{G}\}\}; s_1, \mu_1, \phi_1),$$

that is,

$$G_\nu(\sigma_2) \cdot \phi_1(E(\sigma_2)) + (1 - G_\nu(\sigma_2)) \cdot \phi_1(36) \geq \phi_1(V_1(\text{M} | \{\{\text{In}, \text{G}\}\}; s_1, \mu_1, \phi_1)).$$

Therefore, a necessary condition for 1 to choose M over T is that, for every $\sigma_2 \in \Delta(\{\text{L}, \text{R}\})$,

$$G_\nu(\sigma_2) \cdot \phi_1(E(\sigma_2)) + (1 - G_\nu(\sigma_2)) \cdot k \geq 3.$$

For $\hat{\sigma}_2$ with $E(\hat{\sigma}_2) = 17/2$, it reads

$$G_\nu(\hat{\sigma}_2) \leq \frac{k - 3}{k - \sqrt{17/2}}. \quad (11)$$

The interpretation is the following; to choose M at (In, G), the probability assigned to the models under which action M has an (objective) value lower or equal than 17/2 must be sufficiently small.

Now, we return to the choice of player 1 at the root of the game. Note that

$$\mu_1 = \mu_1(\cdot | \{\emptyset\}) = \mu_1(\cdot | \{\{\text{In}, \text{G}\}\})$$

because the belief about the co-player cannot change upon observing one's own move, or a chance move. The DM is sophisticated, therefore, in choosing between In and Out, he predicts his behavior at (In, G). There are three cases:

⁴¹Given the belief ν and a model σ_2 , $G_\nu(\sigma_2)$ is the probability assigned by ν to the models that, paired with action M, yield an (objective) expected utility lower or equal to the one obtained under model σ_2 . Therefore, the value of action M under ν cannot exceed the value of M under the following belief: the models that, paired with action M, yield an (objective) expected utility equal to the one obtained under model σ_2 have probability $G_\nu(\sigma_2)$, whereas δ_R has probability $(1 - G_\nu(\sigma_2))$.

1. He realizes that μ_1 is such that he will play T at (In, G). Then, since In.T is preferred to Out for every belief, he will play In at $\{\emptyset\}$.
2. He realizes that μ_1 is such that he will play M at (In, G). But then, it must be the case that (11) holds. Therefore, the probability assigned to the models that (given In.G) yield an objective expected utility of choosing action M larger than or equal to $\frac{17}{2}$ is at least $1 - \frac{k-3}{k-\sqrt{17/2}}$. In turn, this implies that the probability assigned to the models that, ex-ante, yields an objective expected utility of strategy In.M larger than or equal to $\frac{17}{4}$ is at least $1 - \frac{k-3}{k-\sqrt{17/2}}$. But then, the evaluation of strategy In.M under belief μ_1 satisfies the following conditions:⁴²

$$\begin{aligned} \phi_1(V_1(\text{In}|\{\emptyset\}; \text{In.M}, \mu_1, \phi_1)) &\geq \phi_1(1/2) \cdot \frac{k-3}{k-\sqrt{17/2}} + \phi_1(17/4) \cdot \left(1 - \frac{k-3}{k-\sqrt{17/2}}\right) \\ &= \frac{\sqrt{17}}{2} - \frac{k-3}{k-\sqrt{17/2}} \cdot \left(\frac{\sqrt{17}}{2} + 1\right). \end{aligned}$$

For $k = \phi_1$ (36) sufficiently close to 3, $\phi_1(V_1(\text{In}|\{\emptyset\}; \text{In.M}, \mu_1, \phi_1))$ is higher than 2, thus player 1 chooses In over Out.

3. He realizes that μ_1 is such that he will play B at (In, G). A similar argument as for the previous case shows that for k sufficiently close to 3 also in this case player 1 chooses In over Out.

Summing up, there is a concave and strictly increasing transformation $\phi_1 = \varphi_1 \circ \bar{\phi}_1$ such that Out is not prescribed by any (μ_1, ϕ_1) -unimprovable strategy for all beliefs μ_1 .

10.3 Monotonicity of the SCE correspondence

Define the set of messages for i consistent with an information set $h \in H_i$:

$$M_i(h) = \{m : \exists (x, z) \in h \times Z, (x \prec z) \wedge (m = f_i(z))\}.$$

Lemma 3 *Ex post perfect recall implies that*

$$S_{0,-i}(h) = \bigcup_{m \in M_i(h)} F_{i,s_i}^{-1}(m)$$

for all $i \in I$, $h \in H_i$, and $s_i \in S_i(h)$.

⁴²The first inequality is due to the following fact. The probability assigned by μ_1 to the models that, ex-ante, paired with strategy In.M, yield an (objective) expected utility larger or equal than $\frac{17}{4}$ is larger or equal than $1 - \frac{k-3}{k-\sqrt{17/2}}$. Therefore, the value of action In given strategy In.M under μ_1 cannot be lower than its value under the following belief: the models that, paired with strategy In.M, yield an (objective) expected utility equal to $\frac{17}{4}$ have probability $1 - \frac{k-3}{k-\sqrt{17/2}}$, whereas δ_L has probability $\frac{k-3}{k-\sqrt{17/2}}$.

In words, $S_{0,-i}(h)$ is the union of the sets of preimages of messages consistent with h , because these messages “record” that h has been reached.

Proof. First note that perfect recall implies

$$S_{0,I}(h) = S_i(h) \times S_{0,-i}(h).$$

We first prove that

$$S_{0,-i}(h) \subseteq \bigcup_{m \in M_i(h)} F_{i,s_i}^{-1}(m).$$

Fix any $s_{0,-i} \in S_{0,-i}(h)$; since $s_i \in S_i(h)$ and $S_{0,I}(h) = S_i(h) \times S_{0,-i}(h)$, then $(s_i, s_{0,-i}) \in S_{0,I}(h)$, that is, $x \prec \zeta(s_i, s_{0,-i})$ for some $x \in h$. Thus, by definition of $M_i(h)$, $f_i(\zeta(s_i, s_{0,-i})) \in M_i(h)$. Hence, $s_{0,-i} \in F_{i,s_i}^{-1}(m)$ for some $m \in M_i(h)$.

Next we prove by contraposition that the converse

$$S_{0,-i}(h) \supseteq \bigcup_{m \in M_i(h)} F_{i,s_i}^{-1}(m)$$

is implied by *ex post* perfect recall. Suppose that we can find some $m \in M_i(h)$ and $s'_{0,-i} \in F_{i,s_i}^{-1}(m) \setminus S_{0,-i}(h)$. We show that this implies a violation of *ex post* perfect recall. Since $m \in M_i(h)$ there is a pair $(x, z) \in h \times Z$ such that $x \prec z$ and $f_i(z) = m$. Fix any $s_{0,-i} \in \text{proj}_{S_{0,-i}} \zeta^{-1}(z)$, so that $(s_i, s_{0,-i}) \in S_i(h) \times S_{0,-i}(h) = S_{0,I}(h)$. Let $z = \zeta(s_i, s_{0,-i})$ and $z' = \zeta(s_i, s'_{0,-i})$. Then, by choice of $s_{0,-i}$ and $s'_{0,-i}$, $f_i(z) = m = f_i(z')$, z is preceded by a node of h and z' is not preceded by any node of h . Hence, there are $z, z' \in Z$ such that z has a predecessor in h , z' has no predecessor in h , and yet $f_i(z) = m = f_i(z')$, thus violating *ex post* perfect recall. \blacksquare

The following result says that the value of equilibrium actions is unambiguous, hence independent of ambiguity attitudes:

Lemma 4 *Let $\bar{\sigma}$ be an SCE of the game with observable payoffs (Γ, f, ϕ) justified by the confirmed beliefs $(\mu_{s_i})_{i \in I, s_i \in \text{Supp} \bar{\sigma}_i}$. For every $i \in I$ and $s_i \in \text{Supp} \bar{\sigma}_i$ and $h \in H_i(\mu_{s_i}) \cap H_i(s_i)$, action $s_{i,h}$ is $\mu_i(\cdot|h)$ -unambiguous, and its value is the conditional objective expected payoff, that is,*

$$V_i(s_{i,h}|h; s_i, \mu_{s_i}, \phi_i) = U_i(s_i, \bar{\sigma}_{-i}|h).$$

Proof. By *ex post* perfect recall and Lemma 3,

$$S_{0,-i}(h) = \bigcup_{m \in M_i(h)} F_{i,s_i}^{-1}(m) \tag{12}$$

for each $h \in H_i$. Also, recall that

$$\hat{\Sigma}_{-i}(s_i, \bar{\sigma}_{-i}) = \left\{ \sigma_{-i} \in \Sigma_{-i} : \forall m, (\sigma_0 \times \sigma_{-i})(F_{i,s_i}^{-1}(m)) = (\sigma_0 \times \bar{\sigma}_{-i})(F_{i,s_i}^{-1}(m)) \right\}$$

is the partially identified set of co-players distributions of strategies **observationally equivalent** for i to $\bar{\sigma}_{-i}$ given s_i .

Fix any $h \in H_i(\mu_{s_i}) \cap H_i(s_i)$. Then, for each $\sigma_{-i} \in \hat{\Sigma}_{-i}(s_i, \bar{\sigma}_{-i})$,

$$(\sigma_0 \times \sigma_{-i})(S_{0,-i}(h)) \stackrel{(\text{obs.eq.})}{=} (\sigma_0 \times \bar{\sigma}_{-i})(S_{0,-i}(h)) \stackrel{(\text{conf.})}{=} (\sigma_0 \times p_{\mu_{s_i}})(S_{0,-i}(h)) > 0, \quad (13)$$

where the first equality follows from eq. (12), the second equality follows from belief **confirmation** ($\mu_{s_i}(\hat{\Sigma}_{-i}(s_i, \bar{\sigma}_{-i})) = 1$) and eq. (12), and the inequality follows from $h \in H_i(\mu_{s_i})$.

Fix any $m \in M_i(h)$; by observable payoffs (obs.p), there is $v^m \in \mathbb{R}$ such that $v_i(\gamma(\zeta(s_0, s_i, s_{-i}))) = v^m$ for all $(s_0, s_{-i}) \in S_0 \times S_{-i}$ with $f_i(\zeta(s_0, s_i, s_{-i})) = m$. Then, observable payoffs and eq.s (12)-(13) imply that, for every $\sigma_{-i} \in \hat{\Sigma}_{-i}(s_i, \bar{\sigma}_{-i})$,

$$U_i(s_i, \sigma_{-i}|h) = U_i(s_i, \bar{\sigma}_{-i}|h). \quad (14)$$

Indeed, we have

$$\begin{aligned} U_i(s_i, \sigma_{-i}|h) &\stackrel{(\text{Def.})}{=} \sum_{(s_0, s_{-i}) \in S_{0,-i}(h)} \frac{\sigma_0(s_0) \cdot \sigma_{-i}(s_{-i})}{(\sigma_0 \times \sigma_{-i})(S_{0,-i}(h))} v_i(\gamma(\zeta(s_0, s_i, s_{-i}))) \\ &\stackrel{(12, \text{obs.p.})}{=} \sum_{m \in M_i(h)} \frac{(\sigma_0 \times \sigma_{-i})(F_{i,s_i}^{-1}(m)) v^m}{(\sigma_0 \times \sigma_{-i})(S_{0,-i}(h))} \stackrel{(\text{obs.eq.})}{=} \sum_{m \in M_i(h)} \frac{(\sigma_0 \times \bar{\sigma}_{-i})(F_{i,s_i}^{-1}(m)) v^m}{(\sigma_0 \times \bar{\sigma}_{-i})(S_{0,-i}(h))} \\ &\stackrel{(12, \text{obs.p.})}{=} \sum_{(s_0, s_{-i}) \in S_{0,-i}(h)} \frac{\sigma_0(s_0) \cdot \bar{\sigma}_{-i}(s_{-i})}{(\sigma_0 \times \bar{\sigma}_{-i})(S_{0,-i}(h))} v_i(\gamma(\zeta(s_0, s_i, s_{-i}))) \stackrel{(\text{Def.})}{=} U_i(s_i, \bar{\sigma}_{-i}|h). \end{aligned}$$

By confirmed beliefs, $\mu_{s_i}(\hat{\Sigma}_{-i}(s_i, \bar{\sigma}_{-i})) = 1$. Hence, $\mu_{s_i}(\hat{\Sigma}_{-i}(s_i, \bar{\sigma}_{-i})|h) = 1$ and

$$\begin{aligned} V_i(s_i, h|h; s_i, \mu_{s_i}, \phi_i) &= \phi_i^{-1} \left(\int_{\Sigma_{-i}(h)} \phi_i(U_i(s_i, \sigma_{-i}|h)) \mu_{s_i}(d\sigma_{-i}|h) \right) \\ &\stackrel{(14)}{=} \phi_i^{-1}(\phi_i(U_i(s_i, \bar{\sigma}_{-i}|h))) = U_i(s_i, \bar{\sigma}_{-i}|h). \end{aligned}$$

■

An increase in ambiguity aversion (weakly) decreases the values of actions:

Lemma 5 *If (Γ, f, ϕ) features more ambiguity aversion than $(\Gamma, f, \bar{\phi})$. Then,*

$$V_i(a_i|h; s_i, \mu_i, \bar{\phi}_i) \geq V_i(a_i|h; s_i, \mu_i, \phi_i)$$

for all $i \in I$, $\mu_i \in \Delta(\Sigma_{-i})$, $s_i \in S_i$, $h \in H_i(\mu_i)$, and $a_i \in A_i(h)$.

Proof. Fix i and s_i arbitrarily. For every $h \in H_i$ and $a_i \in A(h)$, define the following auxiliary function:

$$\begin{aligned} U_{h,a_i} : \Sigma_{-i}(h) &\rightarrow \mathbb{R}, \\ \sigma_{-i} &\mapsto U_i((s_i|h, a_i), \sigma_{-i}|h). \end{aligned}$$

With this, for every $\mu_i \in \Delta(\Sigma_{-i})$ and $h \in H_i(\mu_i)$, the conditional belief $\mu_i(\cdot|h)$ is well defined and

$$\begin{aligned} V_i(a_i|h; s_i, \mu_i, \bar{\phi}_i) &= \bar{\phi}_i^{-1} \left(\mathbb{E}_{\mu_i(\cdot|h)} [\bar{\phi}_i \circ U_{h,a_i}] \right), \\ V_i(a_i|h; s_i, \mu_i, \phi_i) &= \phi_i^{-1} \left(\mathbb{E}_{\mu_i(\cdot|h)} [\phi_i \circ U_{h,a_i}] \right). \end{aligned} \quad (15)$$

Since $\phi_i = \varphi \circ \bar{\phi}_i$ for some concave and strictly increasing $\varphi : \bar{\phi}_i(\mathbb{V}_i) \rightarrow \mathbb{R}$, Jensen's inequality implies

$$\varphi \left(\mathbb{E}_{\mu_i(\cdot|h)} [\bar{\phi}_i \circ U_{h,a_i}] \right) \geq \mathbb{E}_{\mu_i(\cdot|h)} [\varphi \circ \bar{\phi}_i \circ U_{h,a_i}] = \mathbb{E}_{\mu_i(\cdot|h)} [\phi_i \circ U_{h,a_i}].$$

By monotonicity of φ and $\bar{\phi}_i$, and recalling that $\phi_i^{-1} = \bar{\phi}_i^{-1} \circ \varphi^{-1}$

$$\mathbb{E}_{\mu_i(\cdot|h)} [\bar{\phi}_i \circ U_{h,a_i}] \geq \varphi^{-1} \left(\mathbb{E}_{\mu_i(\cdot|h)} [\phi_i \circ U_{h,a_i}] \right),$$

$$\bar{\phi}_i^{-1} \left(\mathbb{E}_{\mu_i(\cdot|h)} [\bar{\phi}_i \circ U_{h,a_i}] \right) \geq \left(\bar{\phi}_i^{-1} \circ \varphi^{-1} \right) \left(\mathbb{E}_{\mu_i(\cdot|h)} [\phi_i \circ U_{h,a_i}] \right) = \phi_i^{-1} \left(\mathbb{E}_{\mu_i(\cdot|h)} [\phi_i \circ U_{h,a_i}] \right).$$

By eq. (15), $V_i(a_i|h; s_i, \mu_i, \bar{\phi}_i) \geq V_i(a_i|h; s_i, \mu_i, \phi_i)$. ■

Lemma 6 *Fix two games with observable payoffs, (Γ, f, ϕ) and $(\Gamma, f, \bar{\phi})$, such that (Γ, f, ϕ) features more ambiguity aversion than $(\Gamma, f, \bar{\phi})$. Fix an SCE $\bar{\sigma}$ of $(\Gamma, f, \bar{\phi})$ and justifying beliefs $(\mu_{s_i})_{i \in I, s_i \in \text{Supp} \bar{\sigma}_i}$. Suppose that for each $i \in I$, every $s_i \in \text{Supp} \bar{\sigma}_i$ is $(\mu_{s_i}, \bar{\phi}_i)$ -sequentially optimal. Then, there exist maps $\bar{s}_i : \text{Supp} \bar{\sigma}_i \rightarrow S_i$, $i \in I$, such that*

- (i) *for every $i \in I$ and $h \in H_i(s_i)$, $\bar{s}_i(s_i)(h) = s_i(h)$, and*
- (ii) *$\sigma = (\bar{\sigma}_i \circ \bar{s}_i^{-1})_{i \in I}$ is an SCE of (Γ, f, ϕ) where for every $i \in I$ and $s_i \in \text{Supp} \bar{\sigma}_i$, μ_{s_i} justifies $\bar{s}_i(s_i)$.*

(So, σ and $\bar{\sigma}$ are realization equivalent and are justified by the same beliefs.)

Proof For every $i \in I$ and $s_i \in \text{Supp} \bar{\sigma}_i$, we construct a (μ_{s_i}, ϕ_i) -unimprovable strategy $\bar{s}_i(s_i)$. Map $\bar{s}_i : \text{Supp} \bar{\sigma}_i \rightarrow S_i$ is such that (1) $\bar{s}_i(s_i)_h = s_{i,h}$ for all $h \in H_i(s_i)$, and (2)

$\bar{s}_i(s_i)$ is derived by folding back on $H_i(\mu_{s_i}) \setminus H_i(s_i)$ given μ_{s_i} and ϕ_i .⁴³ Therefore, by construction, for every $h \in H_i(\mu_{s_i}) \setminus H_i(s_i)$,

$$\bar{s}_i(s_i)_h \in \arg \max_{a_i \in A_i(h)} V_i(a_i|h; \bar{s}_i(s_i), \mu_{s_i}, \phi_i).$$

Now, let $h \in H_i(\mu_{s_i}) \cap H_i(s_i)$. For every $a_i \in A_i(h)$, we have

$$\begin{aligned} V_i(\bar{s}_i(s_i)_h|h; \bar{s}_i(s_i), \mu_{s_i}, \phi_i) &= V_i(s_{i,h}|h; s_i, \mu_{s_i}, \phi_i) \\ &\stackrel{(L4)}{=} U_i(s_i, \bar{\sigma}_{-i}|h) \\ &\stackrel{(L4)}{=} V_i(s_{i,h}|h; s_i, \mu_{s_i}, \bar{\phi}_i) \\ &\stackrel{(s.opt.)}{\geq} V_i(a_i|h; \bar{s}_i(s_i), \mu_{s_i}, \bar{\phi}_i) \\ &\stackrel{(L5)}{\geq} V_i(a_i|h; \bar{s}_i(s_i), \mu_{s_i}, \phi_i), \end{aligned}$$

where the first equality follows from construction of $\bar{s}_i(s_i)$, the second and the third equalities from Lemma 4, the first inequality from sequential (s.opt.) optimality, and the second inequality from Lemma 5. This shows that $\bar{s}_i(s_i)$ is (μ_{s_i}, ϕ_i) -unimprovable.

To conclude, note that the profile $\left(\bar{\sigma}_i \circ \bar{s}_i^{-1}, (\mu_{\hat{s}_i})_{\hat{s}_i \in \text{Supp}(\bar{\sigma}_i \circ \bar{s}_i^{-1})} \right)_{i \in I}$, where $\mu_{\hat{s}_i}(\cdot) = \mu_{s_i}(\cdot)$ for some s_i with $\hat{s}_i = \bar{s}_i(s_i)$, also satisfies the confirmed beliefs condition (because $\left(\bar{s}_i(s_i), (\bar{\sigma}_j \circ \bar{s}_j^{-1})_{j \neq i} \right)$ and $(s_i, \bar{\sigma}_{-i})$, given σ_0 , induce the same distribution over terminal nodes). Therefore $(\bar{\sigma}_i \circ \bar{s}_i^{-1})_{i \in I}$ is an SCE of (Γ, f, ϕ) . \blacksquare

10.4 No player moves more than once

If player i does not move more than once on any path, then every information set of i is consistent with every strategy of i and the value of an action is independent of i 's strategy:

$$\forall s_i \in S_i, H(s_i) = H_i,$$

$$\forall (\mu_i, s'_i, s''_i) \in \Delta(\Sigma_{-i}) \times S_i^2, \forall h \in H_i(\mu_i), \forall a_i \in A_i(h), V_i(a_i|h; s'_i, \mu_i, \bar{\phi}_i) = V_i(a_i|h; s''_i, \mu_i, \bar{\phi}_i).$$

Proof of Corollary 4. Fix any $\bar{\sigma} \in SCE(\Gamma, f, \bar{\phi})$ with justifying confirmed beliefs $(\mu_{\bar{s}_i})_{i \in I, \bar{s}_i \in \text{Supp} \bar{\sigma}_i}$. Fix $i \in I$, $\bar{s}_i \in \text{Supp} \bar{\sigma}_i$, and $h \in H_i(\mu_{\bar{s}_i})$ arbitrarily. By $(\mu_{\bar{s}_i}, \bar{\phi}_i)$ -unimprovability of \bar{s}_i ,

$$\bar{s}_{i,h} \in \arg \max_{a_i \in A_i(h)} V_i(a_i|h; \bar{s}_i, \mu_{\bar{s}_i}, \bar{\phi}_i).$$

⁴³Note that if $h \notin H_i(s_i)$, and $h' \succ h$, $h' \notin H_i(s_i)$, therefore the folding back construction is well defined. See Section 4 for the description of folding back optimality.

Since i does not move again after h , this is equivalent to

$$\bar{s}_i \in \arg \max_{s_i \in S_i} V_i(s_{i,h}|h; s_i, \mu_{\bar{s}_i}, \bar{\phi}_i).$$

Then, by Lemma 6, there exist maps $\bar{s}_i : \text{Supp } \bar{\sigma}_i \rightarrow S_i$, $i \in I$, such that (i) for every $i \in I$ and $h \in H_i(\bar{s}_i)$, $\bar{s}_i(\bar{s}_i)(h) = \bar{s}_i(h)$, and (ii) $\sigma = (\bar{\sigma}_i \circ \bar{s}_i^{-1})_{i \in I}$ is an SCE of (Γ, f, ϕ) , justified by the same beliefs as $\bar{\sigma}$. Since $H_i(\bar{s}_i) = H_i$ when no player moves more than once on any path, $\sigma = \bar{\sigma}$. \blacksquare

10.5 Symmetric equilibria @[P must still check]@

In this section and the following ones, we consider games *without chance moves*. Therefore, the outcome function and the strategic-form feedback and payoff functions of player i are, respectively,

$$\zeta : S \rightarrow Z,$$

$$F_i = f_i \circ \zeta : S \rightarrow M,$$

and

$$U_i = v_i \circ \gamma \circ \zeta : S \rightarrow \mathbb{R}.$$

The proof of Theorem 5 requires some preliminary results. For every $i \in I$, $h \in H_i$, and $s_i \in S_i$, define the collections of information sets and terminal nodes “immediately” follow h within $H_i \cup Z$ given that i sticks to s_i from h onward, that is,

$$\begin{aligned} H_i(h, s_i) &= \{h' \in H_i((s_i|h)) : (h \prec h') \wedge (\nexists h'' \in H_i, h \prec h'' \prec h')\} \\ Z_i(h, s_i) &= \{z \in Z : \exists s_{-i} \in S_{-i}(h), (\zeta((s_i|h), s_{-i}) = z) \wedge (\nexists h'' \in H_i, h \prec h'' \prec z)\} \\ \Upsilon_i(h, s_i) &= H_i(h, s_i) \cup Z_i(h, s_i) \end{aligned}$$

(recall that $(s_i|h)$ is the replacement that coincides with s_i at all information sets of i that do not strictly precede h , $h'' \prec z$ means that $x \prec z$ for some $x \in h''$).

To shorten the expressions, in summations over $\Upsilon_i(h, s_i)$, $\mu_i(\cdot|x)$ will formally appear also for $x \in Z_i(h, s_i)$, @[P: I not like this $\mu_i(\cdot|x)$ thing, and I do not even see it below][E: it can be eliminated, it was used before when some expressions were more compact@ but it will disappear once the summation is broken down to separates summations over $H_i(h, s_i)$ and $Z_i(h, s_i)$. Moreover, let $S_{-i}(z) = \{s_{-i} \in S_{-i} : \exists s_i \in S_i, \zeta(s_i, s_{-i}) = z\}$.

Remark 5 Fix $i \in I$, $h \in H_i$ and $s_i \in S_i$. Then $\{S_{-i}(x)\}_{x \in \Upsilon_i(h, s_i)}$ is a partition of $S_{-i}(h)$.

We say that a belief $\mu_i \in \Delta(\Sigma_{-i})$ is supported by Dirac models if

$$\mu_i(\{\delta_{s_{-i}} : s_{-i} \in S_{-i}\}) = 1.$$

Since $V_i(a_i|h; s_i, \mu_i, \phi_i)$ is a certainty equivalent, we have

$$\phi_i(V_i(a_i|h; s_i, \mu_i, \phi_i)) = \int_{\Sigma_{-i}(h)} \phi_i(U_i((s_i|h, a_i), \sigma_{-i}|h)) \mu_i(d\sigma_{-i}).$$

Under a belief supported by Dirac models, the value of a strategy at an information set is a weighted sum of the values of the strategy at the next information sets, where the weights are the predictive probabilities of reaching them. **[P: Why the hats? The proof seems OK, but unnecessarily long. I do not see why we cannot make Corollary 14 just a remark. Why is it not obvious? Belief supported by Dirac models are isomorphic to predictive beliefs][E: agree, but we must leave the next lemma as intermediate step towards dynamic consistency because it is used directly later]@**

Lemma 7 Fix $i \in I$ and a belief supported by Dirac models $\hat{\mu}_i$. Then, for every $\hat{s}_i \in S_i$ and $h \in H_i(\hat{\mu}_i)$,

$$\begin{aligned} & \phi_i(V_i(\hat{s}_{i,h}|h; \hat{s}_i, \hat{\mu}_i, \phi_i)) \\ = & \sum_{h' \in H_i(h, \hat{s}_i)} p_{\hat{\mu}_i}(S_{-i}(h')|h) \phi_i(V_i((\hat{s}_i|h'), \hat{\mu}_i, \phi_i)) + \sum_{z \in Z_i(h, \hat{s}_i)} p_{\hat{\mu}_i}(S_{-i}(z)|h) \phi_i(v_i(\gamma(z))). \end{aligned}$$

Proof By Remark 5, we can write

$$\begin{aligned} & \phi_i(V_i(\hat{s}_{i,h}|h; \hat{s}_i, \hat{\mu}_i, \phi_i)) \\ = & \sum_{x \in \Upsilon_i(h, \hat{s}_i): p_{\hat{\mu}_i}(S_{-i}(x)|h) > 0} p_{\hat{\mu}_i}(S_{-i}(x)|h) \sum_{s_{-i} \in S_{-i}(x)} \phi_i(U_i((\hat{s}_i|h'), s_{-i})) \frac{\hat{\mu}_i(\delta_{s_{-i}}|h)}{p_{\hat{\mu}_i}(S_{-i}(x)|h)}. \end{aligned}$$

For every $h' \in H_i(h, \hat{s}_i)$ with $p_{\hat{\mu}_i}(S_{-i}(h')|h) > 0$ and for every $s_{-i} \in S_{-i}(h')$,

$$\frac{\hat{\mu}_i(\delta_{s_{-i}}|h)}{p_{\hat{\mu}_i}(S_{-i}(h')|h)} = \hat{\mu}_i(\delta_{s_{-i}}|h').$$

For every $z \in Z_i(h, \hat{s}_i)$ with $p_{\hat{\mu}_i}(S_{-i}(z)|h) > 0$ and for every $s_{-i} \in S_{-i}(z)$, it holds $U_i((\hat{s}_i|h), s_{-i}) = v_i(\gamma(z))$, and

$$\sum_{s'_{-i} \in S_{-i}(z)} \frac{\hat{\mu}_i(\delta_{s'_{-i}}|h)}{p_{\hat{\mu}_i}(S_{-i}(z)|h)} = 1.$$

So we have the desired result. ■

Iterating and using the chain rule, one can prove the following:

Corollary 14 Fix $i \in I$ and a belief supported by Dirac models $\hat{\mu}_i$. Then, for every $\hat{s}_i \in S_i$ and $h \in H_i(\hat{\mu}_i)$,

$$\phi_i(V_i(\hat{s}_{i,h}|h; \hat{s}_i, \hat{\mu}_i, \phi_i)) = \sum_{s_{-i} \in S_{-i}} p_{\hat{\mu}_i}(s_{-i}|h) \phi_i(v_i(\gamma(\zeta((\hat{s}_i|h), s_{-i}))))).$$

@[P: I had to stop here at 7:45pm ITA]@

Thus, the value of a strategy under a belief supported by Dirac models corresponds to its expected payoff under the corresponding predictive belief after transforming the payoffs of the game with ϕ_i . Then, by a well known dynamic programming argument, we can claim the following.

Lemma 8 Fix $i \in I$ and a belief supported by Dirac models $\hat{\mu}_i$. Then \hat{s}_i is a $(\hat{\mu}_i, \phi_i)$ -unimprovable strategy if and only if it is $(\hat{\mu}_i, \phi_i)$ -sequentially optimal.

Consider now a generic belief μ_i that yields the same predictive belief $p_{\mu_i} = p_{\hat{\mu}_i}$ as $\hat{\mu}_i$. For an ambiguity averse agent, the value at h of a (μ_i, ϕ_i) -unimprovable strategy is not lower than the value at h of a $(\hat{\mu}_i, \phi_i)$ -unimprovable strategy.

Lemma 9 Fix $i \in I$, $\mu_i \in \Delta(\Sigma_{-i})$, and concave ϕ_i . Call $\hat{\mu}_i \in \Delta(\Sigma_{-i})$ the belief supported by Dirac models such that for every $h \in H_i(\mu_i)$ and $s_{-i} \in S_{-i}$, $\hat{\mu}_i(\delta_{s_{-i}}|h) = p_{\mu_i}(s_{-i}|h)$. Let s_i^* and \hat{s}_i^* be (μ_i, ϕ_i) - and $(\hat{\mu}_i, \phi_i)$ -unimprovable strategies. Then, for every $h \in H_i(\mu_i) = H_i(\hat{\mu}_i)$,

$$V_i(s_{i,h}^*|h; s_i^*, \mu_i, \phi_i) \geq V_i(\hat{s}_{i,h}^*|h; \hat{s}_i^*, \hat{\mu}_i, \phi_i).$$

Proof Fix $h \in H_i(\mu_i)$. Suppose, by way of induction, that for every $h' \in H_i(\mu_i)$ with $h \prec h'$,

$$V_i(s_{i,h'}^*|h'; s_i^*, \mu_i, \phi_i) \geq V_i(\hat{s}_{i,h'}^*|h'; \hat{s}_i^*, \hat{\mu}_i, \phi_i).$$

If there is no $h' \in H_i(\mu_i)$ with $h' \succ h$ (basis step), the proof will not require an inductive hypothesis.

Define the replacement plan $s_i^h = (s_i^*|h, \hat{s}_{i,h}^*)$. Note that s_i^h prescribes $\hat{s}_{i,h}^*$ instead of $s_{i,h}^*$ at h . By unimprovability of s_i^* at h we have

$$V_i(s_{i,h}^*|h; s_i^*, \mu_i, \phi_i) \geq V_i(\hat{s}_{i,h}^*|h; s_i^*, \mu_i, \phi_i).$$

By Remark 5 we can write

$$\begin{aligned} & \phi_i(V_i(\hat{s}_{i,h}^*|h; s_i^*, \mu_i, \phi_i)) \\ &= \int_{\Sigma_{-i}(h)} \phi_i \left(\sum_{x \in \Upsilon_i(h, s_i^h)} \sum_{s_{-i} \in S_{-i}(x)} \sigma_{-i}(s_{-i}|h) U_i(s_i^h, s_{-i}) \right) \cdot \mu_i(d\sigma_{-i}|h). \end{aligned}$$

By concavity of ϕ_i and Jensen's inequality, we have

$$\begin{aligned} & \int_{\Sigma_{-i}(h)} \phi_i \left(\sum_{x \in \Upsilon_i(h, s_i^h)} \sum_{s_{-i} \in S_{-i}(x) \cap \text{Supp} \sigma_{-i}} \frac{\sigma_{-i}(S_{-i}(x)|h)}{\sigma_{-i}(S_{-i}(x)|h)} \sigma_{-i}(s_{-i}|h) U_i(s_i^h, s_{-i}) \right) \cdot \mu_i(d\sigma_{-i}|h) \\ & \geq \int_{\Sigma_{-i}(h)} \sum_{x \in \Upsilon_i(h, s_i^h)} \sigma_{-i}(S_{-i}(x)|h) \cdot \phi_i \left(\sum_{s_{-i} \in S_{-i}(x) \cap \text{Supp} \sigma_{-i}} \frac{\sigma_{-i}(s_{-i}|h)}{\sigma_{-i}(S_{-i}(x)|h)} U_i(s_i^h, s_{-i}) \right) \cdot \mu_i(d\sigma_{-i}|h) \end{aligned}$$

The last expression can be rewritten as:

$$\begin{aligned} & \sum_{\substack{x \in \Upsilon_i(h, s_i^h), \\ \mu_i(\Sigma_{-i}(x)|h) \neq 0}} \int_{\Sigma_{-i}(x)} \sigma_{-i}(S_{-i}(x)|h) \cdot \phi_i \left(\sum_{s_{-i} \in S_{-i}(x)} \frac{\sigma_{-i}(s_{-i}|h) U_i(s_i^h, s_{-i})}{\sigma_{-i}(S_{-i}(x)|h)} \right) \cdot \frac{p_{\mu_i}(S_{-i}(x)|h)}{p_{\mu_i}(S_{-i}(x)|h)} \mu_i(d\sigma_{-i}|h) = \\ & \sum_{\substack{x \in \Upsilon_i(h, s_i^h), \\ \mu_i(\Sigma_{-i}(x)|h) \neq 0}} p_{\mu_i}(S_{-i}(x)|h) \int_{\Sigma_{-i}(x)} \phi_i \left(\sum_{s_{-i} \in S_{-i}(x)} \frac{\sigma_{-i}(s_{-i}|h) U_i(s_i^h, s_{-i})}{\sigma_{-i}(S_{-i}(x)|h)} \right) \cdot \frac{\sigma_{-i}(S_{-i}(x)|h) \mu_i(d\sigma_{-i}|h)}{p_{\mu_i}(S_{-i}(x)|h)}. \end{aligned}$$

For every $h' \in H_i(h, s_i^h)$ with $\mu_i(\Sigma_{-i}(h')|h) \neq 0$, $\sigma_{-i} \in \Sigma_{-i}(h')$, $f \in \mathbb{R}^{\Sigma_{-i}(h')}$ measurable and $s_{-i} \in S_{-i}(h')$,

$$\begin{aligned} & \frac{\sigma_{-i}(s_{-i}|h)}{\sigma_{-i}(S_{-i}(h')|h)} = \sigma_{-i}(s_{-i}|h'); \\ & \int_{\Sigma_{-i}(h')} f(\sigma'_{-i}) \frac{\sigma'_{-i}(S_{-i}(h')|h)}{p_{\mu_i}(S_{-i}(h')|h)} \mu_i(d\sigma'_{-i}|h) = \int_{\Sigma_{-i}(h')} f(\sigma'_{-i}) \mu_i(d\sigma'_{-i}|h). \end{aligned}$$

For every $z \in Z_i(h, s_i)$ with $\mu_i(\Sigma_{-i}(z)|h) \neq 0$, $\sigma_{-i} \in \Sigma_{-i}(z)$ and $s_{-i} \in S_{-i}(z) \cap \text{Supp} \sigma_{-i}$, it holds $U_i(s_i^h, s_{-i}) = v_i(\gamma(z))$. Moreover,

$$\begin{aligned} & \sum_{s'_{-i} \in S_{-i}(z)} \frac{\sigma_{-i}(s'_{-i}|h)}{\sigma_{-i}(S_{-i}(z)|h)} = 1; \\ & \int_{\Sigma_{-i}(z)} \frac{\sigma_{-i}(S_{-i}(z)|h) \mu_i(d\sigma_{-i}|h)}{p_{\mu_i}(S_{-i}(z)|h)} = 1. \end{aligned}$$

So we have

$$\begin{aligned} & \phi_i(V_i(s_{i,h}^*|h; s_i^*, \mu_i, \phi_i)) \geq \\ & \sum_{h' \in H_i(h, s_i^h)} p_{\mu_i}(S_{-i}(h')|h) \phi_i(V_i(s_{i,h'}^*|h'; s_i^*, \mu_i, \phi_i)) + \sum_{z \in Z_i(h, s_i^h)} p_{\mu_i}(S_{-i}(z)|h) \phi_i(v_i(\gamma(z))). \end{aligned}$$

By Lemma 7,

$$\begin{aligned} & \phi_i (V_i (\hat{s}_{i,h}^* | h; \hat{s}_i^*, \hat{\mu}_i, \phi_i)) = \\ & \sum_{h' \in H_i(h, \hat{s}_i^*)} p_{\hat{\mu}_i} (S_{-i}(h') | h) \phi_i (V_i (\hat{s}_{i,h'}^* | h'; \hat{s}_i^*, \hat{\mu}_i, \phi_i)) + \sum_{z \in Z_i(h, \hat{s}_i^*)} p_{\hat{\mu}_i} (S_{-i}(z) | h) \phi_i (v_i (\gamma (z))). \end{aligned}$$

Note that by $s_{i,h}^h = \hat{s}_{i,h}^*$, we have $H_i(h, s_i^h) = H_i(h, \hat{s}_i^*)$ and $Z_i(h, s_i^h) = Z_i(h, \hat{s}_i^*)$. Moreover, recall that $p_{\hat{\mu}_i} = p_{\mu_i}$. Then, by $H_i(h, s_i^h) = \emptyset$ for the basis step and by the inductive hypothesis for the inductive step:

$$V_i (s_{i,h}^* | h; s_i^*, \mu_i, \phi_i) \geq V_i (\hat{s}_{i,h}^* | h; \hat{s}_i^*, \hat{\mu}_i, \phi_i).$$

■

Proof of Theorem 5. Let $z \in \zeta (\text{symSCE} (\Gamma, f, \bar{\phi}))$. Then, there is a symmetric SCE $(s_i^*)_{i \in I}$ of $(\Gamma, f, \bar{\phi})$ such that $z = \zeta (s^*)$. Fix $i \in I$, let μ_i be the confirmed belief that justifies s_i^* , and let $m_i^* = f_i (z)$. Call $\hat{\mu}_i \in \Delta (\Sigma_{-i})$ the belief supported by Dirac models such that for every $s_{-i} \in S_{-i}$, $\hat{\mu}_i (\delta_{s_{-i}}) = p_{\mu_i} (s_{-i})$. Fix any $(\hat{\mu}_i, \bar{\phi}_i)$ -unimprovable strategy \hat{s}_i^* and define \hat{s}_i as

$$\hat{s}_{i,h} = \begin{cases} s_{i,h}^* & h \in H_i, f^{-1} (m^*) \succ h, \\ \hat{s}_{i,h}^* & h \in H_i, f^{-1} (m^*) \not\succeq h. \end{cases}$$

Clearly, for every $h \in H_i$ with $f^{-1} (m^*) \not\succeq h$, we have

$$V_i (\hat{s}_{i,h} | h; \hat{s}_i, \hat{\mu}_i, \bar{\phi}_i) = V_i (\hat{s}_{i,h}^* | h; \hat{s}_i^*, \hat{\mu}_i, \bar{\phi}_i).$$

Fix $h \in H_i$ with $f^{-1} (m^*) \succ h$. By confirmed beliefs $p_{\mu_i} (S_{-i} (f^{-1} (m^*)) | h) = 1$. So, by observable payoffs:

$$V_i (\hat{s}_{i,h} | h; \hat{s}_i, \hat{\mu}_i, \bar{\phi}_i) = v_i (\gamma (z)) = V_i (s_{i,h}^* | h; s_i^*, \mu_i, \bar{\phi}_i).$$

By Lemma 9,

$$V_i (s_{i,h}^* | h; s_i^*, \mu_i, \bar{\phi}_i) \geq V_i (\hat{s}_{i,h}^* | h; \hat{s}_i^*, \hat{\mu}_i, \bar{\phi}_i).$$

So we have:

$$V_i (\hat{s}_{i,h} | h; \hat{s}_i, \hat{\mu}_i, \bar{\phi}_i) \geq V_i (\hat{s}_{i,h}^* | h; \hat{s}_i^*, \hat{\mu}_i, \bar{\phi}_i).$$

Then, by Lemma 8,

$$\hat{s}_{i,h} \in \arg \max_{a_i \in A_i(h)} V_i (a_i | h; \hat{s}_i, \hat{\mu}_i, \bar{\phi}_i).$$

Repeating for every player, \hat{s} is a symmetric SCE for $(\Gamma, f, \bar{\phi})$ justified by the confirmed beliefs $(\hat{\mu}_i)_{i \in I}$. By Lemma 8, each \hat{s}_i is $(\hat{\mu}_i, \bar{\phi}_i)$ -sequentially optimal. Since $\zeta (\hat{s}) = z$, by Lemma 6, there exists a symmetric SCE of (Γ, f, ϕ) , $s = (s_i)_{i \in I}$, such that $\zeta (s) = z$. ■

10.6 Full Unimprovability and Rationalizable SCE

For the reader's convenience we recall some notation and definitions of Section 7. $\bar{H}_i = H_i \cup \{\{\emptyset\}\}$ is the extended collection of information sets that includes the initial information set $\{\emptyset\}$ even if i is not a first mover ($\bar{H}_i = H_i$ if i is a first mover). With this, we let $\bar{H}_i(\mu_i) = H_i(\mu_i) \cup \{\{\emptyset\}\}$ be the subcollection of possible information sets given prior μ_i . A CPS $\bar{\mu}_i(\cdot|\cdot) \in \Delta^{\bar{H}_i}(\Sigma_{-i}) \subseteq [\Delta(\Sigma_{-i})]^{\bar{H}_i}$ for player i specifies an initial belief $\bar{\mu}_i(\cdot|\{\emptyset\}) \in \Delta(\Sigma_{-i})$ and a conditional belief $\bar{\mu}_i(\cdot|h)$ for each $h \in H_i$. Full $(\bar{\mu}_i(\cdot|\cdot), \phi_i)$ -unimprovability requires value-maximization over actions at each $h \in H_i$, that is, at the information sets where i is active. The initial belief $\bar{\mu}_i(\cdot|\{\emptyset\})$ matters to compare full $(\bar{\mu}_i(\cdot|\cdot), \phi_i)$ -unimprovability with (μ_i, ϕ_i) -unimprovability, because in the latter μ_i has to be interpreted as an initial belief. Furthermore $\bar{\mu}_i(\cdot|\{\emptyset\})$ captures how i strategically analyzes the game before playing it. Finally, recall that here we assume that there are *no chance moves*.

10.6.1 Equivalence of SCE and fully unimprovable SCE

Proof of Proposition 9. First note that, for every prior μ_i on Σ_{-i} , one can find a CPS $\bar{\mu}_i(\cdot|\cdot)$ on (Σ_{-i}, \bar{H}_i) such that $\mu_i(\cdot|h) = \bar{\mu}_i(\cdot|h)$ for all information sets $h \in \bar{H}_i(\mu_i)$: Let $\bar{\mu}_i(\cdot|\{\emptyset\}) = \mu_i$ and derive $\bar{\mu}_i(\cdot|h)$ by conditioning for all $h \in H_i(\mu_i)$. Next, for every $h \in H_i \setminus H_i(\mu_i)$ whose immediate predecessor h in (\bar{H}_i, \preceq) belongs to $\bar{H}_i(\mu_i)$, fix some $\bar{\mu}_i(\cdot|h) \in \Delta(\Sigma_{-i}(h))$ such that $p_{\bar{\mu}_i(\cdot|h)}(S_{-i}(h')) > 0$ for all the information sets h' that weakly follow h (e.g., $\bar{\mu}_i(\cdot|h) = \frac{1}{|S_{-i}(h)|} \sum_{s_{-i} \in S_{-i}(h)} \delta_{\delta_{s_{-i}}}$) and derive $\bar{\mu}_i(\cdot|h')$ from $\bar{\mu}_i(\cdot|h)$ by conditioning. One can check that the constructed array $(\bar{\mu}_i(\cdot|h))_{h \in \bar{H}_i} \in [\Delta(\Sigma_{-i})]^{\bar{H}_i}$ is a CPS.

Fix an SCE σ justified by confirmed beliefs $(\mu_{s_i})_{i \in I, s_i \in \text{Supp}\sigma_i}$. For each $i \in I$ and $s_i \in \text{Supp}\sigma_i$, let $\bar{\mu}_{s_i}(\cdot|\cdot) \in \Delta^{\bar{H}_i}(\Sigma_{-i})$ be a CPS such that $\bar{\mu}_{s_i}(\cdot|h) = \mu_{s_i}(\cdot|h)$ for every $h \in H_i(\mu_{s_i})$ (see above). Now construct a new strategy $\hat{s}_i(s_i)$ so that (1) $\hat{s}_i(s_i)(h) = s_i(h)$ for all $h \in H_i(\mu_{s_i})$, and (2) $\hat{s}_i(s_i)$ is derived by folding back on $H_i \setminus H_i(\mu_{s_i})$ given $\bar{\mu}_{s_i}(\cdot|\cdot)$. Since s_i is (μ_{s_i}, ϕ_i) -unimprovable, $\hat{s}_i(s_i)$ must be *fully* $(\bar{\mu}_{s_i}(\cdot|\cdot), \phi_i)$ -unimprovable. By construction, (s_i, μ_{s_i}) and $(\hat{s}_i(s_i), \bar{\mu}_{s_i}(\cdot|\{\emptyset\}))$ imply the same probabilities of terminal nodes, because, they reach the same information sets, $H_i(s_i) \cap H_i(\mu_{s_i}) = H_i(\hat{s}_i(s_i)) \cap H_i(\bar{\mu}_{s_i}(\cdot|\{\emptyset\}))$, where s_i and $\hat{s}_i(s_i)$ prescribe the same actions.

By ex post perfect recall and the self-confirming conditions for σ , for every $i \in I$ and $s_i \in \text{Supp}\sigma_i$, (s_i, μ_{s_i}) and (s_i, σ_{-i}) yield the same probabilities of reaching information sets of i : $p_{\mu_{s_i}}(S_{-i}(h)) = \sigma_{-i}(S_{-i}(h))$ for every $h \in H_i(s_i)$. Therefore, by construction, $\times p_{\bar{\mu}_{s_i}}(S_{-i}(h)|\{\emptyset\}) = \sigma_{-i}(S_{-i}(h))$ for every $i \in I$, $s_i \in \text{Supp}\sigma_i$, and $h \in H_i(s_i)$. Now, for every $i \in I$, consider the pushforward measure $\hat{\sigma}_i = \sigma_i \circ \hat{s}_i^{-1}$, that is,

$$\forall s_i \in S_i, \hat{\sigma}_i(s_i) = \sum_{s'_i: \hat{s}_i(s'_i) = s_i} \sigma_i(s'_i).$$

By construction σ and $\hat{\sigma}$ yield the same distribution over terminal nodes, because they reach the same information sets and, for each $i \in I$, the pure strategies in the support of σ_i take the same actions as the associated pure strategies in the support of $\hat{\sigma}_i$ at all reachable information sets. Furthermore, the profile $\left(\hat{\sigma}_i, \left(\bar{\mu}_{\bar{s}_i}(\cdot|\cdot)\right)_{\bar{s}_i \in \text{Supp}\hat{\sigma}_i}\right)_{i \in I}$, where $\bar{\mu}_{\bar{s}_i}(\cdot|\cdot) = \bar{\mu}_{s_i}(\cdot|\cdot)$ for some $s_i \in \hat{s}_i^{-1}(s_i)$,⁴⁴ also satisfies the confirmed beliefs condition on top of the full unimprovability condition. Therefore $\hat{\sigma}$ is a fully unimprovable SCE. \blacksquare

10.6.2 Monotonicity of the symmetric RSCE correspondence

Recall that in the analysis of symmetric RSCE, we assume that there are no chance moves.

Proof of Remark 4 (Only if) Let $\bar{B}_i = \bigcap_{k \in \mathbb{N}} B_i^k$ for each $i \in I$; it can be checked that

$(\bar{B}_i)_{i \in I}$ satisfies the required property.

(If) We show by induction that if $(\bar{B}_i)_{i \in I} \in \times_{i \in I} 2^{S_i \times M}$ has the required property, then $\bar{B}_i \subseteq B_i^k$ for every $i \in I$ and $k \in \mathbb{N}_0$. The claim is trivially true for $k = 0$. Suppose it is true for $k \in \mathbb{N}_0$; fix $i \in I$ and $(\hat{s}_i, \hat{m}_i) \in \bar{B}_i$ arbitrarily. Then $(\hat{s}_i, \hat{m}_i) \in B_i^k$ (inductive hypothesis), and there is $p_i(\cdot|\cdot) \in \Delta^{\bar{H}_i}(S_{-i})$ with $\hat{s}_i \in r_i(p_i(\cdot|\cdot), \phi_i)$ such that eq. (10) holds. By the inductive hypothesis, $\bar{B}_{-i} \subseteq B_{-i}^k$; hence

$$\begin{aligned} & p_i \left(F_{\hat{s}_i}^{-1}(\hat{m}_i) \cap \left\{ s_{-i} : (s_j, F_j(\hat{s}_i, s_{-i}))_{j \neq i} \in B_{-i}^k \right\} \mid \{\emptyset\} \right) \\ & \geq p_i \left(F_{\hat{s}_i}^{-1}(\hat{m}_i) \cap \left\{ s_{-i} : (s_j, F_j(\hat{s}_i, s_{-i}))_{j \neq i} \in \bar{B}_{-i} \right\} \mid \{\emptyset\} \right), \end{aligned}$$

and (10) implies

$$p_i \left(F_{\hat{s}_i}^{-1}(\hat{m}_i) \cap \left\{ s_{-i} : (s_j, F_j(\hat{s}_i, s_{-i}))_{j \neq i} \in B_{-i}^k \right\} \mid \{\emptyset\} \right) = 1.$$

Therefore, $(\hat{s}_i, \hat{m}_i) \in B_i^{k+1}$. \blacksquare

Proof of Lemma 2 The statement is trivially true for $k = 0$. Suppose, by way of induction, that

$$\text{symSCE}^k(\Gamma, f, \phi) = \left\{ \bar{s} : (\bar{s}_i, F_i(\bar{s}))_{i \in I} \in B^k \right\}.$$

We first show that, for every fixed \bar{s} in the above set and $i \in I$

$$F_{\bar{s}_i}^{-1}(F_i(\bar{s})) \cap \text{symSCE}^k(\Gamma, f, \phi)_{\bar{s}_i} = F_{\bar{s}_i}^{-1}(F_i(\bar{s})) \cap \left\{ s_{-i} : (s_j, F_j(\bar{s}_i, s_{-i}))_{j \in I \setminus \{i\}} \in B_{-i}^k \right\}. \quad (16)$$

⁴⁴It can be checked that it does not matter which $s_i \in \hat{s}_i^{-1}(s_i)$ we pick.

(*Proof of \subseteq*) By definition of section and the inductive hypothesis,

$$\begin{aligned} \text{symSCE}^k(\Gamma, f, \phi)_{\bar{s}_i} &= \left\{ s_{-i} : (\bar{s}_i, s_{-i}) \in \text{symSCE}^k(\Gamma, f, \phi) \right\} \\ &= \left\{ s_{-i} : ((\bar{s}_i, F_i(\bar{s}_i, s_{-i})), (s_j, F_j(\bar{s}_i, s_{-i})))_{j \in I \setminus \{i\}} \in B^k \right\} \\ &\subseteq \left\{ s_{-i} : (s_j, F_j(\bar{s}_i, s_{-i}))_{j \neq i} \in B_{-i}^k \right\}. \end{aligned}$$

Hence

$$F_{\bar{s}_i}^{-1}(F_i(\bar{s})) \cap \text{symSCE}^k(\Gamma, f, \phi)_{\bar{s}_i} \subseteq F_{\bar{s}_i}^{-1}(F_i(\bar{s})) \cap \left\{ s_{-i} : (s_j, F_j(\bar{s}_i, s_{-i}))_{j \in I \setminus \{i\}} \in B_{-i}^k \right\}.$$

(*Proof of \supseteq*) Let $\hat{s}_{-i} \in F_{\bar{s}_i}^{-1}(F_i(\bar{s})) \cap \left\{ s_{-i} : (s_j, F_j(\bar{s}_i, s_{-i}))_{j \in I \setminus \{i\}} \in B_{-i}^k \right\}$. Since $\hat{s}_{-i} \in F_{\bar{s}_i}^{-1}(F_i(\bar{s}))$ and $(\bar{s}_i, F_i(\bar{s})) \in B_i^k$, then $F_i(\bar{s}_i, \hat{s}_{-i}) = F_i(\bar{s})$ and

$$((\bar{s}_i, F_i(\bar{s}_i, \hat{s}_{-i})), (\hat{s}_j, F_j(\bar{s}_i, \hat{s}_{-i})))_{j \in I \setminus \{i\}} \in B^k.$$

Hence, $\hat{s}_{-i} \in \text{symSCE}^k(\Gamma, f, \phi)_{\bar{s}_i}$. Thus

$$F_{\bar{s}_i}^{-1}(F_i(\bar{s})) \cap \left\{ s_{-i} : (s_j, F_j(\bar{s}_i, s_{-i}))_{j \in I \setminus \{i\}} \in B_{-i}^k \right\} \subseteq F_{\bar{s}_i}^{-1}(F_i(\bar{s})) \cap \text{symSCE}^k(\Gamma, f, \phi)_{\bar{s}_i}.$$

This completes the proof of (16). \square

Now, let $\bar{s} \in \text{symSCE}^{k+1}(\Gamma, f, \phi)$; then – by definition – $\bar{s} \in \text{symSCE}^k(\Gamma, f, \phi)$ and the inductive hypothesis implies that $((\bar{s}_i, F_i(\bar{s})))_{i \in I} \in B^k$. Let $(p_i(\cdot))_{i \in I}$ be as in the definition of $\text{symSCE}^{k+1}(\Gamma, f, \phi)$. Then, by eq. (16),

$$p_i \left(F_{\bar{s}_i}^{-1}(F_i(\bar{s})) \cap \left\{ s_{-i} : (s_j, F_j(\bar{s}_i, s_{-i}))_{j \in I \setminus \{i\}} \in B_{-i}^k \right\} \mid \{\emptyset\} \right) = 1$$

for every $i \in I$, and $(\bar{s}_i, F_i(\bar{s}))_{i \in I} \in B^{k+1}$. Similarly, let $(\bar{s}_i, F_i(\bar{s}))_{i \in I} \in B^{k+1}$; then, by eq. (16),

$$p_i \left(F_{\bar{s}_i}^{-1}(F_i(\bar{s})) \cap \text{symSCE}^k(\Gamma, f, \phi)_{\bar{s}_i} \mid \{\emptyset\} \right) = 1$$

for every $i \in I$, and $\bar{s} \in \text{symSCE}^{k+1}(\Gamma, f, \phi)$. \blacksquare

Our results about monotonicity of the RSCE correspondence rely on the following result of *monotonicity of the justifiability correspondence*: Consider a decision problem under uncertainty $(\hat{A}, \hat{S}, \hat{u})$ where \hat{A} and \hat{S} are finite sets of actions and states, and $\hat{u} : \hat{A} \times \hat{S} \rightarrow \mathbb{R}$ is a vNM utility function. Fix a non-empty and *compact* set of distributions $\hat{\Sigma} \subseteq \Delta(\hat{S})$. Let $\mathcal{IC} \subseteq \mathbb{R}^{\hat{\Sigma}}$ denote the set of *strictly increasing* and *continuous* functions. For each $\phi \in \mathcal{IC}$ and $\mu \in \Delta(\hat{\Sigma})$ define the set of (μ, ϕ) -best replies:

$$\hat{r}(\mu, \phi) = \arg \max_{a \in \hat{A}} \int_{\hat{\Sigma}} \phi(\hat{u}(a, \sigma)) \mu(d\sigma).$$

Similarly, for each $p \in \Delta(\hat{S})$ we write

$$\hat{r}(p, \phi) = \arg \max_{a \in A} \sum_{s \in \hat{S}} \phi(\hat{u}(a, s)) p(s),$$

which is the special case when μ is supported by Dirac beliefs and p is the corresponding predictive belief. If an action is a (μ, ϕ) -best reply we say that it is ϕ -**justified** by μ ; we say that it is ϕ -**justifiable** if it is ϕ -justified by some μ . Battigalli et al. (2016a) proved that the ϕ -justifiability correspondence is monotone with respect to concave and strictly increasing transformations (see also Weinstein, 2016):

Lemma 10 *For all $\bar{\phi}, \tau \in \mathcal{IC}$, if $\tau \in \mathcal{IC}$ is a concave transformation then*

$$\bigcup_{\mu \in \Delta(\hat{S})} \hat{r}(\mu, \tau \circ \bar{\phi}) \supseteq \bigcup_{\bar{\mu} \in \Delta(\hat{S})} \hat{r}(\bar{\mu}, \bar{\phi}).$$

Since $\hat{S} \cong \{\delta_s : s \in \hat{S}\} \subseteq \Delta(\hat{S})$, letting $\hat{\Sigma} = \{\delta_s : s \in \hat{S}\}$ in Lemma 10 we obtain the following:

Corollary 15 *For all $\bar{\phi}, \tau \in \mathcal{IC}$, if $\tau \in \mathcal{IC}$ is a concave transformation then*

$$\bigcup_{p \in \Delta(\hat{S})} \hat{r}(p, \tau \circ \bar{\phi}) \supseteq \bigcup_{\bar{p} \in \Delta(\hat{S})} \hat{r}(\bar{p}, \bar{\phi}).$$

From now on, whenever we refer to games in which every player moves at most once along every path, strategies will be omitted from the value formulas. In this class of games, the value of an action at an information set depends only on the action itself and not on the overall strategy of the agent.

Proof of Theorem 11 To prove the result, we will show that

$$\forall k \in \mathbb{N}, \text{symSCE}^k(\Gamma, f, \bar{\phi}) \subseteq \text{symSCE}^k(\Gamma, f, \phi).$$

Then the claim follows from Lemma 2. The statement is trivially true for $k = 0$. Suppose, by way of induction, that

$$\text{symSCE}^k(\Gamma, f, \bar{\phi}) \subseteq \text{symSCE}^k(\Gamma, f, \phi). \quad (\text{I.H.})$$

Fix $\bar{s} \in \text{symSCE}^{k+1}(\Gamma, f, \bar{\phi})$ and $i \in I$ arbitrarily. By definition of symSCE^{k+1} and the inductive hypothesis (I.H.), there is $\bar{p}_i(\cdot|\cdot) \in \Delta^{\bar{H}_i}(S_{-i})$ such that \bar{s}_i is fully $(\bar{p}_i(\cdot|\cdot), \bar{\phi}_i)$ -unimprovable and

$$\text{Supp} \bar{p}_i(\cdot|\{\emptyset\}) \subseteq \left(F_{\bar{s}_i}^{-1}(F_i(\bar{s})) \cap \text{symSCE}^k(\Gamma, f, \bar{\phi})_{\bar{s}_i} \right) \stackrel{(\text{I.H.})}{\subseteq} \left(F_{\bar{s}_i}^{-1}(F_i(\bar{s})) \cap \text{symSCE}^k(\Gamma, f, \phi)_{\bar{s}_i} \right).$$

We construct $p_i(\cdot|\cdot) \in \Delta^{\bar{H}_i}(\Sigma_{-i})$ such that $p_i(\cdot|h) = \bar{p}_i(\cdot|h)$ for each $h \in \bar{H}_i(\bar{p}_i(\cdot|\{\emptyset\}))$ and \bar{s}_i is fully $(p_i(\cdot|\cdot), \phi_i)$ -unimprovable. Since

$$\text{Supp}\bar{p}_i(\cdot|\{\emptyset\}) = \text{Supp}p_i(\cdot|\{\emptyset\}) \subseteq \left(F_{\bar{s}_i}^{-1}(F_i(\bar{s})) \cap \text{symSCE}^k(\Gamma, f, \phi)_{\bar{s}_i}\right)$$

and the construction holds for each i , this implies that $\bar{s} \in \text{symSCE}^{k+1}(\Gamma, f, \phi)$.

(*Construction of $p_i(\cdot|\cdot)$*) Since i moves at most once on every path, $H_i(s_i) = H_i$; furthermore, the value of any action $a_i \in A_i(h)$ ($h \in H_i$) is independent of i 's strategy. To construct $p_i(\cdot|\cdot)$, keep the same initial belief as \bar{p}_i : $p_i(\cdot|\{\emptyset\}) = \bar{p}_i(\cdot|\{\emptyset\})$. By symmetry (pure equilibrium), no chance moves, confirmed beliefs and ex-post perfect recall, there is a unique $h \in H_i(\bar{p}_i(\cdot|\{\emptyset\})) = H_i(p_i(\cdot|\{\emptyset\}))$ that i expects to reach with probability 1 under \bar{p}_i . Thus, let $p_i(\cdot|h) = \bar{p}_i(\cdot|\{\emptyset\}) = \bar{p}_i(\cdot|h)$ for $h \in H_i(\bar{p}_i(\cdot|\{\emptyset\}))$ even if $h \neq \{\emptyset\}$. It follows that for the unique $h \in H_i(p_i(\cdot|\{\emptyset\}))$ and for every $a_i \in A_i(h)$,

$$\begin{aligned} V_i(\bar{s}_{i,h}|h; p_i, \phi_i) &\stackrel{\text{(L4)}}{=} U_i(\bar{s}) \\ &\stackrel{\text{(L4)}}{=} V_i(\bar{s}_{i,h}|h; p_i, \bar{\phi}_i) \\ &\stackrel{\text{(uprv.)}}{\geq} V_i(a_i|h; p_i, \bar{\phi}_i) \\ &\stackrel{\text{(L5)}}{\geq} V_i(a_i|h; p_i, \phi_i), \end{aligned}$$

where the equalities follow from $h \in H_i(\bar{p}_i(\cdot|\{\emptyset\})) = H_i(p_i(\cdot|\{\emptyset\}))$ and (given the observability of payoffs) Lemma 4, the first inequality follows from $(p_i(\cdot|\{\emptyset\}), \bar{\phi}_i)$ -unimprovability of \bar{s}_i (uprv.) and the second one from Lemma 5.

Now, consider any $h \in H_i \setminus H_i(\bar{p}_i(\cdot|\{\emptyset\}))$. Since \bar{s}_i is fully $(\bar{p}_i(\cdot|\cdot), \bar{\phi}_i)$ -unimprovable we have

$$\bar{s}_{i,h} \in \arg \max_{a_i \in A_i(h)} V_i(a_i|h; \bar{p}_i, \bar{\phi}_i) = \arg \max_{a_i \in A_i(h)} \mathbb{E}_{\bar{p}_i(\cdot|h)} [\bar{\phi}_i \circ U_{h,a_i}].$$

By Lemma 10 there exists some $p_{i,h} \in \Delta(S_{-i}(h))$ such that

$$\bar{s}_{i,h} \in \arg \max_{a_i \in A_i(h)} \mathbb{E}_{p_{i,h}} [\varphi \circ \bar{\phi}_i \circ U_{h,a_i}] = \arg \max_{a_i \in A_i(h)} \mathbb{E}_{p_{i,h}} [\phi_i \circ U_{h,a_i}].$$

Let $p_i(\cdot|h) = p_{i,h}$. By the one-move assumption, h has no strict predecessors or followers in H_i . Then, the array $(p_i(\cdot|h))_{h \in \bar{H}_i}$ is a CPS. By construction \bar{s}_i is fully $(p_i(\cdot|\cdot), \phi_i)$ -unimprovable. \square

We conclude that $\bar{s} \in \text{symSCE}^{k+1}(\Gamma, f, \phi)$. \blacksquare

@[**Proof modified: Check carefully**]@

Proof of Theorem 12 Throughout the proof, for each $s = (s_j)_{j \in I}$, let $[s] = \times_{j \in I} [s_j]$ and $[s_{-i}] = \times_{j \neq i} [s_j]$. To prove the result, we will show that for every $k \in \mathbb{N}$ and $\bar{s} \in \text{symSCE}^k(\Gamma, f, \bar{\phi})$, there exists $s \in \text{symSCE}^k(\Gamma, f, \phi)$ with $s \in [\bar{s}]$. Then the claim follows from Lemma 2. The statement is trivially true for $k = 0$.

Suppose, by way of induction, that the statement is true for k . Let

$$\bar{s} = (\bar{s}_i)_{i \in I} \in \text{symSCE}^{k+1}(\Gamma, f, \bar{\phi}) \subseteq \text{symSCE}^k(\Gamma, f, \bar{\phi})$$

(the inclusion holds by definition). By the inductive hypothesis, for every $i \in I$, there exists some $s^* = (s_i^*)_{i \in I} \in \text{symSCE}^k(\Gamma, f, \phi)$ such that $s^* \in [\bar{s}]$. Since $\bar{s} \in \text{symSCE}^{k+1}(\Gamma, f, \bar{\phi})$, for every $i \in I$ there is some CPS $\bar{p}_i(\cdot|\cdot) \in \Delta^{\bar{H}_i}(S_{-i})$ such that \bar{s}_i is fully $(\bar{p}_i(\cdot|\cdot), \bar{\phi}_i)$ -unimprovable and

$$\text{Supp}\bar{p}_i(\cdot|\{\emptyset\}) \subseteq \left(F_{\bar{s}_i}^{-1}(F_i(\bar{s})) \cap \text{symSCE}^k(\Gamma, f, \bar{\phi})_{\bar{s}_i} \right).$$

It can be checked that, for any $s = (s_i, s_{-i}) \in \text{symSCE}^k(\Gamma, f, \phi)$, if $s_i \in [s_i^*]$, then $(s_i^*, s_{-i}) \in \text{symSCE}^k(\Gamma, f, \phi)$ as well. Then, by the inductive hypothesis, for every $s_{-i} \in \text{symSCE}^k(\Gamma, f, \bar{\phi})_{\bar{s}_i}$, there exists $s'_{-i} \in \text{symSCE}^k(\Gamma, f, \phi)_{s_i^*}$ such that $s'_{-i} \in [s_{-i}]$. Moreover, since strategic feedback depends only on the realization equivalence classes of strategies, for each $s_{-i} \in F_{\bar{s}_i}^{-1}(F_i(\bar{s}))$ and $s'_{-i} \in [s_{-i}]$, we have $s'_{-i} \in F_{s_i^*}^{-1}(F_i(s^*))$. So, we can construct a belief $p_i^* \in \Delta(S_{-i})$ with

$$\text{Supp}p_i^* \subseteq \left(F_{s_i^*}^{-1}(F_i(s^*)) \cap \text{symSCE}^k(\Gamma, f, \phi)_{s_i^*} \right)$$

such that for each $\hat{s}_{-i} \in \text{Supp}p_i^*$,

$$p_i^*([\hat{s}_{-i}]) = \bar{p}_i([\hat{s}_{-i}]|\{\emptyset\}). \quad (17)$$

Since $\bar{p}_i(\cdot|\{\emptyset\})$ is the predictive probability isomorphic to a belief over Dirac models, by Corollary 14 player i 's preferences are dynamically consistent. Thus,⁴⁵

$$\bar{s}_i \in \arg \max_{s_i \in S_i} \sum_{s_{-i}} \bar{\phi}_i(U_i(s_i, s_{-i}) \bar{p}_i(s_{-i}|\{\emptyset\})).$$

But then, by eq. 17, we also have

$$\bar{s}_i \in \arg \max_{s_i \in S_i} \sum_{s_{-i}} \bar{\phi}_i(U_i(s_i, s_{-i}) p_i^*(s_{-i})). \quad (18)$$

Consider the decision problem $(\hat{A}, \hat{S}, \hat{u})$ where $\hat{A} = S_i$, $\hat{S} = F_{s_i^*}^{-1}(F_i(s^*)) \cap \text{symSCE}^k(\Gamma, f, \phi)_{s_i^*}$, and \hat{u} is the restriction of U_i on $\hat{A} \times \hat{S}$. By eq. 18, \bar{s}_i is a $\bar{\phi}_i$ -justified by $p_i^* \in \Delta(\hat{S})$. By Lemma 10, there exists $p_i \in \Delta(\hat{S})$ that ϕ_i -justifies \bar{s}_i . Let $p_i(\cdot|\{\emptyset\}) = p_i$. Then

$$p_i(\cdot|\{\emptyset\}) \in \Delta \left(F_{s_i^*}^{-1}(F_i(s^*)) \cap \text{symSCE}^k(\Gamma, f, \phi)_{s_i^*} \right)$$

⁴⁵ $\bar{\phi}_i^{-1}$ and ϕ_i^{-1} are omitted in front of the next maximands, because they are strictly increasing transformations, and thus they do not affect the maximizers.

and

$$\bar{s}_i \in \arg \max_{s_i \in S_i} \sum_{s_{-i}} \phi_i(U_i(s_i, s_{-i}) p_i(s_{-i} | \{\emptyset\})).$$

By dynamic consistency of player i 's preferences and standard arguments, this implies that for every $s_i \in [\bar{s}_i]$, for every $h' \in H_i(p_i(\cdot | \{\emptyset\})) \cap H_i(\bar{s}_i)$,

$$s_{i,h'} \in \arg \max_{a_i \in A_i(h')} \sum_{s_{-i}} \phi_i(U_i((s_i | h', a_i), s_{-i} | h') p_i(s_{-i} | h')),$$

where $p_i(s_{-i} | h')$ is derived by conditioning.

Next, fix $h \in H_i(\bar{s}_i) \setminus H_i(p_i(\cdot | \{\emptyset\}))$ with $h' \in H_i(p_i(\cdot | \{\emptyset\}))$ for all $h' \prec h$. By dynamic consistency, we can frame the continuation of the game as a (static) decision problem $(\hat{A}, \hat{S}, \hat{u})$, with $\hat{A} = S_i(h)$, $\hat{S} = S_{-i}(h)$ and \hat{u} is the restriction of $\bar{\phi}_i \circ U_i$ on $\hat{A} \times \hat{S} = S_i(h) \times S_{-i}(h)$. By the same argument as above, since \bar{s}_i is $\bar{\phi}_i$ -justified by $\bar{p}_i(\cdot | h)$, Lemma 10 implies that \bar{s}_i is ϕ_i -justified by some $p_i(\cdot | h) \in \Delta(S_{-i}(h))$. By dynamic consistency of player i 's preferences and standard arguments, this implies that for every $s_i \in [\bar{s}_i]$, for every $h' \in H_i(p_i(\cdot | h)) \cap H_i(\bar{s}_i)$ with $h \preceq h'$,

$$s_{i,h'} \in \arg \max_{a_i \in A_i(h')} \sum_{s_{-i}} \phi_i(U_i((s_i | h', a_i), s_{-i} | h') p_i(s_{-i} | h')),$$

where $p_i(s_{-i} | h')$ is derived from $p_i(s_{-i} | h)$ by conditioning.

Repeating iteratively the operation at all information sets that are not reached with positive probability under the probability measures already constructed, one can finally construct a CPS $\hat{p}_i(\cdot | \cdot) \in \Delta^{\bar{H}_i}(\Sigma_{-i})$ such that $\hat{p}_i(\cdot | h) = p_i(\cdot | h)$ for each $h \in H_i(\bar{s}_i)$ with $p_i(S_{-i}(h) | h') = 0$ for all $h' \prec h$. Since every $s_i \in [\bar{s}_i]$ is $(\hat{p}_i(\cdot | \cdot), \phi_i)$ -unimprovable at every $h \in H_i(\bar{s}_i)$, there exists a fully $(\hat{p}_i(\cdot | \cdot), \phi_i)$ -unimprovable $s_i \in [\bar{s}_i] = [s_i^*]$. Let $s = \times_{i \in I} s_i$. By construction,

$$\hat{p}_i(\cdot | \{\emptyset\}) \in \Delta \left(F_{s_i^*}^{-1}(F_i(s^*)) \cap \text{symSCE}^k(\Gamma, f, \phi)_{s_i^*} \right).$$

By $s_i \in [s_i^*]$ and $s \in [s^*]$, $F_{s_i^*}^{-1}(F_i(s^*)) = F_{s_i}^{-1}(F_i(s))$. So, $s \in \text{symSCE}^1(\Gamma, f, \phi)$. If $k > 0$, for all $s'_{-i} \in S_{-i}$, $(s_i^*, s'_{-i}) \in \text{symSCE}^1(\Gamma, f, \phi)$ if and only if $(s_i, s'_{-i}) \in \text{symSCE}^1(\Gamma, f, \phi)$. So, $\text{symSCE}^1(\Gamma, f, \phi)_{s_i^*} = \text{symSCE}^1(\Gamma, f, \phi)_{s_i}$. But then, $\hat{p}_i(\cdot | \{\emptyset\}) \in \Delta \left(F_{s_i}^{-1}(F_i(s)) \cap \text{symSCE}^1(\Gamma, f, \phi)_{s_i} \right)$ and $s \in \text{symSCE}^2(\Gamma, f, \phi)$. Inductively,

$$\hat{p}_i(\cdot | \{\emptyset\}) \in \Delta \left(F_{s_i}^{-1}(F_i(s)) \cap \text{symSCE}^k(\Gamma, f, \phi)_{s_i} \right),$$

and so $s \in \text{symSCE}^{k+1}(\Gamma, f, \phi)$. ■

10.6.3 Monotonicity of the RSCE correspondence

We first use Lemma 10 to prove a preliminary monotonicity result for the set of fully unimprovable SCEs.

Lemma 11 *Fix two games with observable payoffs where no player moves more than once, (Γ, f, ϕ) and $(\Gamma, f, \bar{\phi})$, so that (Γ, f, ϕ) features more ambiguity aversion than $(\Gamma, f, \bar{\phi})$. Then $SCE^1(\Gamma, f, \bar{\phi}) \subseteq SCE^1(\Gamma, f, \phi)$.*

Proof of Lemma 11 Let $\bar{\sigma} \in SCE^1(\Gamma, f, \bar{\phi})$. Fix $i \in I$, $s_i \in \text{Supp} \bar{\sigma}_i$ and let $\bar{\mu}_{s_i}(\cdot|\cdot) \in \Delta^{\bar{H}_i}(\Sigma_{-i})$ be such that s_i is fully $(\bar{\mu}_{s_i}(\cdot|\cdot), \bar{\phi}_i)$ -unimprovable and

$$\text{Supp} \bar{\mu}_{s_i}(\cdot|\{\emptyset\}) \subseteq \hat{\Sigma}_{-i}(s_i, \bar{\sigma}_{-i}) \stackrel{(\text{def.})}{=} \left\{ \sigma_{-i} \in \Sigma_{-i} : \hat{F}_i(s_i, \sigma_{-i}) = \hat{F}_i(s_i, \bar{\sigma}_{-i}) \right\}.$$

We will construct $\mu_{s_i}(\cdot|\cdot) \in \Delta^{\bar{H}_i}(\Sigma_{-i})$ such that $\mu_{s_i}(\cdot|h) = \bar{\mu}_{s_i}(\cdot|h)$ for each $h \in \bar{H}_i(\bar{\mu}_{s_i}(\cdot|\{\emptyset\}))$, and s_i is fully $(\mu_{s_i}(\cdot|\cdot), \phi_i)$ -unimprovable. Since

$$\text{Supp} \bar{\mu}_{s_i}(\cdot|\{\emptyset\}) = \text{Supp} \mu_{s_i}(\cdot|\{\emptyset\}) \subseteq \hat{\Sigma}_{-i}(s_i, \bar{\sigma}_{-i}),$$

this implies that $\bar{\sigma} \in SCE^1(\Gamma, f, \phi)$.

Recall that $H_i(s_i) = H_i$ because i moves at most once on every path. To construct $\mu_{s_i}(\cdot|\cdot)$, keep the same prior belief: $\mu_{s_i}(\cdot|\{\emptyset\}) = \bar{\mu}_{s_i}(\cdot|\{\emptyset\})$. For each $h \in H_i(\bar{\mu}_{s_i}(\cdot|\{\emptyset\})) = H_i(\mu_{s_i}(\cdot|\{\emptyset\}))$ define $\mu_{s_i}(\cdot|h)$ by conditioning. Hence, for each $h \in \bar{H}_i(\bar{\mu}_{s_i}(\cdot|\{\emptyset\}))$, $\mu_{s_i}(\cdot|h) = \bar{\mu}_{s_i}(\cdot|h)$. Thus, for every $a_i \in A_i(h)$,

$$\begin{aligned} V_i(s_{i,h}|h; \mu_{s_i}, \phi_i) &\stackrel{(\text{L4})}{=} U_i(s_i, \bar{\sigma}_{-i}|h) \\ &\stackrel{(\text{L4})}{=} V_i(s_{i,h}|h; \mu_{s_i}, \bar{\phi}_i) \\ &\stackrel{(\text{uprv.})}{\geq} V_i(a_i|h; \mu_{s_i}, \bar{\phi}_i) \\ &\stackrel{(\text{L5})}{\geq} V_i(a_i|h; \mu_{s_i}, \phi_i), \end{aligned}$$

where the equalities follow from Lemma 4, the first inequality from full $(\bar{\mu}_{s_i}(\cdot|\cdot), \bar{\phi}_i)$ -unimprovability of s_i (uprv.), and the second one from Lemma 5.

Now, consider any $h \in H_i \setminus H_i(\bar{\mu}_{s_i}(\cdot|\{\emptyset\}))$. Since s_i is fully $(\bar{\mu}_{s_i}(\cdot|\cdot), \bar{\phi}_i)$ -unimprovable and $\bar{\mu}_{s_i}(\Sigma_{-i}(h)|h) = 1$ by condition (1) of Definition 6, we have

$$\begin{aligned} s_{i,h} &\in \arg \max_{a_i \in A_i(h)} \int_{\Sigma_{-i}} \bar{\phi}_i(U_i((s_i|h, a_i), \sigma_{-i}|h)) \bar{\mu}_{s_i}(d\sigma_{-i}|h) \\ &= \arg \max_{a_i \in A_i(h)} \int_{\Sigma_{-i}(h)} \bar{\phi}_i \left(\sum_{s_{-i} \in S_{-i}(h)} \sigma_{-i}(s_{-i}|h) v_i(\gamma(\zeta(a_i, s_{-i}))) \right) \bar{\mu}_{s_i}(d\sigma_{-i}|h), \end{aligned}$$

where we abuse notation and write $\zeta(a_i, s_{-i})$, because i does not move before h , nor after h , hence the terminal node reached depends only on $s_{-i} \in S_{-i}(h)$ and the action chosen by i at h . Next, we consider the problem of choice under uncertainty $(\hat{A}, \hat{S}, \hat{u})$ where $A = A_i(h)$, $\hat{S} = S_{-i}(h)$, $\hat{u}(a_i, s_{-i}) = v_i(\gamma(\zeta(a_i, s_{-i})))$. Recall that we assumed $S_{-i}(h) = \times_{j \neq i} S_j(h)$ (observable deviators), therefore it makes sense to consider the following compact set of product distributions

$$\Sigma_{-i|h} = \{\sigma_{-i} \in \Sigma_{-i} : \sigma_{-i}(\times_{j \neq i} S_j(h)) = 1\} \subseteq \Sigma_{-i}(h).$$

The following map associates each $\sigma_{-i} \in \Sigma_{-i}(h)$ with the corresponding updated distribution $\sigma_{-i}(\cdot|h)$ defined in (1):

$$\begin{aligned} \varsigma_{h,-i} : \Sigma_{-i}(h) &\rightarrow \Sigma_{-i|h} \\ \sigma_{-i} &\mapsto \frac{\sigma_{-i}(\cdot)}{\sigma_{-i}(\times_{j \neq i} S_j(h))}. \end{aligned}$$

Note that $\Sigma_{-i}(h)$ is a (relatively) open subset of the Polish space Σ_{-i} and $\varsigma_{h,-i}$ is continuous. Furthermore, the restriction $\varsigma_{h,-i}|_{\Sigma_{-i|h}}$ is the identity on $\Sigma_{-i|h}$, because $\sigma_{-i}(\times_{j \neq i} S_j(h)) = 1$ implies

$$\sigma_{-i}(s_{-i}|h) = \frac{\sigma_{-i}(s_{-i})}{\sigma_{-i}(\times_{j \neq i} S_j(h))} = \sigma_{-i}(s_{-i})$$

for each s_{-i} . Therefore $\varsigma_{h,-i}$ is also onto and $\varsigma_{h,-i}$ is a measurable surjection that yields the *onto* pushforward map

$$\begin{aligned} \hat{\varsigma}_{h,-i} : \Delta(\Sigma_{-i}(h)) &\rightarrow \Delta(\Sigma_{-i|h}), \\ \mu_i &\mapsto \mu_i \circ \varsigma_{h,-i}^{-1}. \end{aligned}$$

Let $\bar{\mu}_{s_i|h} = \bar{\mu}_{s_i}(\cdot|h) \circ \varsigma_{h,-i}^{-1} \in \Delta(\Sigma_{-i|h})$, that is,

$$\forall E \in \mathcal{B}(\Sigma_{-i|h}), \quad \bar{\mu}_{s_i,h}(E) = \bar{\mu}_{s_i}(\varsigma_{h,-i}^{-1}(E)|h).$$

For every $a_i \in A_i(h)$, we have:

$$\begin{aligned} &\int_{\Sigma_{-i}(h)} \bar{\phi}_i \left(\sum_{s_{-i} \in S_{-i}(h)} \sigma_{-i}(s_{-i}|h) v_i(\gamma(\zeta(a_i, s_{-i}))) \right) \bar{\mu}_{s_i}(\mathrm{d}\sigma_{-i}|h) = \\ &\int_{\Sigma_{-i|h}} \bar{\phi}_i \left(\sum_{s_{-i} \in S_{-i}(h)} \sigma_{-i}(s_{-i}|h) v_i(\gamma(\zeta(a_i, s_{-i}))) \right) \bar{\mu}_{s_i|h}(\mathrm{d}\sigma_{-i}(\cdot|h)). \end{aligned}$$

Lemma 10, (??) and the above equality imply that $s_{i,h}$ is ϕ_i -justified by some belief $\mu_{i|h} \in \Delta(\Sigma_{-i|h})$, that is,

$$s_{i,h} \in \arg \max_{a_i \in A_i(h)} \int_{\Delta(\Sigma_{-i|h})} \phi_i \left(\sum_{s_{-i} \in S_{-i}(h)} \sigma_{-i}(s_{-i}|h) v_i(\gamma(\zeta(a_i, s_{-i}))) \right) \mu_{i|h}(\mathrm{d}\sigma_{-i}(\cdot|h)).$$

Now go back to a belief on $\Sigma_{-i}(h)$: Since the pushforward map $\hat{\zeta}_{h,-i}$ is onto we can find some belief $\mu_{s_i}(\cdot|h) \in \hat{\zeta}_{h,-i}^{-1}(\mu_{i|h}) \subseteq \Delta(\Sigma_{-i}(h))$ such that

$$\begin{aligned} s_{i,h} &\in \arg \max_{a_i \in A_i(h)} \int_{\Delta(\Sigma_{-i}(h))} \phi_i \left(\sum_{s_{-i} \in S_{-i}(h)} \sigma_{-i}(s_{-i}|h) v_i(\gamma(\zeta(a_i, s_{-i}))) \right) \mu_{s_i}(d\hat{\sigma}_{-i}|h) \\ &= \arg \max_{a_i \in A_i(h)} \int_{\Sigma_{-i}} \phi_i(U_i((s_i|h, a_i), \sigma_{-i}|h)) \mu_{s_i}(d\sigma_{-i}|h). \end{aligned}$$

We can do this for every off-path information set $h \in H_i \setminus H_i(\bar{\mu}_{s_i}(\cdot|\{\emptyset\}))$ and obtain an array $(\mu_{s_i}(\cdot|h))_{h \in \bar{H}_i} \in \times_{h \in \bar{H}_i} \Delta(\Sigma_{-i}(h))$ (recall that μ_{s_i} coincides with $\bar{\mu}_{s_i}$ on $H_i(\bar{\mu}_{s_i}(\cdot|\{\emptyset\}))$) that satisfies condition (1) of the definition of CPS by construction. Since every player moves at most once on every path, condition (2) trivially holds. Finally, by construction, s_i is fully $(\mu_i(\cdot), \phi_i)$ -unimprovable. ■

Proof of Theorem 13 Lemma 11 shows that the result holds for $k = 1$. Suppose by way of induction that $SCE^k(\Gamma, f, \bar{\phi}) \subseteq SCE^k(\Gamma, f, \phi)$. Fix $\sigma \in SCE^{k+1}(\Gamma, f, \bar{\phi})$, $i \in I$, $s_i \in \text{Supp}\sigma_i$ arbitrarily and let $\bar{\mu}_{s_i}(\cdot|\cdot) \in \Delta^{\bar{H}_i}(\Sigma_{-i})$ be a confirmed CPS that $\bar{\phi}_i$ -justifies s_i , that is, s_i is fully $(\bar{\mu}_{s_i}(\cdot|\cdot), \bar{\phi})$ -unimprovable and

$$\bar{\mu}_{s_i} \left(\hat{\Sigma}_{-i}(s_i, \sigma_{-i}) \cap \text{proj}_{\Sigma_{-i}} SCE^k(\Gamma, f, \bar{\phi}) | \{\emptyset\} \right) = 1.$$

Since $\sigma \in SCE^{k+1}(\Gamma, f, \bar{\phi}) \subseteq SCE^k(\Gamma, f, \bar{\phi})$, the inductive hypothesis implies $\sigma \in SCE^k(\Gamma, f, \phi)$, so there is $\hat{\mu}_{s_i}(\cdot|\cdot) \in \Delta^{\hat{H}_i}(\Sigma_{-i})$ such that s_i is fully $(\hat{\mu}_{s_i}(\cdot|\cdot), \phi_i)$ -unimprovable. Let $\mu_{s_i}(\cdot|\cdot) \in \Delta^{\bar{H}_i}(\Sigma_{-i})$ be such that $\mu_{s_i}(\cdot|h) = \bar{\mu}_{s_i}(\cdot|h)$ for each $h \in \bar{H}_i(\bar{\mu}_{s_i}(\cdot|\{\emptyset\}))$ and $\mu_{s_i}(\cdot|h) = \hat{\mu}_{s_i}(\cdot|h)$ for each $h \in H_i \setminus \bar{H}_i(\bar{\mu}_{s_i}(\cdot|\{\emptyset\}))$. By same argument employed in the proof of Lemma 11, $\mu_{s_i}(\cdot|\cdot)$ is a CPS and s_i is $(\mu_{s_i}(\cdot|\cdot), \phi_i)$ -unimprovable. Moreover, by construction and the inductive hypothesis,

$$\begin{aligned} &\mu_{s_i} \left(\hat{\Sigma}_{-i}(s_i, \sigma_{-i}) \cap \text{proj}_{\Sigma_{-i}} SCE^k(\Gamma, f, \phi) | \{\emptyset\} \right) \\ &\geq \mu_{s_i} \left(\hat{\Sigma}_{-i}(s_i, \sigma_{-i}) \cap \text{proj}_{\Sigma_{-i}} SCE^k(\Gamma, f, \bar{\phi}) | \{\emptyset\} \right) \\ &= \bar{\mu}_{s_i} \left(\hat{\Sigma}_{-i}(s_i, \sigma_{-i}) \cap \text{proj}_{\Sigma_{-i}} SCE^k(\Gamma, f, \bar{\phi}) | \{\emptyset\} \right) = 1. \end{aligned}$$

Therefore s_i is ϕ_i -justified by a confirmed CPS that initially believes $\text{proj}_{\Sigma_{-i}} SCE^k(\Gamma, f, \phi)$. This holds for every $i \in I$ and $\text{Supp}\sigma_i$, thus, $\sigma \in SCE^{k+1}(\Gamma, f, \phi)$. ■

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